Entropy-minimizing Mechanism for Differential Privacy of Discrete-time Linear Feedback Systems

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General Question

Trade-off between "privacy" and "accuracy": a common strategy to protect some data private is to randomize it, but this undermines the accuracy of the data. **Example**¹:



Figure: Block Diagram for $\epsilon\text{-Differentially}$ Private Discrete-time Linear Feedback System

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¹Huang et al., HiCoNS 14.

Preliminaries

In this work, we use the concept of ϵ -differential privacy as a measure of privacy. It originates from the study of privacy-preserving queries of datasets ² and later extends to dynamic systems.

Definition

The mechanism \mathcal{M} is ϵ -differentially private if the inequality

$$\mathbb{P}\left[\mathcal{M}(\boldsymbol{x}_1) \subseteq O\right] \le \exp\left(\epsilon \|\boldsymbol{x}_1 - \boldsymbol{x}_2\|_1\right) \mathbb{P}\left[\mathcal{M}(\boldsymbol{x}_2) \subseteq O\right] \qquad (1)$$

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holds for any inputs x_1, x_2 and a set of possible outputs O, where $\|\mathbf{x}\|_1 = \sum_{i=1}^n |x_i|$.

²C. Dwork, 2006.

Preliminaries



Accuracy is measured by Shannon entropy. For a random variable X on \mathbb{R}^n with probability distribution function $f(\mathbf{x})$,

$$\mathbf{H}(X) = -\int_{\mathbb{R}^n} f(\mathbf{x}) \ln(\mathbf{x}) d\mathbf{x}$$
(2)

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One-shot Query



Figure: Block Diagram for a ϵ -Differentially Private Mechanism

Conditions:

- $\blacktriangleright X, Y \in (\mathbb{R}^n, \|\cdot\|_1)$
- the joint p.d.f. p(x, y) is absolute continuous;
- the noise N(X) is zero-mean;
- the accuracy is measured by $\mathbf{H}(\mathcal{M}) = \sup_{X} \mathbf{H}(Y)$.

Theorem

For an ϵ -differentially private mechanism \mathcal{M} with input set $(\mathbb{R}^n, \|\cdot\|_1)$, we have $\mathbf{H}(\mathcal{M}) \ge n - n \ln(\epsilon/2)$ and the minimum is achieved by $p(\mathbf{x}, \mathbf{y}) = (\frac{\epsilon}{2})^n \exp(-\epsilon \|\mathbf{y} - \mathbf{x}\|_1) = \prod_{i=1}^n (\frac{\epsilon}{2} e^{-\epsilon |y_i - x_i|})$.



Trade-off: Privacy $\uparrow \implies \epsilon \downarrow \implies \mathbf{H}(\mathcal{M}) \uparrow \implies \mathsf{Accuracy} \downarrow$

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Conditions:

- ► $X(t), Y(t), Z(t), U(t), V(t) \in (\mathbb{R}^n, \|\cdot\|_1)$
- > zero input: U(t) = 0
- unit gain feedback: V(t) = Y(t) = Z(t)
- dynamics: X(t+1) = AX(t) + BV(t).



The adversary A only has access to the randomized outputs $\{Z(i) \mid i \in [t]\}$. Since

$$X(t) = A^{t}X(0) + \sum_{i=0}^{t-1} A^{t-i-1}BZ(i),$$
(3)

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protecting the ϵ -differential privacy of the initial system state is equivalent to protecting the ϵ -differential privacy of the whole trajectory.

