



Advancing Security and Privacy of Bluetooth IoTs via Formal Protocol Analysis

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Outline

- 1 Introduction
- 2 Background
- 3 Our Prior Works
- 4 Proposed Research

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What is Bluetooth Low Energy



Power Consumption: High
Communication Distance: Short (10+ m)



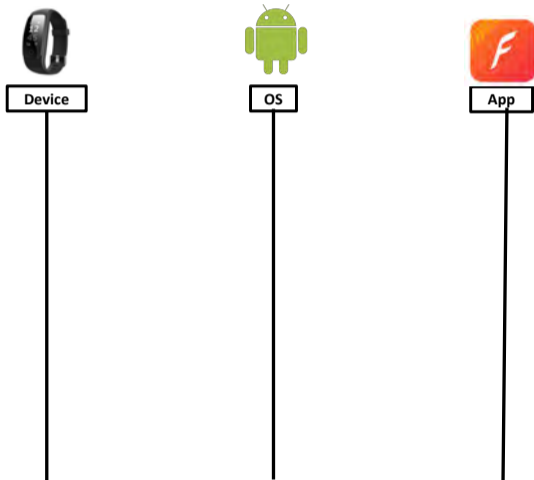
Power Consumption: Low
Communication Distance: Long (100+ m)



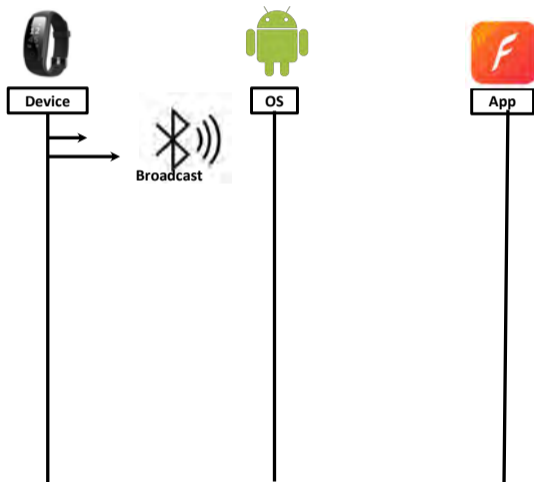
Bluetooth Low Energy Applications



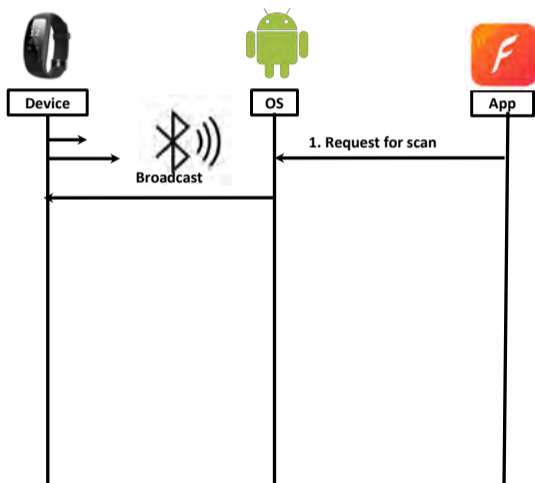
Bluetooth Low Energy Communication



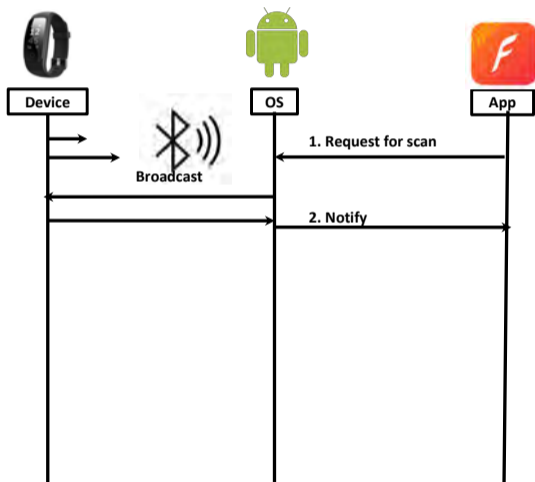
