

Metadata Conscious Anonymous Messaging

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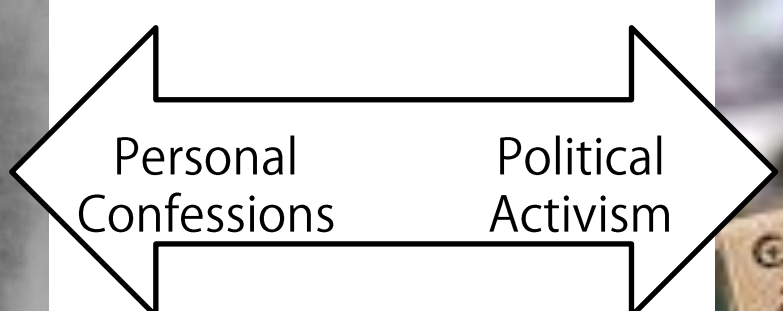
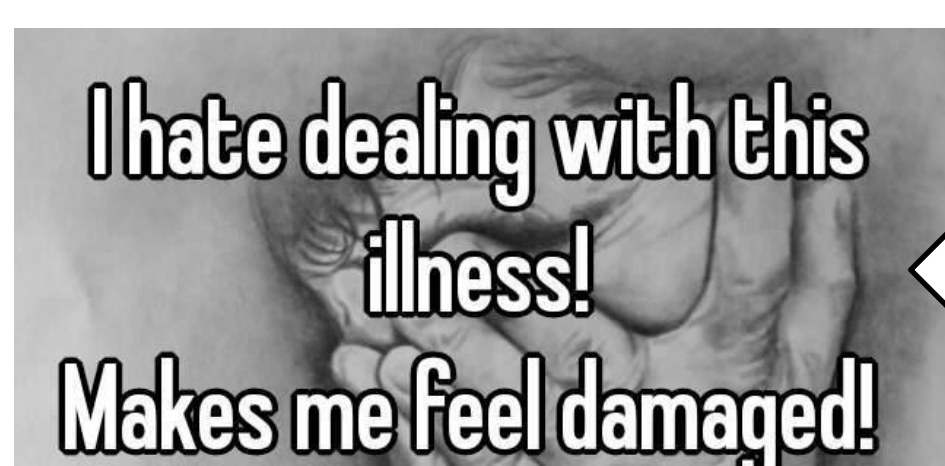
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Anonymity matters



Jason Rezaian's Year of Imprisonment in Iran

Wednesday marks the one-year anniversary of the Washington Post reporter's detention in the Islamic Republic.

