



A Retrospective on Constructive Verification



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Contents

- Well...it was was 21(ish) years ago today...
- Retrospective vs Constructive verification
- Getting to 4th base...
- A future?

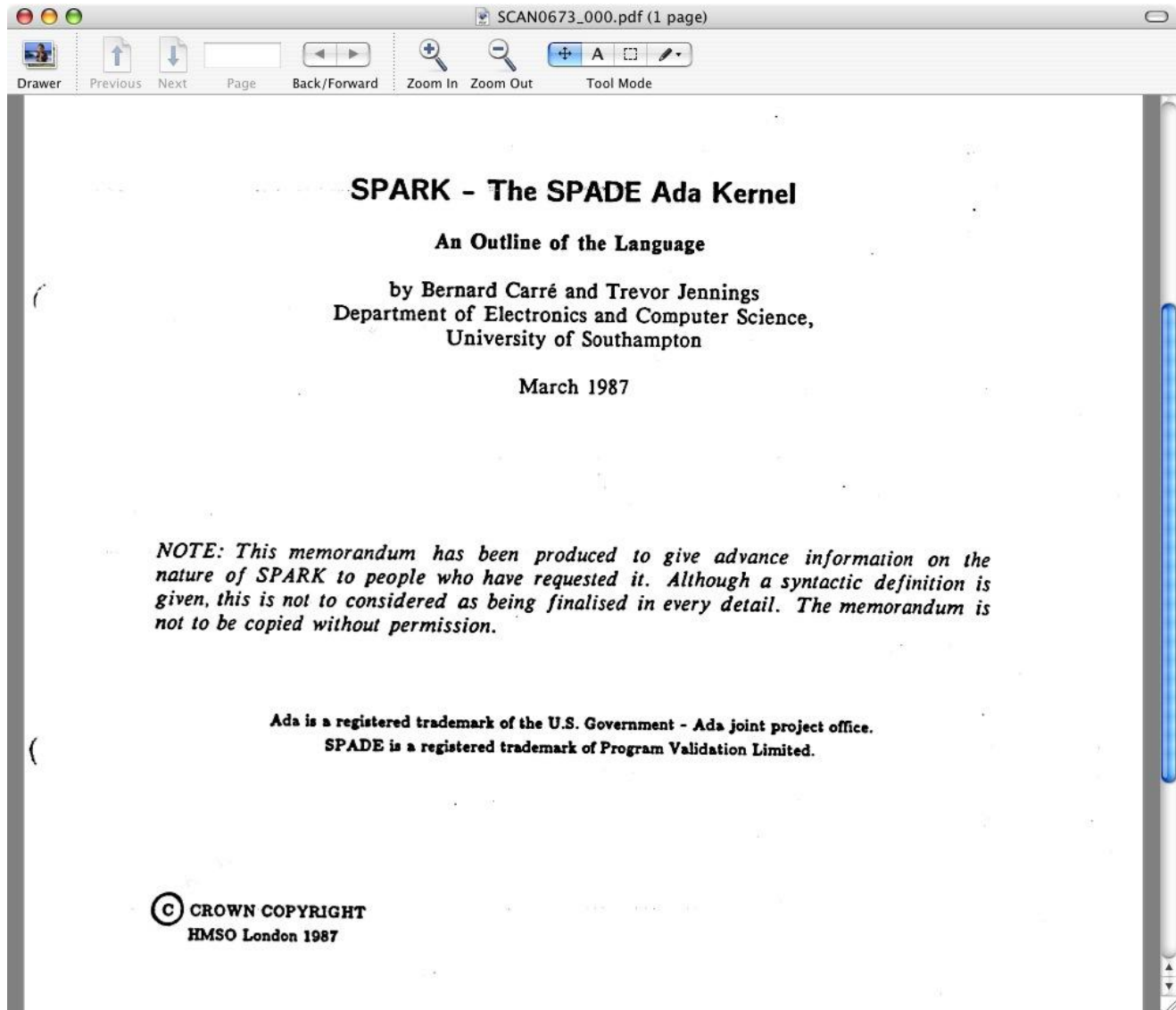


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Well it was 21(ish) years ago today





Well it was 21(ish) years ago today

SPARK - The SPADE Ada Kernel

An Outline of the Language

**by Bernard Carré and Trevor Jennings
Department of Electronics and Computer Science,
University of Southampton**

March 1987



So what...