Bluetooth Low Energy Communication



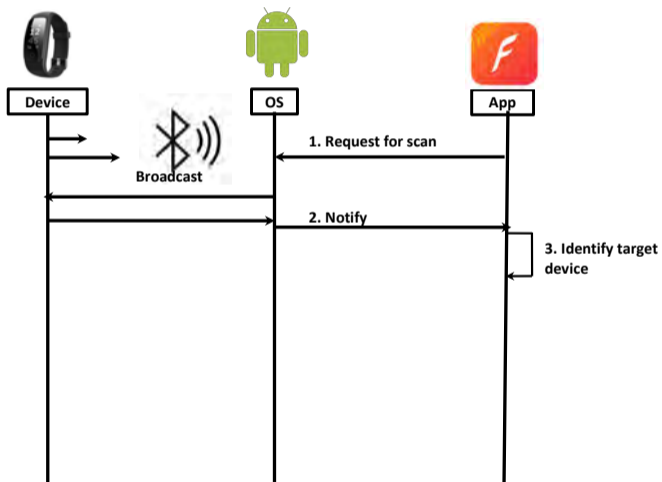
Bluetooth Low Energy Communication



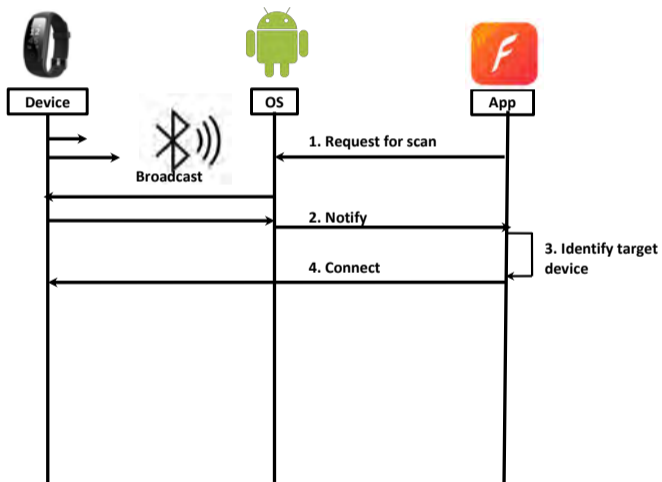
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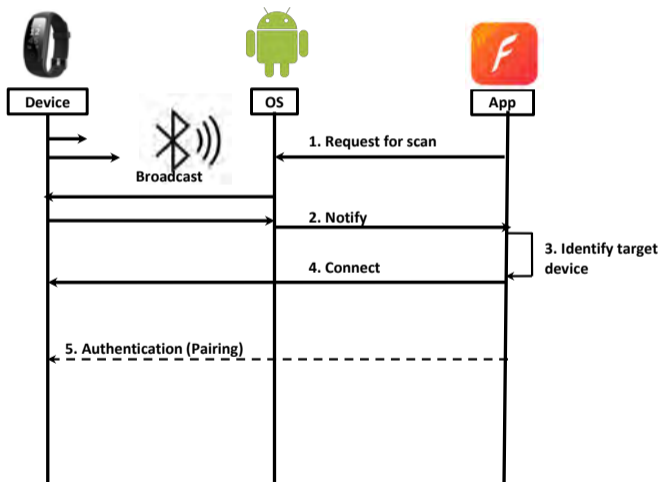
Bluetooth Low Energy Communication



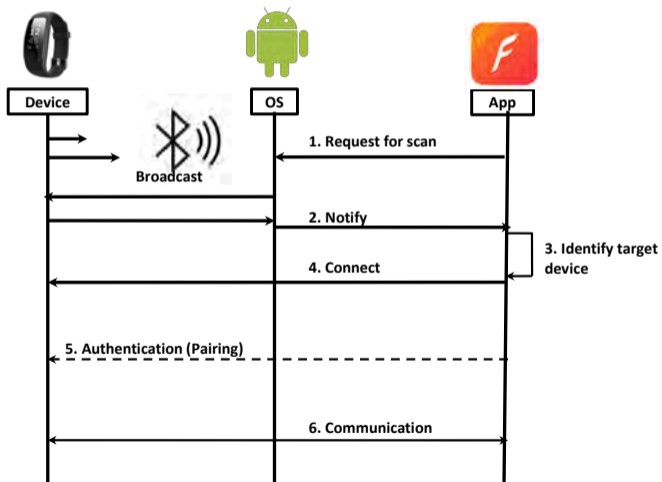
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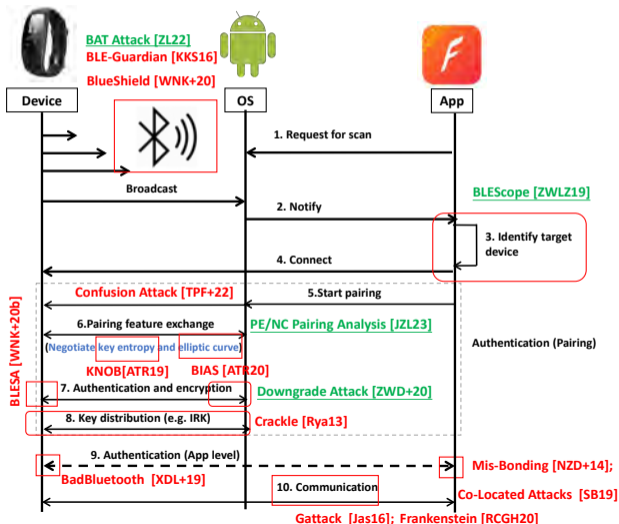
Bluetooth Low Energy Communication



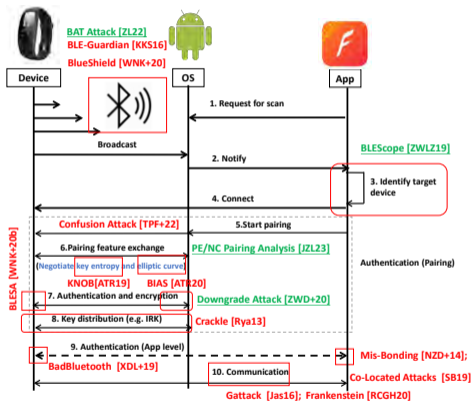
Bluetooth Low Energy Communication



Bluetooth Low Energy Communication



Bluetooth Low Energy Communication



The goal of this project is to systematically uncover the attacks via formal methods