Russian Activists and Journalists Attacked at Chechen Border

Saudi Man Gets 10 Years, 2,000 Lashes Over Atheist Tweets

By THE ASSOCIATED PRESS - RIYADH, Saudi Arabia - Feb 27, 2016, 8:26 AM ET

Syria blocks Facebook in Internet crackdown

DAMASCUS | BY KHALED YACOB OWES

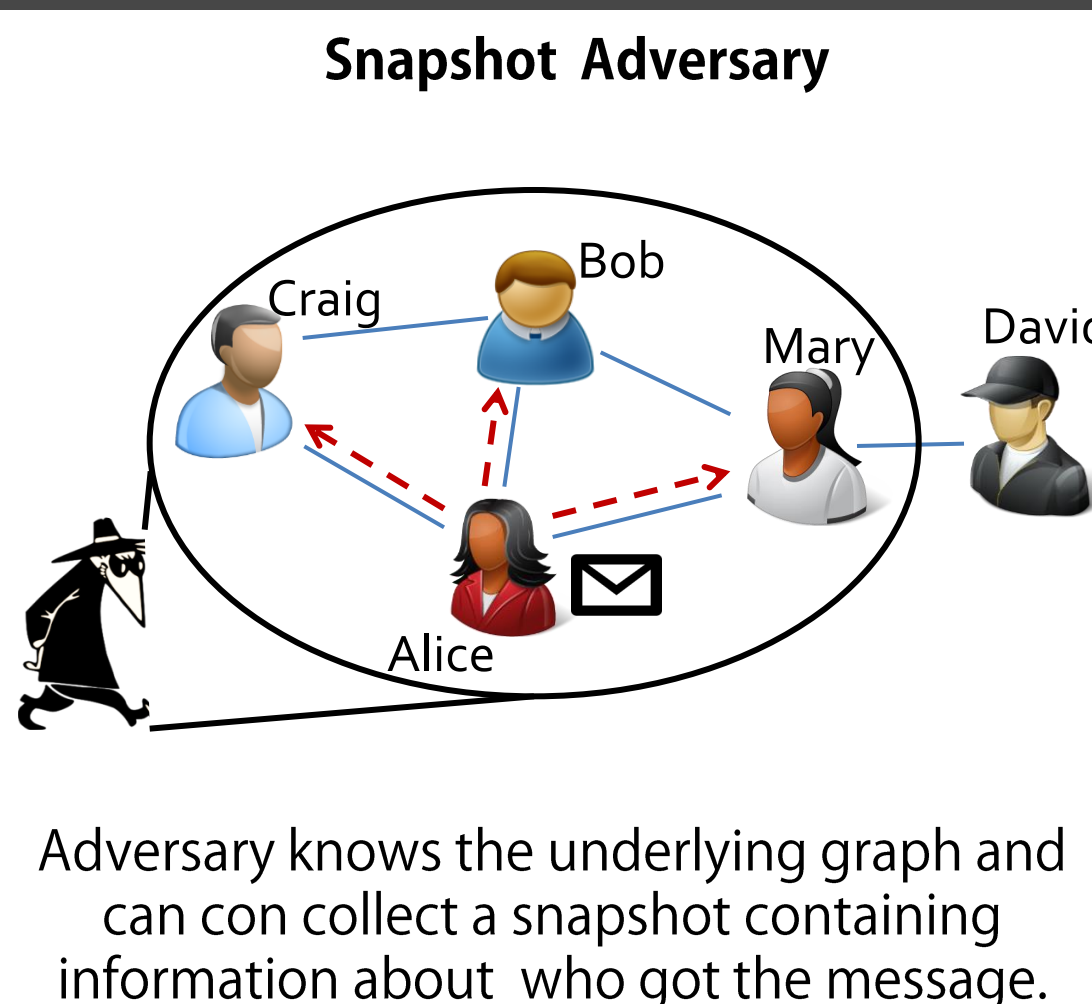
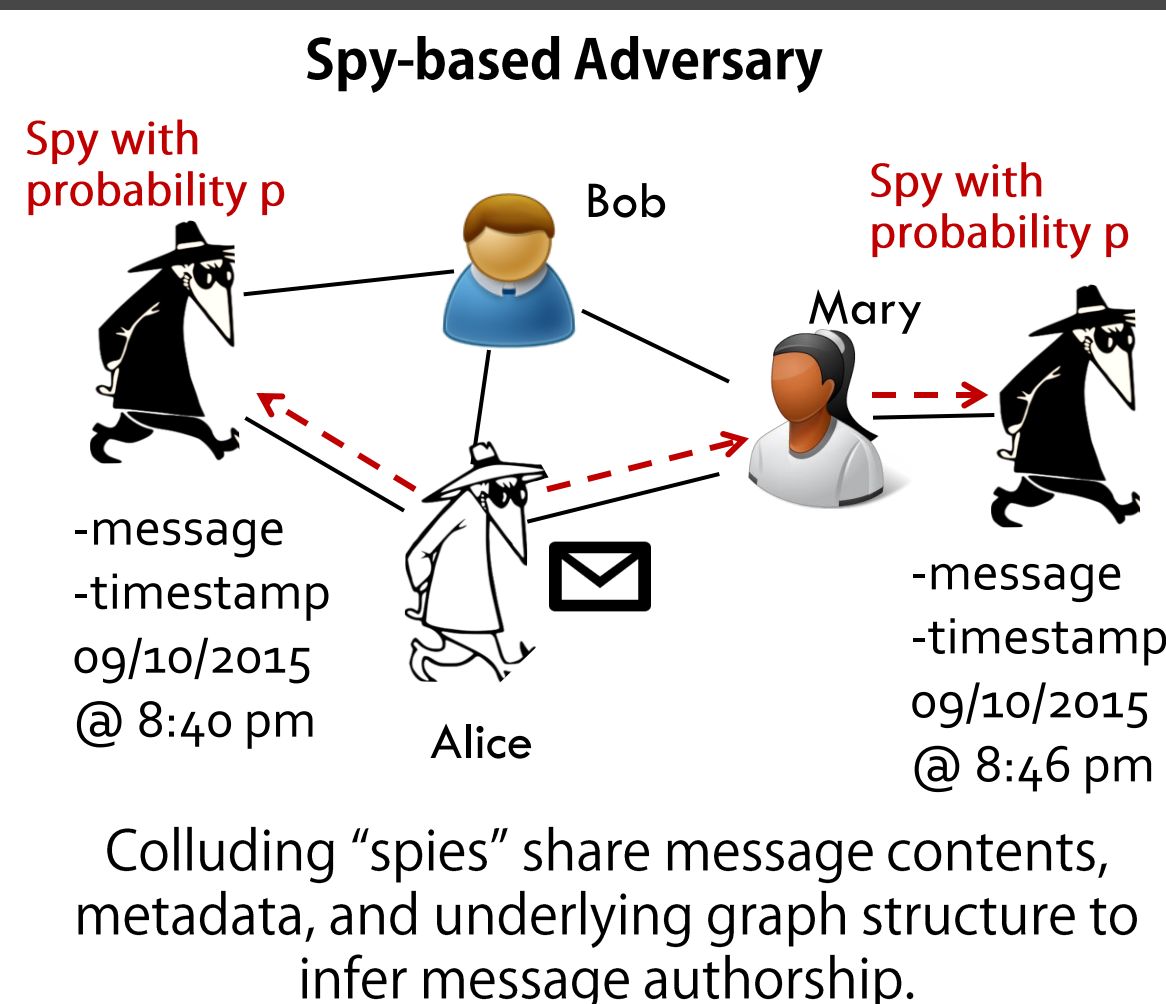
How can we empower people to speak without fear of social or political retribution?