- PVL/Praxis/SPARK team have been
 - Designing programming languages...
 - Building static analysis tools...
 - Actually trying to use them on real projects...

 - ...for what seems like a long time.
- Here are a few reflections on what we've learnt and what's going on now...



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Early days...

- UK Military Aerospace
 - Software begins to appear in military aircraft in about 1985-1990
 - No tools exist...what can you do?
 - This led to *retrospective* style of analysis



Retrospective analysis

- Typical process
 - Procure/take delivery/pay for box containing software
 - Peer at it for a long time
 - Report bugs that you find. Hope they might be fixed.
 - Try to decide to fly aeroplane or not...
 - (In the mean time, some tools get developed – SPADE, MALPAS etc)



Retrospective analysis

- Analysis typically carried out by buyer/evaluator *after* development and “test”.
- Observations:
 - Perception of limited utility: analysis is hard, slow, and subject to human frailty.
 - Little motivation for the developer to change their ways or do better



Retrospective analysis

- Key observation
 - Utility of retrospective analysis critically depends not only on quality/power of tools, but also on the quality of the software under analysis.
 - Poorly designed programs defy analysis by any method, tool or person.
 - Example: Chinook HC2 FADEC
 - Even programs which “seem to work” and “pass testing” defy analysis.



Retrospective analysis

- Is this still true?
 - Huge increase in tool power + 15 iterations of Moore's Law.
- But...
 - Massive increase in program size and complexity
 - Programming languages didn't help...they got bigger, more ambiguous, more dynamic...
- Who is winning this race?



Constructive Analysis

- The big idea:
 - Place tools in hands of developers, to be used all the time...
 - Use *discipline* to manage utility
 - Deliver system with static analysis evidence
 - Regulator and/or customer can reproduce evidence if they want.



Constructive Analysis

- Adoption is hard – requires major change of lifestyle for most developers.
- We encounter enormous resistance to the adoption of discipline.
 - Nobody likes being told what to do...



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Getting to 4th base...

- A “playing field” for static analysis tools
 - 1st base: basic dumb mistakes – subset/coding standard etc.
 - 2nd base: absence of undefined behaviour (e.g uninitialized variables)
 - 3rd base: type safety
 - 4th base: partial correctness, safety and security properties, application and domain specific properties
 - 5th base+: stuff we haven’t even thought of yet...



Getting to 4th base...

- Note: at 4th base and above, desired properties are application and domain specific.
 - There is no “list of vulnerabilities” that can be enumerated or can be “built in” to a tool.
 - Overly generic description (e.g. “SQL Injection”) leads to hopeless false-positive rate from tools.
 - Many languages allow for user-defined properties, via assertions/contracts (e.g. SPARK, Eiffel) or via user-written “Rules” or “Checkers” (e.g. Coverity)



Getting to 4th base...

- SPARK gets to 4th base (just...)
- How?
 - Careful (some would say Draconian...) subset and contractualization of language.
 - Favour soundness above all other design goals.
 - Build soundness – base N+1 depends on base N analyses being OK first.



A worrying conversation

- Customer: “What list of bugs does your tools find?”
- Rod: “There’s no such list – it’s a general-purpose verification framework”
- Customer: “What list of bugs does your tools find?”
- Rod: “Anything that you can express as a predicate in first-order logic”
- Customer: “Eh?” (and leaves...)



A worrying conversation

- Where tools and languages support verification of user-defined properties:
- Perhaps we might ask:
 - “What properties can be expressed? What properties can’t?”
 - “What is the soundness, completeness, and efficiency of the checking algorithm?”
- Many tool vendors don’t seem to be very forthcoming with this information.