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Mathematical Proof: A Simple Example

Assume a set of even integers

$$E = \{\dots, -2, 0, 2, 4, \dots\}$$

And a set of odd integers

$$O = \{\dots, -1, 1, 3, 5, \dots\}$$

Property (Lemma)

P = "Elements of E and O are distinct"

Mathematical Proof: A Simple Example

Math Algebra

$E = 2x$ for $x \in \text{Integer}$

$O = 2y + 1$ for $y \in \text{Integer}$

Proof by Contradiction

not P = "There exist some common elements among E and O"

Mathematical Proof: A Simple Example

Proof

Start with assuming for some x and y

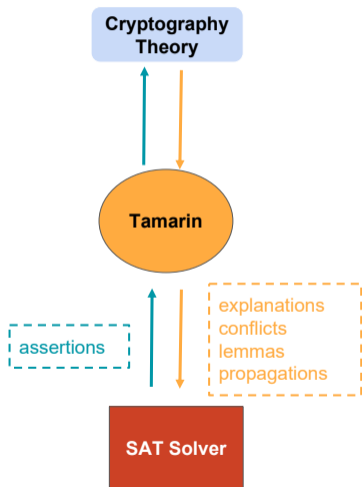
$2x = 2y + 1$ holds true

$\Rightarrow 2(x - y) = 1$

Mathematical Proof: A Simple Example

Possible Sub-Cases of X-Y	Generally known axioms (Generalized Constraint Resolution Methods)	applying axioms to $2(x-y) = 1$	Conclusion (Contradicting all not P system states)
0	0 is the only neutral difference of Integers	$2*0=1$	Contradict
negatives	all negatives differences of integers are -1 or less $\{\dots, -3, -2, -1\}$	$2*(-1 \text{ or less})= 1$	Contradict
positive	all positive differences of Integers are 1 or more $\{1, 2, 3, \dots\}$	$2*(1 \text{ or more})= 1$	Contradict

How does Protocol Verification (Tamarin) Work?



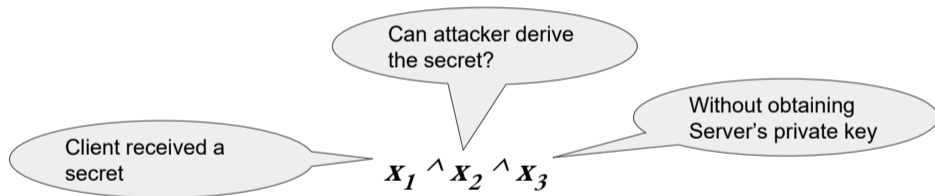
- $\text{Dec}(\text{Enc}(\text{msg}, \text{key})) = \text{msg}$
- $\text{Sign}(\text{msg}, \text{privKey}) = \text{Verify}(\text{msg}, \text{pubKey})$
- Adversary Replay
- ...

- Find next proof requirements
- Choose which SAT problems to solve first
- Convert problem algebra theory into SAT problems
- ...

$$X_1 \wedge X_2 \wedge X_3$$

- Does SAT problem have a solution?

How does Protocol Verification (Tamarin) Work?



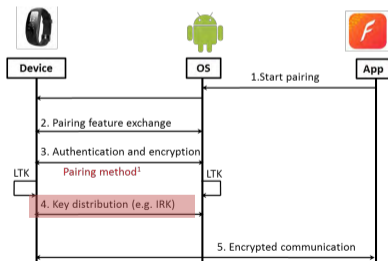
(Privacy) How to Avoid Being Tracked: MAC Address Randomization



Identity Resolving Key (irk_p)



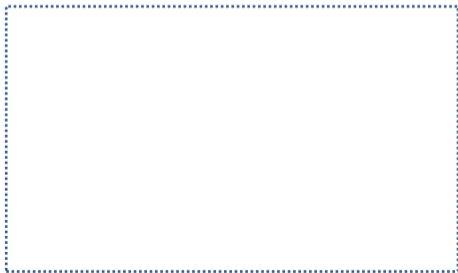
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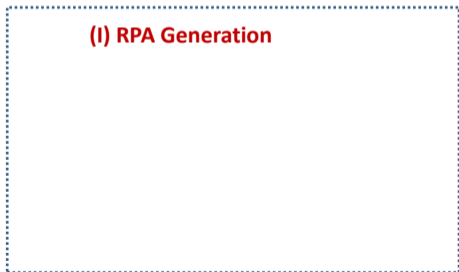


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(Privacy) How to Avoid Being Tracked: MAC Address Randomization



Identity Resolving Key (irk_p)



Identity Resolving Key (irk_c)

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Identity Resolving Key (irk_p)