The problem

Design a distributed messaging algorithm that:

- Prevents a powerful adversary from identifying the true message source,
- Spreads content quickly over contact graphs.

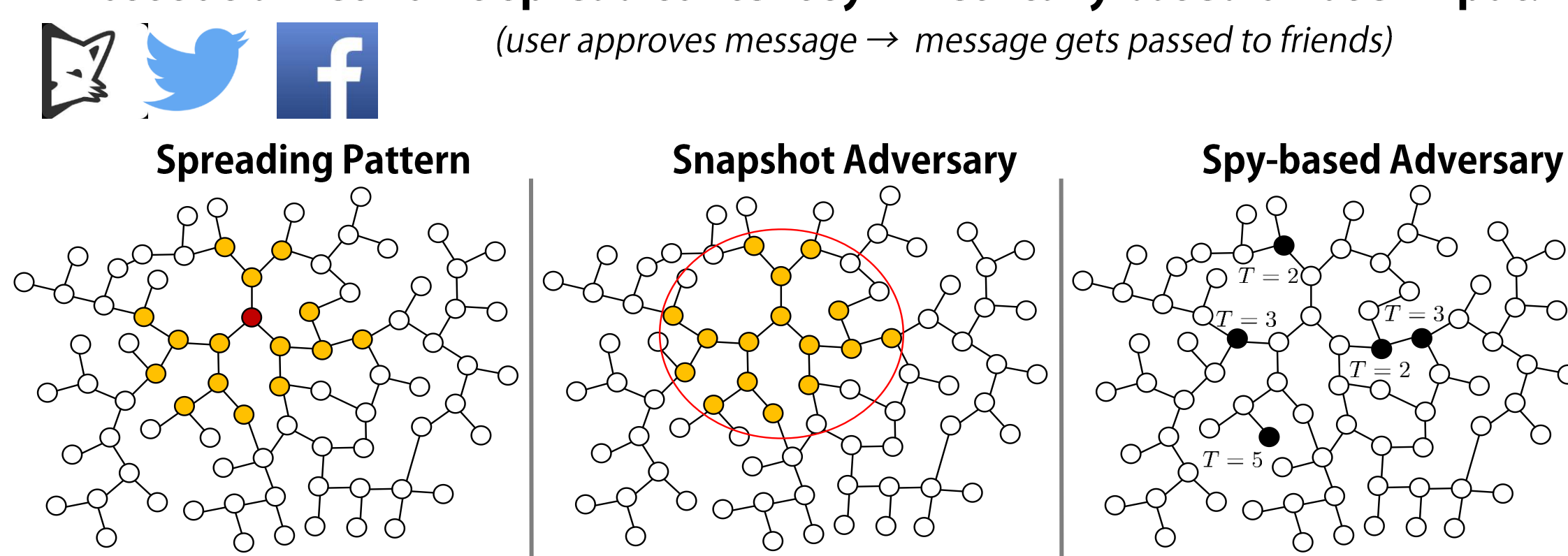
The adversaries



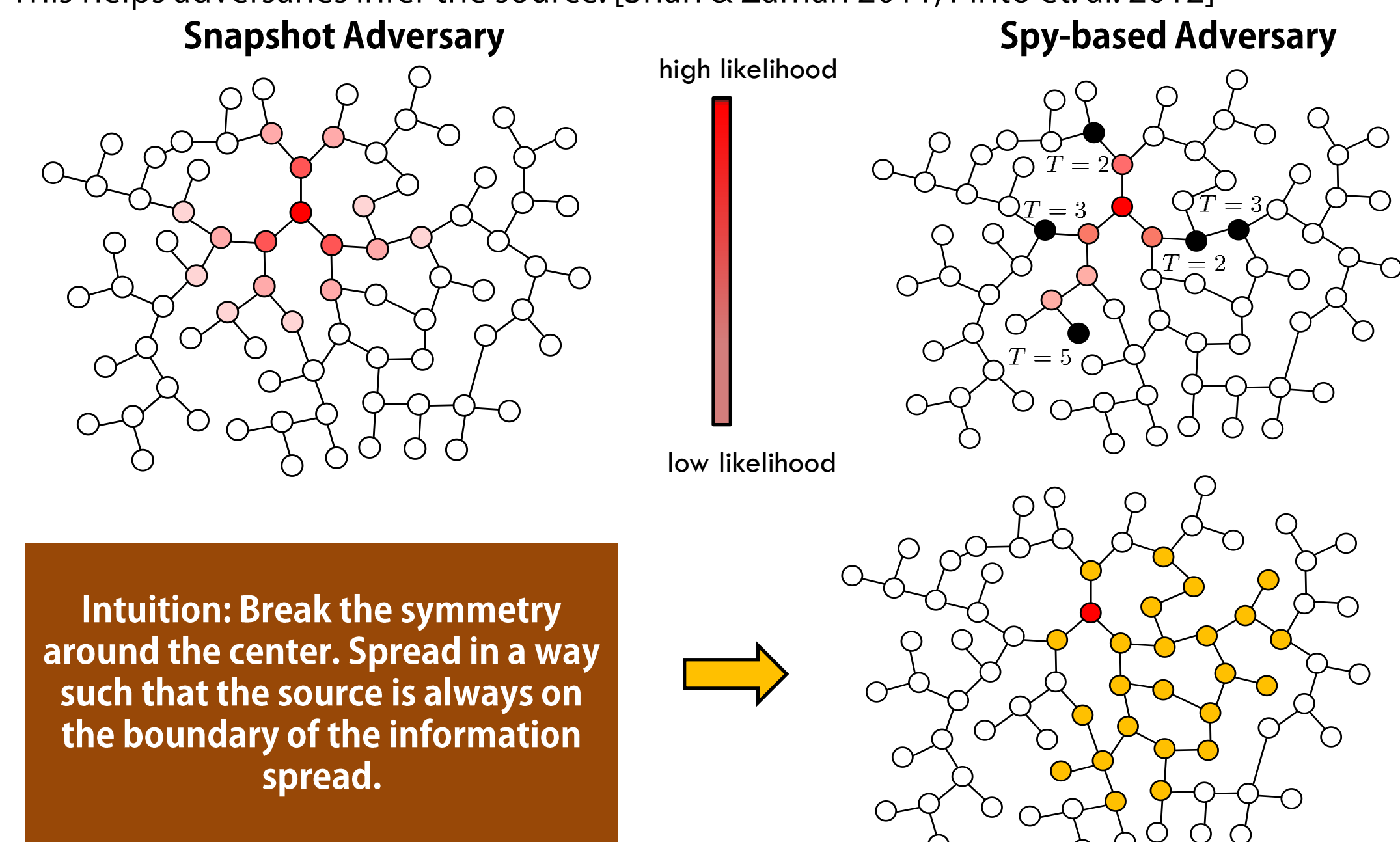
Information flow in social networks

Most social networks spread content symmetrically based on user input.