Does Soundness Matter?

- “Soundness doesn’t matter”
 - Who says?
 - Well...err...All tool vendors whose tools are unsound.

- Or does it...?



Does Soundness Matter?

- In retrospective analysis mode, it appears not
 - finding 90% of bugs is better than none!
- But...if we are to move to constructive evidence-based assurance, soundness will matter
 - Would you present evidence to an evaluator if you *know* the tool that generated it can be unsound?
 - As an evaluator, would you accept such evidence?



Does Soundness Matter?

- A warning...
- Soundness is a one-way trip...
- Once achieved, customers will get used to it very rapidly, and come to depend on it.
 - You'll never go back...



Intermission...

- Enough moaning...
- Here's comes some code...



An example “4th Base” verification in SPARK.

- SQL Injection
 - Actually, just a special-case of input data validity.

 - It’s both easy, and very hard...



SQL Injection

- Imagine a simple SPARK package that is used to query a database:

```
package DB
--# own State;
--# initializes State;
is
    procedure Query (SQL_String : in      String;
                    Result      : out   String);
--# global in State;
--# derives Result from State, SQL_String;
end DB;
```



SQL Injection

- Dumb implementation of user-generated query:

```
-- get input from user, whatever it is..  
Read_Input (User_String);
```

```
-- construct SQL query string from user input  
Form_Query (User_String, SQL_String);
```

```
-- Chuck the resulting query at the database  
DB.Query (SQL_String, Result);
```

- This implementation is *weak* in that there is no checking that the user-provided string is not malicious, mal-formed, or just wrong.



SQL Injection

- A better SPARK Database Interface:

```
package DB
--# own State;
--# initializes State;
is
    function Valid_Query
        (SQL_String : in String) return Boolean;
--# global in State;

    procedure Query (SQL_String : in      String;
                    Result      : out String);
--# global in State;
--# derives Result from State, SQL_String;
--# pre Valid_Query (SQL_String, State);
end DB;
```



SQL Injection

- Now what happens?

```
-- get input from user, whatever it is...  
Read_Input (User_String);
```

```
-- construct SQL query string from user input  
Form_Query (User_String, SQL_String);
```

```
-- Chuck the resulting query at the database  
DB.Query (SQL_String, Result);
```

- You get an unprovable precondition VC for the call to DB.Query



SQL Injection

- The unprovable VC “reminds” you to bother to check, so I re-write the code:

```
-- get input from user, whatever it is...
Read_Input (User_String);

-- construct SQL query string from user input
Form_Query (User_String, SQL_String);

-- Check validity of generated query
if DB.Valid_Query (SQL_String) then
    -- Chuck the resulting query at the database
    DB.Query (SQL_String, Result);
else
    Error_Handler;
end if;
```



SQL Injection

- The offending precondition VC is now provable.
- Easy huh?
- Well...not quite...there's still no free lunch...



SQL Injection – The Catch...

- You have to write the bodies of DB.Valid_Query and Error_Handler
- What defines a “Valid” query anyway?
 - Look in your specification or security policy
 - You *have* got a specification, right?
- You end up *having* to specify error-handling behaviour as well...



Static analysis for engineer “behaviour modification”

- The upshot of all this:
 - A disciplined/formal/design-by-contract implementation style forces robustness.
 - This leads you to resolve issues in security policy, requirements, and specification.
 - The *behaviour* of engineers (eventually) changes to deal with these issue “up-front” rather than post-hoc.



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A future?

- Who will win? Constructive or Retrospective tools?
 - Hopefully...*both*...
- All systems have many components – some new, some highly critical, some re-used, some COTS, some firmware, written in multiple languages.
 - There must be room for both styles of analysis.



A future?

- Why not use *architecture* to separate the really critical stuff from the rest?
 - Use sound constructive techniques where soundness and assurance really matter.
 - Use other techniques for the remainder.
 - (Assuming we can make logical arguments for separation and isolation of such components...)



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