(I) RPA Generation

$$rpa_p = prand_{24} || H_{24}(Prand_{24} || irk_p)$$



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(II) RPA Resolution

(Privacy) How to Avoid Being Tracked: MAC Address Randomization



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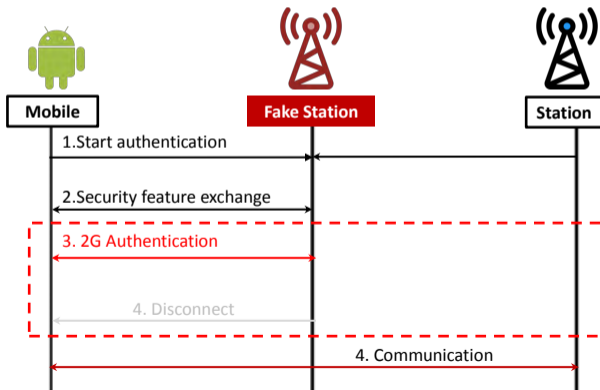
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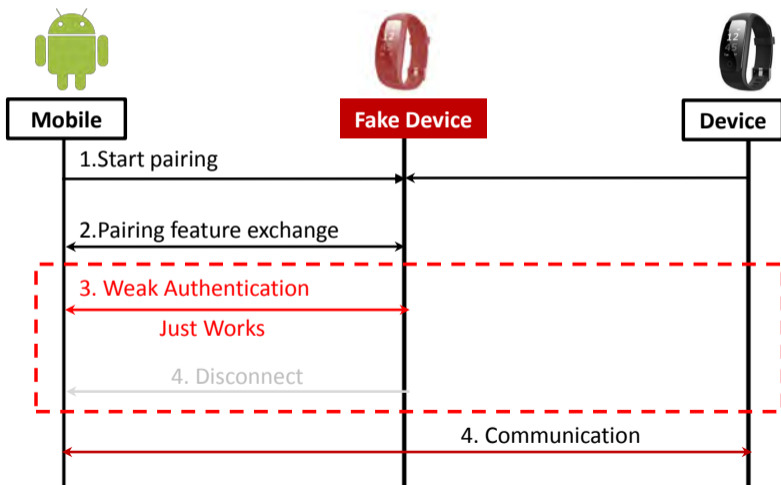
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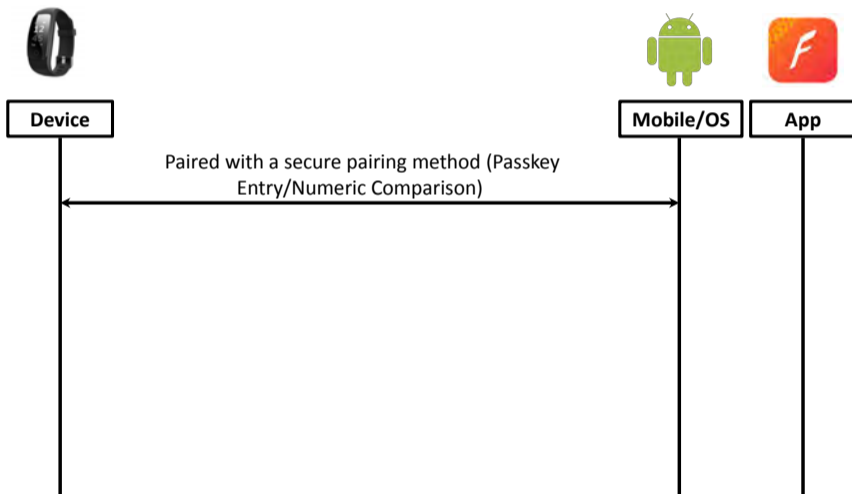
Our Downgrade Attacks against Bluetooth Low Energy [USENIX'20]