(user approves message \rightarrow message gets passed to friends)



This spreading model is known as the **diffusion model**. Messages flow in all directions at the same rate. With high probability, diffusion places the **true source in the center** of the graph. This helps adversaries infer the source. [Shah & Zaman 2011, Pinto et. al. 2012]

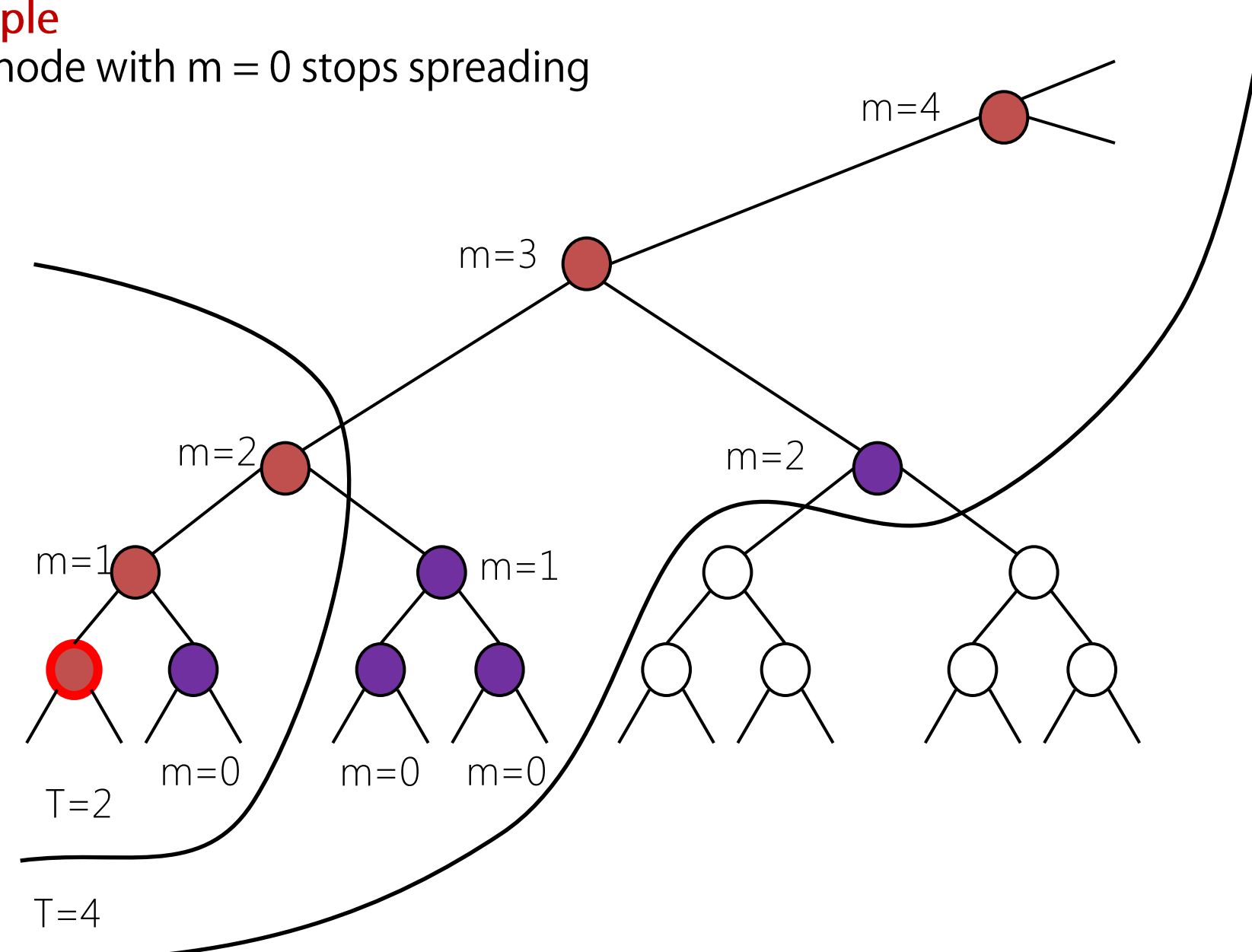


Adaptive diffusion

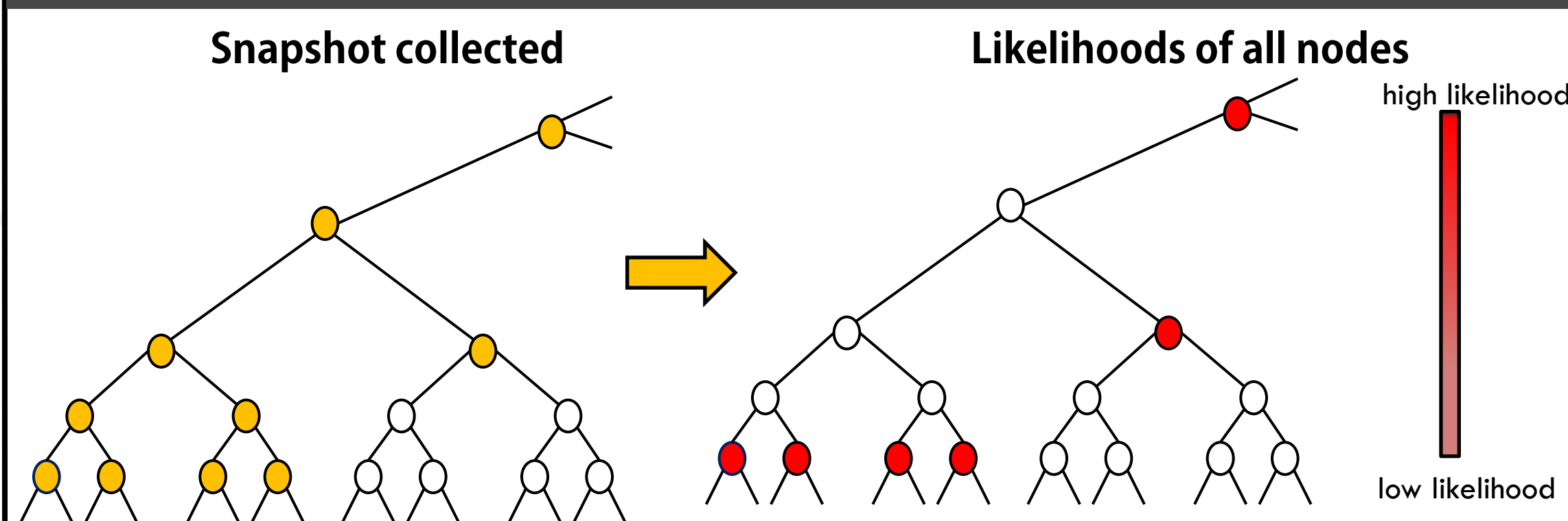
Adaptive diffusion breaks symmetry to provide strong anonymity.
Intuition: carefully adapt the information flow rate and direction

Adaptive diffusion protocol:

- The source node chooses a neighbor at random, and passes the message to it with $m = 1$ & **color = orange**
- An orange node
 - passes the message to a randomly chosen neighbor with an **incremented m & color = orange**,
 - and then to all other neighbors with a **decremented m & color = purple**.
- A purple node with $m > 1$ passes the message to all its neighbors with a **decremented m & color = purple**
- A purple node with $m = 0$ stops spreading