The adversary A estimates the initial system state from the past history of randomized outputs $\{Z(i) \mid i \in [t]\}$ by

$$\tilde{X}(t) = \mathbb{E}\left[X(0) \mid Z(0), Z(1), \dots, Z(t)\right],\tag{4}$$

The accuracy of the output of the mechanism \mathcal{M} at time $t \in \mathbb{N}$ is measured by

$$\mathbf{H}(\mathcal{M},t) = \mathbf{H}\left(\tilde{X}(t)\right).$$
(5)

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The mechanism \mathcal{L} is ϵ -differentially private up to time $t \in \mathbb{N}$, if for any pair of initial states $\mathbf{x}_1, \mathbf{x}_2 \in \mathbb{R}^n$, and output history $\{\mathbf{z}(i) \mid i \in [t]\}$,

$$\mathbb{P}[Z(1) = \mathbf{z}(1), \dots, Z(t) = \mathbf{z}(t) \mid X(0) = \mathbf{x}_1] \\
\mathbb{P}[Z(1) = \mathbf{z}(1), \dots, Z(t) = \mathbf{z}(t) \mid X(0) = \mathbf{x}_2] \\
\leq \exp(\epsilon \|\mathbf{x}_1 - \mathbf{x}_2\|).$$
(6)

By Bayes formula, (6) is equivalent to

$$\tilde{h}_t(\boldsymbol{x}_1) \le \exp\left(\epsilon \|\boldsymbol{x}_1 - \boldsymbol{x}_2\|\right) \tilde{h}_t(\boldsymbol{x}_2). \tag{7}$$

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where \tilde{h}_t is the probability density function of $\tilde{X}(t)$.

Theorem

If a mechanism is ϵ -differentially private up to time $t \ge 0$, then

$$\mathbf{H}(\mathcal{L},i) \ge n - n \ln(\frac{\epsilon}{2}) \tag{8}$$

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for $i \in 1, ..., t$. The equality holds when $N(0) \sim Lap(1/\epsilon)$, and for $t \geq 1$, N(t) = AN(t-1). In this case

$$\mathbf{H}(\mathcal{L},1) = \mathbf{H}(\mathcal{L},2) = \ldots = \mathbf{H}(\mathcal{L},t) = n - n \ln(\frac{\epsilon}{2}).$$
(9)

Proof of Theorem

Assume $X, Y \in \mathbb{R}$.

Problem

Minimize: $H(\mathcal{M})$ subject to: $\mathbb{P}[\mathcal{M}(x_1) \subseteq O] \leq \exp(\epsilon ||x_1 - x_2||_1) \mathbb{P}[\mathcal{M}(x_2) \subseteq O]$

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Claim 1: for fixed x, p(x, y - x) is even.

$$\mathbf{H}_{1}^{+}(\mathcal{M}) = \sup_{x \in \mathbb{R}} \int_{[x,\infty)} -p(x,y) \ln p(x,y) \mathrm{d}y, \quad (10)$$
$$\mathbf{H}_{1}^{-}(\mathcal{M}) = \sup_{x \in \mathbb{R}} \int_{(-\infty,x]} -p(x,y) \ln p(x,y) \mathrm{d}y. \quad (11)$$

$$q(x,y) = \begin{cases} p(x,y) & \text{if } y > x, \mathbf{H}_{1}^{+}(\mathcal{M}) \leq \mathbf{H}_{1}^{-}(\mathcal{M}) \\ \text{or } y < x, \mathbf{H}_{1}^{+}(\mathcal{M}) > \mathbf{H}_{1}^{-}(\mathcal{M}), \\ p(x,2x-y) & \text{if } y > x, \mathbf{H}_{1}^{+}(\mathcal{M}) > \mathbf{H}_{1}^{-}(\mathcal{M}) \\ \text{or } y < x, \mathbf{H}_{1}^{+}(\mathcal{M}) \leq \mathbf{H}_{1}^{-}(\mathcal{M}). \end{cases}$$
(12)

 $\mathbf{H}(\mathcal{N}) = 2\min\{\mathbf{H}_{1}^{+}(\mathcal{M}), \mathbf{H}_{1}^{-}(\mathcal{M})\} \leq \mathbf{H}_{1}^{+}(\mathcal{M}) + \mathbf{H}_{1}^{-}(\mathcal{M}) = \mathbf{H}(\mathcal{M}),$ (13)

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Claim 2: for any x, p(x, y) = p(2a - x, 2a - y).