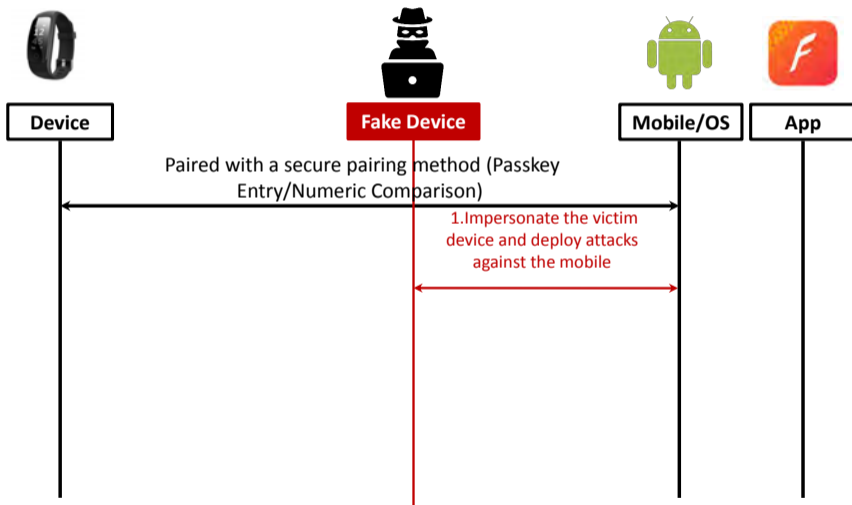
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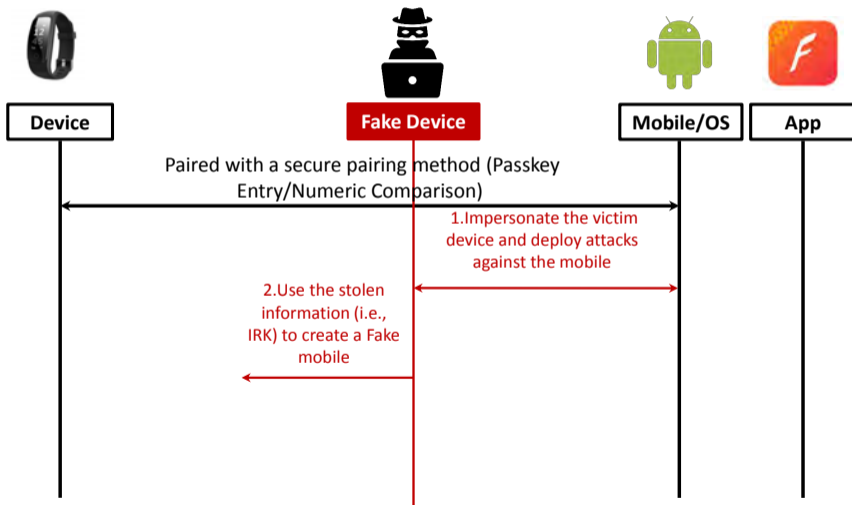
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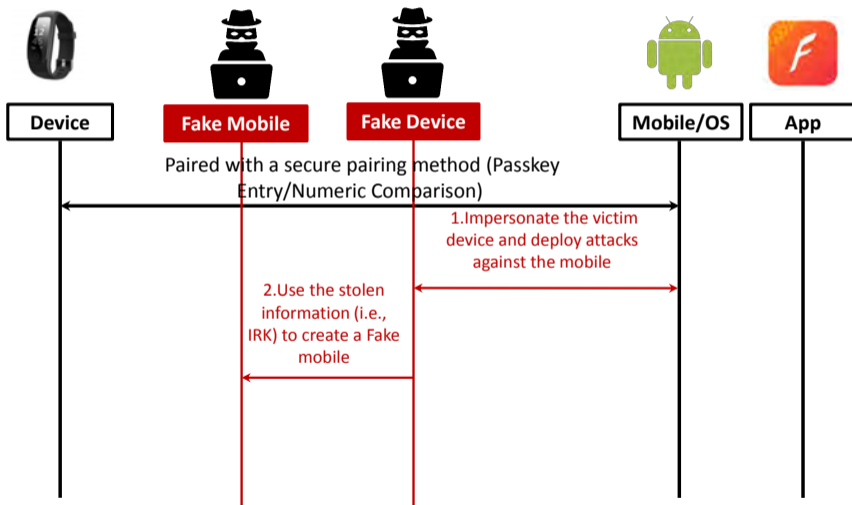
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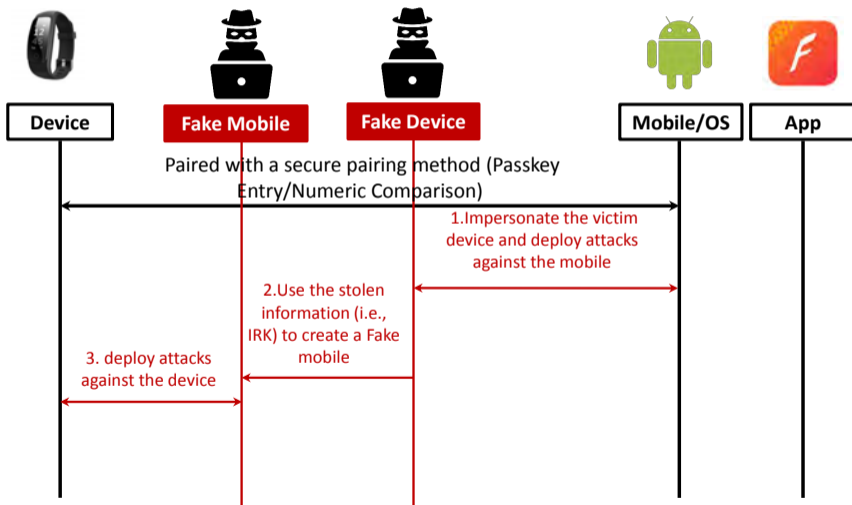
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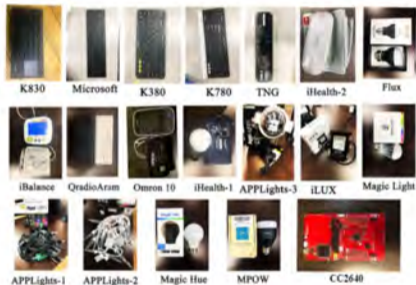
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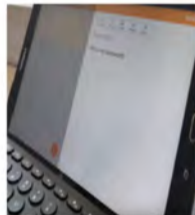
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Our Downgrade Attacks against Bluetooth Low Energy [USENIX'20]



The Tested BLE devices



User



Attacker

MITM attack against BLE keyboards



CVE-2020-9770

Our BAT Privacy Attacks [CCS'22]: Allowlist-based Side Channel



58:D7:8E:C7:8e:31

NO.	Time	Source	Destination	TYPE
1	00:00:04	58:D7:8E:C7:8e:31	Broadcast	ADV_IND

Our BAT Privacy Attacks [CCS'22]: Allowlist-based Side Channel



58:D7:8E:C7:8e:31



7e:D7:8E:C7:8e:51

NO.	Time	Source	Destination	TYPE
1	00:00:04	58:D7:8E:C7:8e:31	Broadcast	ADV_IND
2	00:00:08	7e:D7:8E:C7:8e:51	58:D7:8E:C7:8e:31	SCAN_REQ
3	00:00:12	58:D7:8E:C7:8e:31	Broadcast	SCAN_RSP