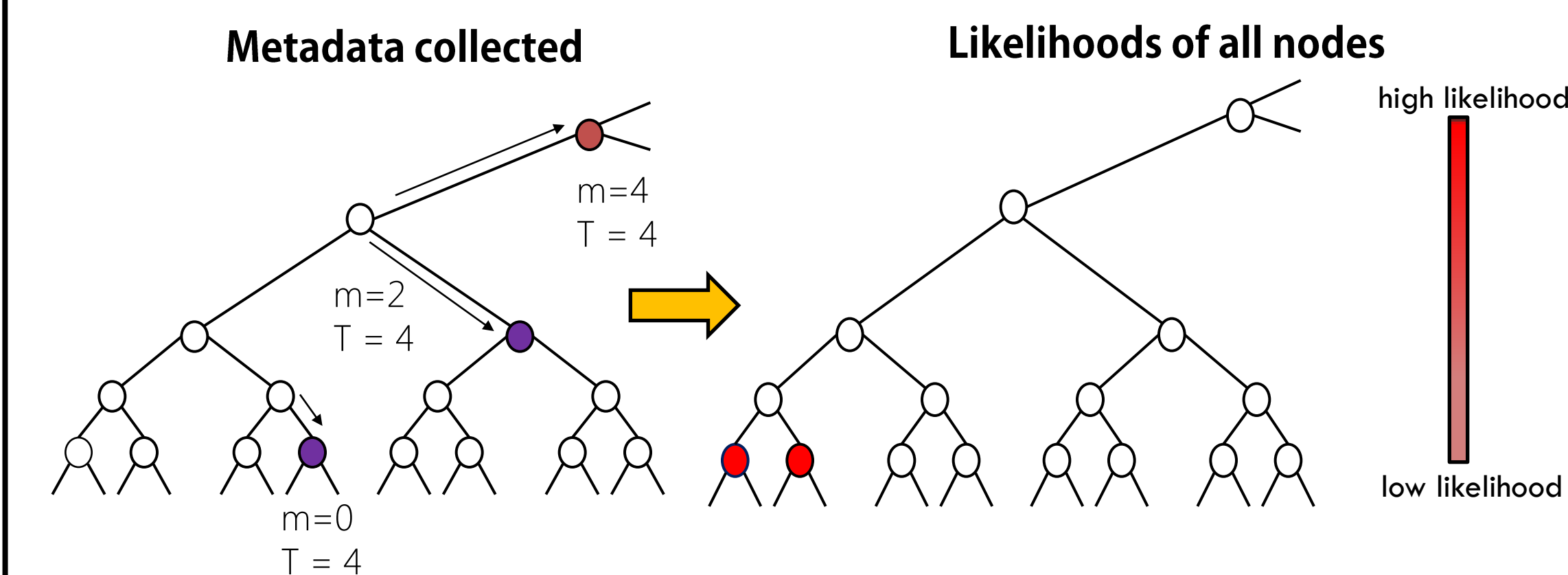


Regular trees: snapshot adversary



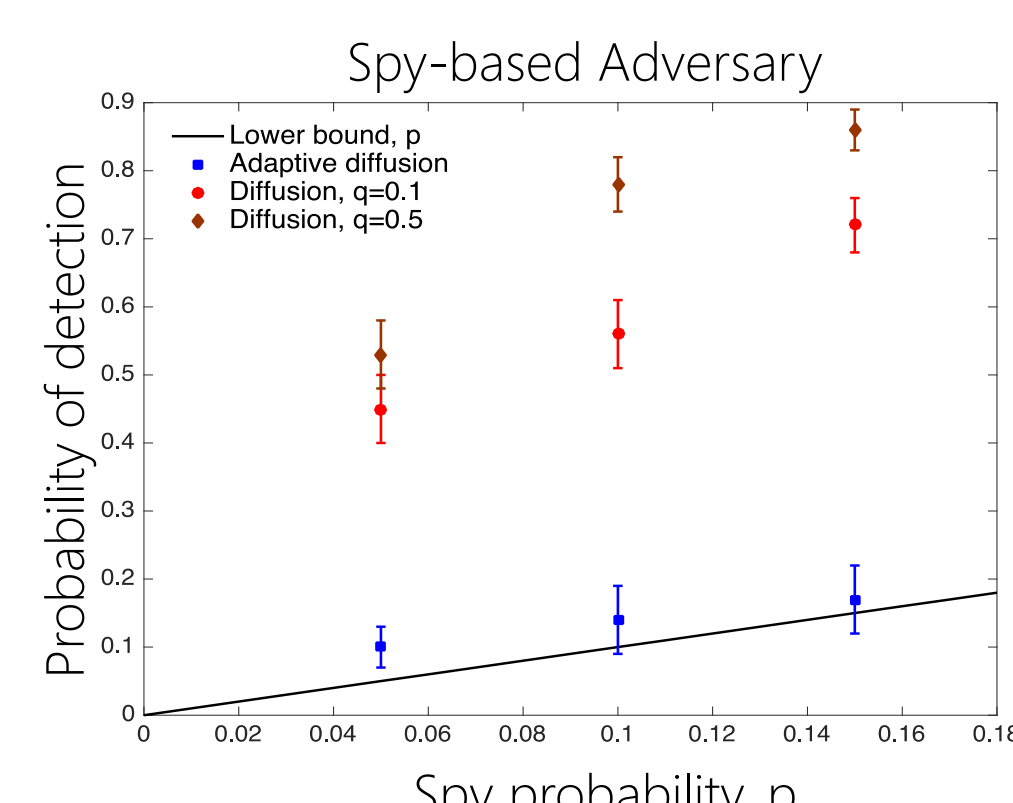
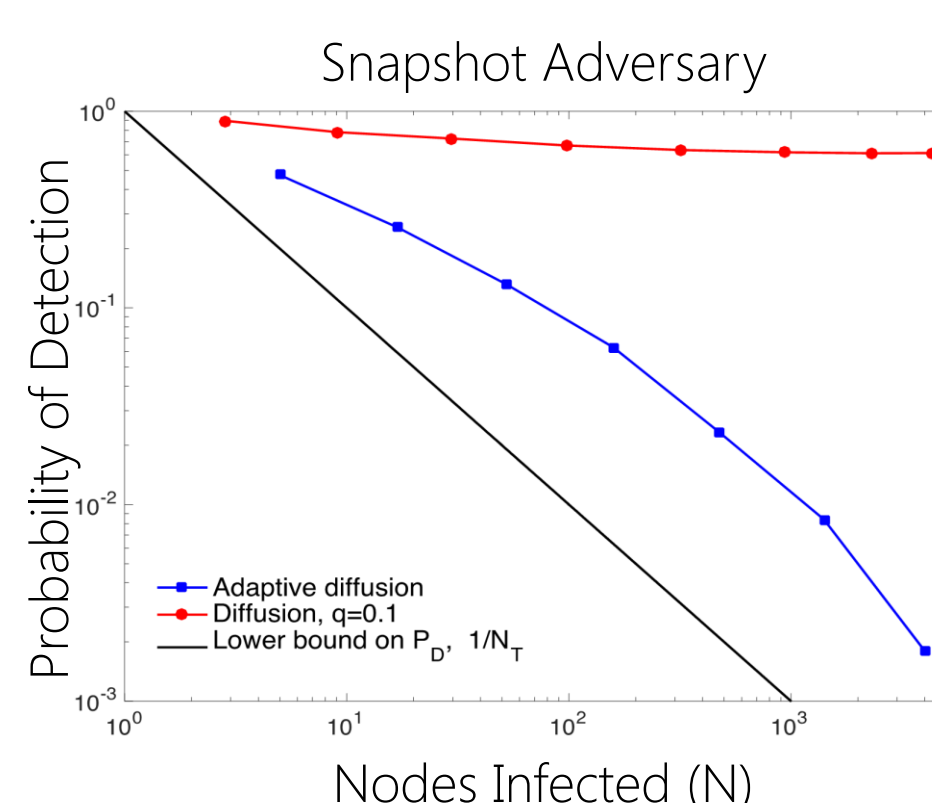
Theorem: On regular trees, adaptive diffusion hides the source **in all the leaf nodes** and spreads the message **exponentially quickly**. Therefore, the probability of rumor source detection is inversely proportional to the number of nodes with the message.

Regular trees: spy-based adversary



Theorem: On regular trees, the probability of detection is always greater than or equal to p . Moreover, the probability of detection under adaptive diffusion is $p + o(p)$. The limit of the probability of detection, as the degree of the tree goes to infinity, is equal to p .

Adaptive diffusion on real graphs



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