

A Quantitative Methodology for Security Monitor Deployment

Uttam Thakore, Gabriel A. Weaver, William H. Sanders University of Illinois at Urbana-Champaign

ITI.ILLINOIS.EDU





Security monitoring today

Administrators must decide what data to collect and how it should be analyzed



Problem

Difficult to determine which monitors are **necessary** to meet intrusion detection requirements

Risks:

- Overprovisioned monitors large volumes of poorly actionable logs
- Underprovisioned monitors insufficient ability to detect or investigate security incidents

We help administrators determine exactly where they stand

- Can expose weaknesses in monitoring



Our contribution

We have developed a **quantitative**, **cost-sensitive** methodology for monitor selection that **meets intrusion detection requirements**

Guiding principles

 Monitors and computing assets can be compromised

- Monitor compromise can affect ability to detect intrusions

• **Redundant monitoring** can mitigate the effect of compromise or unavailability

Outline



O INFORMATIONTRUST

INSTITUTE

Model: Data model

Monitors

Indicators

Sensors that collect information about the system





Primitives representing information provided by monitors about events

Events

0

Intrusions or actions symptomatic of attacks



Case study: E-commerce web service

0



INSTITUTE

Model: Case study data model

0.00



Outline

I



INFORMATIONTRUST

Monitoring metrics

- **Goal of metrics**: *quantify* utility and cost of monitors in supporting intrusion detection
- Three monitor utility metrics:
 - Coverage
 - Redundancy
 - Confidence
- One cost metric:
 - Monitor cost



Metrics: Coverage

Definition: overall fraction of select events that are detectable given a set of monitors



$$\mathbf{Cov}(\{\phi_1, \phi_2, \phi_3\}, \\ \{m_2, m_3\}) = 67\%$$

Metrics: Redundancy

Definition: the number of ways an event can be detected given a set of monitors



Monitors

$$\operatorname{Red}(\phi_1, \{m_1, m_2\}) = 2$$

0

$$\operatorname{Red}(\phi_2, \{m_1, m_2\}) = 1$$

 $\operatorname{Red}(\phi_3, \{m_1, m_2\}) = 0$

13

 $\operatorname{Red}(\phi, M_d) = \sum_{\sigma \in \varsigma(\phi, M_d)} \min_{\iota \in \sigma} \left| \left\{ m \mid m \in M_d, \ \iota \in \alpha(m) \right\} \right|$

Metrics: Confidence

Definition: belief in the ability to detect events accurately, even when monitors are compromised



Monitors

 $\mathbf{Conf}(\phi, M_d) = \max_{\sigma \in \beta(\phi)} \min_{\iota \in \sigma} \gamma_I(\iota, M_d)$

14

0

Metrics: Cost model

- Resource utilization cost
 - CPU utilization
 - Memory utilization
 - Disk storage
 - Network communication
- Amortized purchase price and recurring maintenance cost

Outline

I



INFORMATIONTRUST

Optimal selection methodology

Goal: to be able to use methodology to answer a variety of monitor selection questions

- Minimum set of monitors that can detect a given attack/set of attacks (assuming no compromise)?
- Under cost constraints, what set of monitors will maximize ability to detect a high-priority attack?

Capturing intrusion detection requirements

Represent detection requirements as **weights on metric values** and **minimum metric value constraints**



Requirements:

- Must detect exfiltration
- Exfiltration, then SQL injection, are high priority (decreasing priority)
- Best effort for all others

$$\min_{\operatorname{Red}_{\phi_2}} = 1$$
 $\mathbf{w}_{\operatorname{Red}_{\phi_2}} = 2$ $\mathbf{w}_{\operatorname{Cov}} = 1$

$$\mathbf{w}_{\operatorname{Red}_{\phi_3}} = 1$$

Optimal selection methodology: Constrained-cost monitor selection

 $\underset{M_d}{\operatorname{arg\,max}}$

$$\mathbf{w}_{\text{Cov}} \mathbf{Cov} (\Phi, M_d) + \sum_{\phi \in \Phi} \mathbf{w}_{\text{Red}_{\phi}} \mathbf{Red} (\phi, M_d) + \mathbf{w}_{\text{Conf}_{\phi}} \mathbf{Conf} (\phi, M_d)$$

Objective function: monitoring utility, defined as weighted sum of metric values

 Parameterized by user-specified weight parameters

- Cost function to minimize
- User-specified minimum detection metric requirements

0-1 integer nonlinear programming problem, with monitors as input variables

$$\begin{split} \mathbf{Cost}\left(\boldsymbol{M}_{d}\right) &\leq \mathbf{maxCost} \\ \text{s.t.} \quad \mathbf{Cov}\left(\boldsymbol{\Phi},\boldsymbol{M}_{d}\right) &\geq \mathbf{min}_{\text{Cov}} \\ \mathbf{Red}\left(\boldsymbol{\phi},\boldsymbol{M}_{d}\right) &\geq \mathbf{min}_{\text{Red}_{\phi}}, \quad \forall \boldsymbol{\phi} \in \boldsymbol{\Phi} \\ \mathbf{Conf}\left(\boldsymbol{\phi},\boldsymbol{M}_{d}\right) &\geq \mathbf{min}_{\text{Conf}_{\phi}}, \quad \forall \boldsymbol{\phi} \in \boldsymbol{\Phi} \\ \boldsymbol{M}_{d} &\in \left\{0,1\right\}^{|\mathcal{M}|} \end{split}$$

Solving for optimal monitor selection

- Branch-and-bound algorithm
 - Searches over space of possible selections, pruning suboptimal sets of monitor selections
- Greedy heuristic algorithm
 - Maximizes effective utility increase by incrementally adding monitors until constraints are met



EVALUATION

Experiment Setup

- Parameters: number of *monitors* and *events*
- Randomly generated 100 models for each set of parameters
- Created 4 sets of intrusion detection requirements optimal deployment programs

Goal: Observe scalability and accuracy of greedy solution

Evaluation: Greedy algorithm

Runtime complexity: $O(|I|(|M|^3 + |B||M|^2))$

– Polynomial in the number of monitors (|M|)



Evaluation: Greedy algorithm

Runtime complexity: $O(|I|(|M|^3 + |B||M|^2))$

– Linear in the number of minimal indicator sets (|B|) and indicators (|I|)



Conclusions



responded to by

- We help administrators make model-driven monitor placement decisions
- Administrators can more easily evaluate deployments
- Our methodology is expressive and scalable

Future work:

Preemptive monitoring