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3	00:00:12	58:D7:8E:C7:8e:31	Broadcast	SCAN_RSP
4	00:00:16	4f:b7:8E:C7:8e:38	58:D7:8E:C7:8e:31	SCAN_REQ
5	00:00:24	58:D7:8E:C7:8e:31	Broadcast	ADV_IND

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4	00:00:16	4f:b7:8E:C7:8e:38	58:D7:8E:C7:8e:31	SCAN_REQ
5	00:00:24	58:D7:8E:C7:8e:31	Broadcast	ADV_IND
.....				
200	00:15:08	73:D7:8E:C7:8e:45	58:D7:8E:C7:8e:31	SCAN_REQ
201	00:15:12	58:D7:8E:C7:8e:31	Broadcast	SCAN_RSP

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- ① Cache
- ② Timing
- ③ Power
- ④ Votage
- ⑤ Electromagnetic
- ⑥ Acoustic
- ⑦ Allow-list
- ⑧ ...

Our BAT Privacy Attacks [CCS'22]: MAC Address Replay



Identity Resolving Key (irk_p)



Identity Resolving Key (irk_c)

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rpa_p

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No Identity Resolving Key

RPA Replay (rpa'_p)

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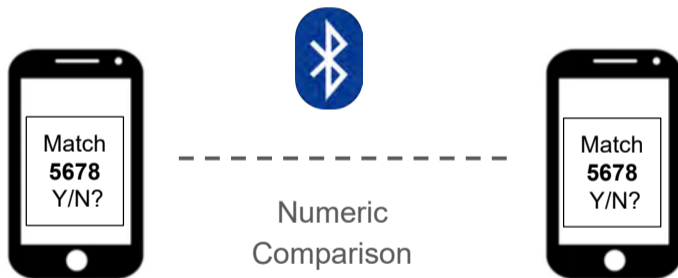


rpa_p

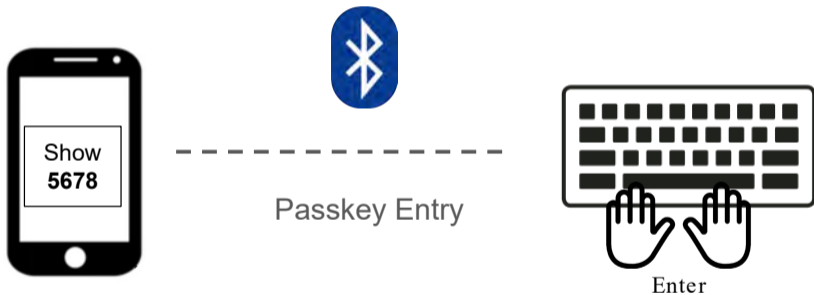


rpa'_p

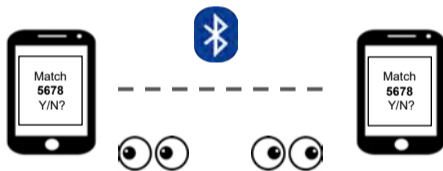
Using Protocol Verification to Identify Confusion Attacks [NDSS'23]



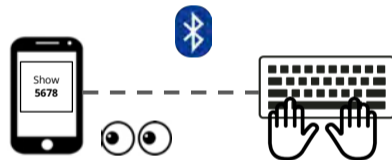
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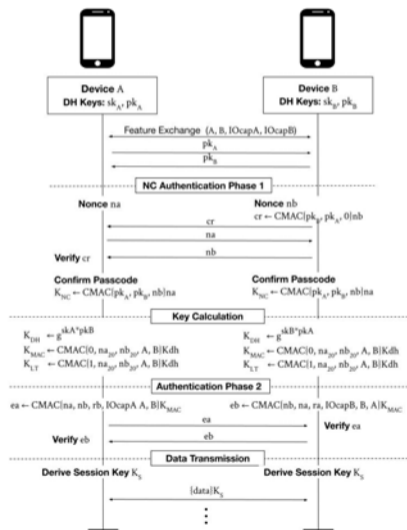
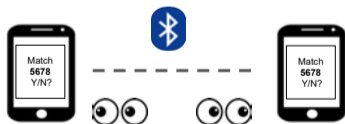