$$\mathbf{H}^{+}(\mathcal{M}) = \sup_{x > a} \int_{\mathbb{R}} -p(x, y) \ln p(x, y) \mathrm{d}y, \qquad (14)$$

$$\mathbf{H}^{-}(\mathcal{M}) = \sup_{x \leq a} \int_{\mathbb{R}} -p(x, y) \ln p(x, y) \mathrm{d}y.$$
(15)

If $\mathbf{H}^{+}\left(\mathcal{M}
ight)\leq\mathbf{H}^{-}\left(\mathcal{M}
ight)$, then define

$$q(x,y) = \begin{cases} p(x,y), & x > a, \\ p(2a - x, 2a - y), & x \le a, \end{cases}$$
(16)

otherwise, define

$$q(x,y) = \begin{cases} p(2a - x, 2a - y), & x > a, \\ p(x,y), & x \le a. \end{cases}$$
(17)

 $\mathbf{H}(\mathcal{N}) = \min\{\mathbf{H}^{+}(\mathcal{M}), \mathbf{H}^{-}(\mathcal{M})\} \leq \max\{\mathbf{H}^{+}(\mathcal{M}), \mathbf{H}^{-}(\mathcal{M})\} = \mathbf{H}(\mathcal{M})$ (18)
(18)

Claim 3: p(x, y) = f(y - x). Let q(x, y) = p(x, y - x). By Claim 2, q(x, y) = q(2a - x, -y). By Claim 1, q(2a - x, -y) = q(2a - x, y). Now the problem becomes,

Problem

Minimize:
$$H(f) = -\int_{[0,\infty)} f(x) \ln f(x) dx$$
,
subject to: $f(x)$ is absolutely continuous,

$$f(x) \ge 0,$$

$$|f'(x)| \le \epsilon f(x) \text{ a.e.},$$

$$\int_{[0,\infty)} f(x) dx = \frac{1}{2}.$$

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Claim 4: f(x) is decreasing.

Let x^* be a local minimum on (0, 1). Then there exists $x^* \in [a, b]$ such that f(a) = f(b) > f(x) for $x \in (a, b)$. Let $d = \frac{1}{f(a)} \int_a^b f(x) dx$ and

$$h(x) = \begin{cases} f(x), & x \in [0, a], \\ f(b), & x \in [a, a+d], \\ f(x+b-a-d), & x \in [a+d, \infty]. \end{cases}$$
(19)

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Then H(h) < H(f).

Let
$$F(x) = \int_{x}^{\infty} f(y) dy$$
.

$$F(x) \ge \int_{x}^{\infty} \frac{|f'(x)|}{\epsilon} dy \ge \frac{1}{\epsilon} |\int_{x}^{\infty} f'(x) dy|$$

$$= \frac{1}{\epsilon} |f(\infty) - f(x)| = \frac{f(x)}{\epsilon}$$
(20)

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In particular, $f(0) \ge \epsilon F(0) = \frac{\epsilon}{2}$.

$$H(f) = -\int_{0}^{\infty} f(x) \ln f(x) dx$$

= $-\int_{0}^{\infty} f(x) \left(\ln f(0) + \int_{0}^{x} \frac{f'(y)}{f(y)} dy \right) dx$
= $-\frac{1}{2} \ln f(0) - \int_{0}^{\infty} \frac{f'(y)}{f(y)} \left(\int_{x}^{\infty} f(x) dx \right) dy$ (21)
= $-\frac{1}{2} \ln f(0) - \int_{0}^{\infty} \frac{f'(y)F(y)}{f(y)} dy$
 $\geq -\frac{1}{2} \ln f(0) - \int_{0}^{\infty} \frac{f'(y)}{\epsilon} dy$
= $\frac{f(0)}{\epsilon} - \frac{1}{2} \ln f(0) \geq \frac{1}{2} - \ln(\frac{\epsilon}{2}),$

The minimum is achieved by

$$f(x) = \frac{\epsilon}{2} \exp(-\epsilon x). \tag{22}$$

Thanks!