Numeric Comparison

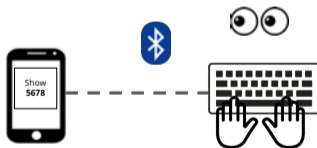


Passkey Entry

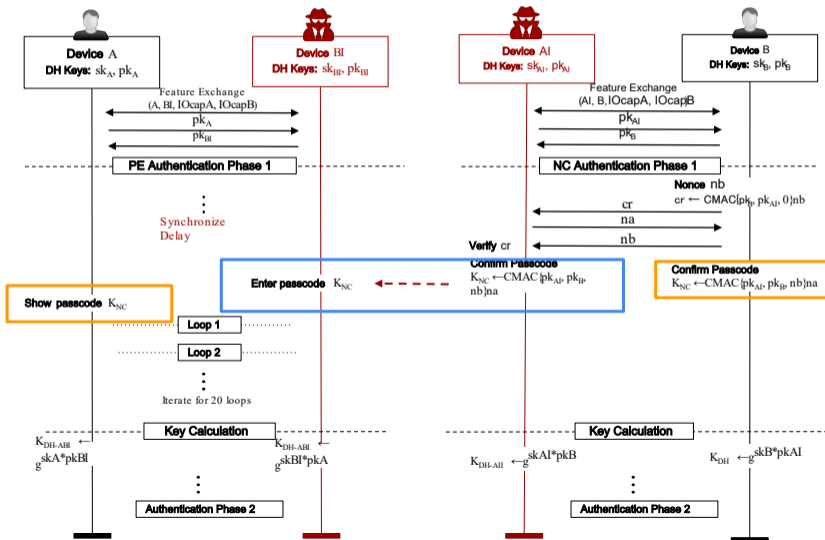
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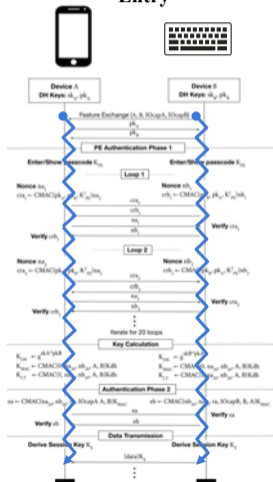


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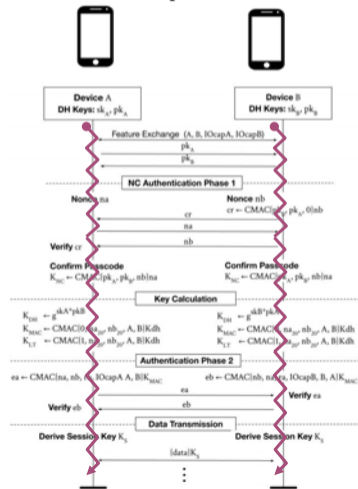


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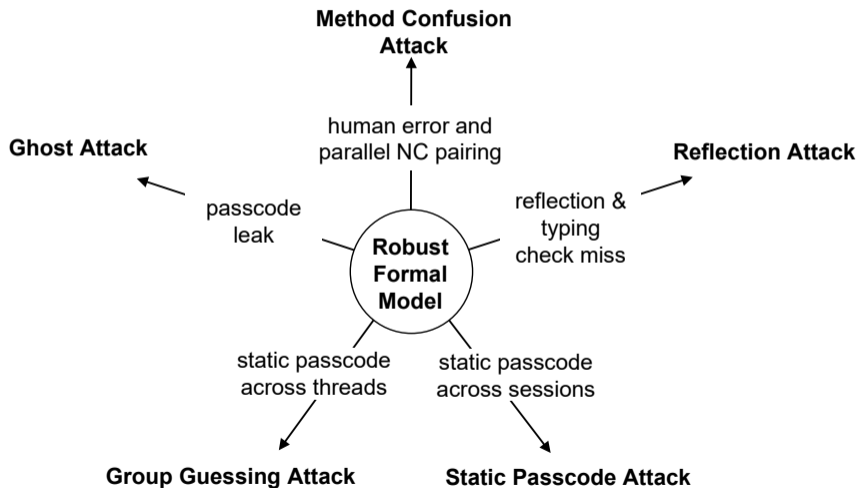
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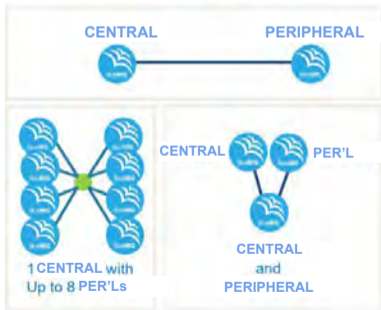


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Task 1: Developing a Formal Model for Full Spectrum of the Protocols

PAIRING one-to-one



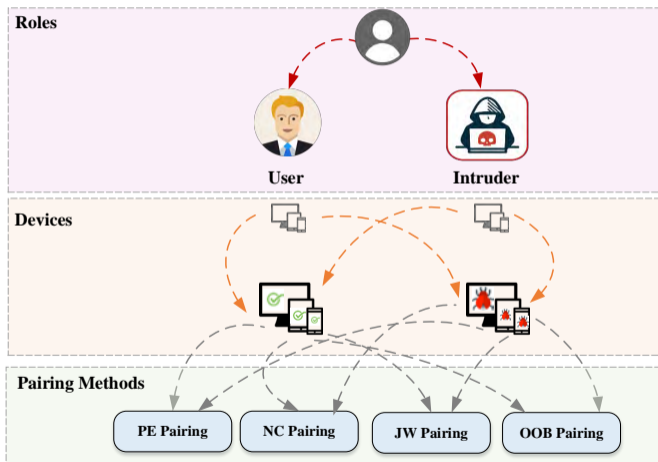
BROADCASTING one-to-many



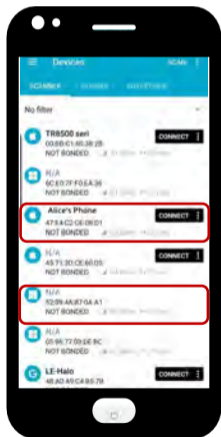
MESH many-to-many



Task 2: Developing a Formal Model for All Pairing Methods (Security)



Task 3: Modelling Linkability of BLE Devices for Privacy



Alice's
phone

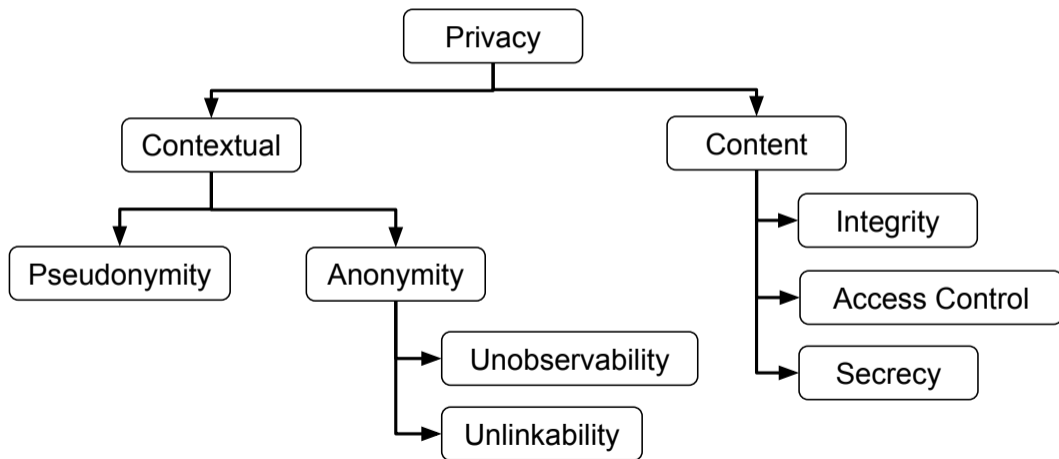
Bob's
phone

T1: 52:09:4A:87:0A:A1



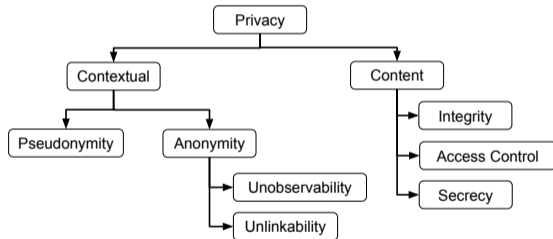
T2: 52:09:4A:87:0A:A1

Task 3: Modelling Linkability of BLE Devices for **Privacy**

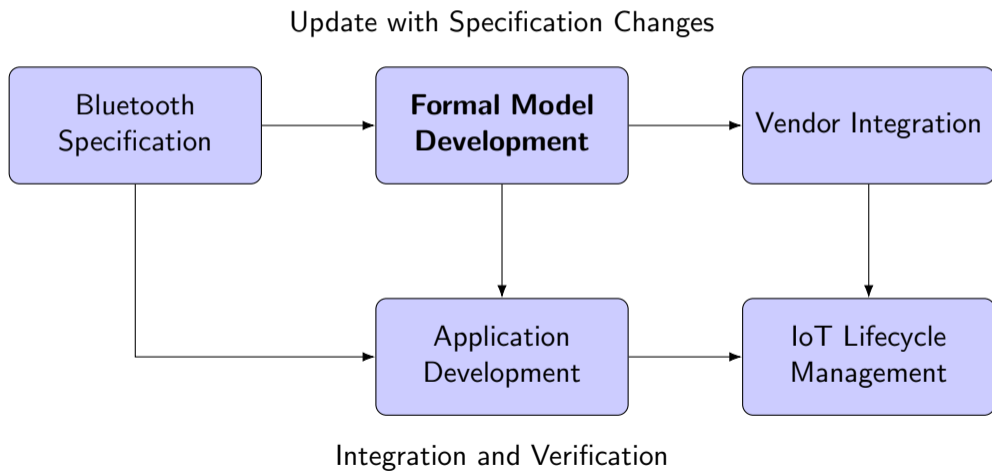


Task 3: Modelling Linkability of BLE Devices for **Privacy**

- ▶ **Unlinkability** implies that an attacker cannot relate multiple observations of user actions.
- ▶ Proposed Solution: Using Observational Equivalence
 - ▶ Finding differences between all possible execution traces of two annotated systems: left and right.



Task 4: Integrating Formal Verification into the Supply-chain



Deliverables

- ① **Formal models of the Bluetooth protocol:** Complete formal models for the Bluetooth Low Energy protocol, covering its various aspects, including device pairing, authentication, and communication.
- ② **Analysis of the discovered vulnerabilities:** A report detailing the identified vulnerabilities in Bluetooth, based on the formal models developed.
- ③ **Open-source implementation:** A prototype implementation of the proposed security enhancements for the Bluetooth protocol, released as an open-source project for the community.
- ④ **Research publications:** Publish findings in peer-reviewed venues to contribute to the global knowledge base on Bluetooth IoT security and privacy.