Ad Hoc Data: An Opportunity for Domain-Specific Languages

> Kathleen Fisher AT&T Labs Research



Joint work with Anne Rogers and Bob Gruber

Abundance of valuable ad hoc data

- Call-detail records (fixed-width binary records)
 - Are these calls typical for *this* customer?
 - Are these two numbers owned by the same person?
- Provisioning data (per-order ASCII event sequences)
 - How long does supplier X take to fulfill orders?
 - How many orders sent to supplier X end up being fulfilled?
- Billing system audits (thousands of Cobol data files)
 Are we paying appropriate taxes in all jurisdictions?
- Internet data (variable-width bin. packets, ASCII messages)
 Can this packet cause a buffer overflow?

Technical challenges

- Data analysts vary widely in programming ability.
- Data arrives "as is."
 - Format determined by data source, not consumers.
 - Documentation is often out-of-date or nonexistent.
 - Some percentage of data is "buggy."
- Often data sources have high volume.
 - Data may not fit into main memory.
 - Data may contain large number of records and "entities of interest."
 - Processing must detect *relevant* errors and respond in *application-specific* ways.

Why might a language help?

- Languages provide very expressive ways of specifying what to do with data, e.g., SQL, XQuery.
- By providing infrastructure, data-processing languages
 - Enable a broader class of people to manipulate data effectively.
 - Shift focus from "How can I compute the information I want?" to "What information do I want?"
 - Shorten programs: easier to write, read, maintain, and reason about.
 - Facilitate error detection.
- Declarative data-specification languages enable generation of a wide variety of tools for manipulating data.

Overview

- Introduction
- Thesis: domain-specific languages facilitate data processing.
 - Hancock: A domain-specific data-processing language for consuming streams of transaction records to maintain customer *signatures*.
 - PADS: A declarative data-specification language for describing physical data formats.
- Conclusions

Hancock: Support for whole data analysis

Individualized analysis: Signatures

- Anomaly detection: fraud, access arbitrage, etc.
- Classification problems: target marketing, biz/res, etc.

Technical challenge:

 Massive data sets and real-time queries ⇒ Hard I/O and storage requirements ⇒ Complex programs (hard to write, read, and maintain).

Solution:

• A system that reduces the complexity of signature programs.

Processing transactions



Evolution of fraud detection

Country-based thresholds:

- Aggregate calls in 1/4/24 hour windows.
- Compare aggregates to fixed thresholds.
- Exclude common false positives.

International signatures:

- Signature is an evolving profile.
- Match calls against the customer's and known fraud signatures.

Domestic signatures?

• Much larger scale...

Problem scale



Efficiently managing communications-scale data requires substantial programming expertise.



Hancock

- Identified abstractions for processing large data streams.
 - Iterated design, meeting with data analysts to get feedback, buy-in.
 - "Wow, you can talk about the things that matter!"
- Embodied these abstractions in Hancock, a C-based domain-specific programming language.

- Embedding avoided reinventing the wheel, fit user's comfort zone.

• Built experimental and production signatures using a number of different data streams.

⁻ All AT&T signature programs are now written in Hancock.

Abstraction overview



Question: Is this calling behavior normal for *this* number? **Approach**: Compute evolving signature for each number to capture normal behavior.

Concrete results



Communities of Interest

Question: Are two numbers owned by the same person? **Approach**: Compare calling circles:



- Core number
- Inbound calls
- Outbound calls

Hancock summary

- Hancock provides infrastructure for signature programs.
 - Language model reifies abstractions described by analysts.
 - Data analysts can easily write efficient signature programs.
 - Programs highlight per-customer computation.
 - Program brevity makes them easy to write, read, maintain, and reason about.
- Embedding Hancock in C
 - Avoided having to design an entire language from scratch.
 - Worked within comfort zone of users.
- Paper: March 2004 issue of TOPLAS

Overview

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PADS System (In Progress)

One person writes declarative description of data source:

- Physical format information
- Semantic constraints

Many people use description and generated library.

- Description serves as maintainable documentation.
- Semantic constraints allow library to detect data errors.
- Bonus: From declarative specification, we can generate (many) auxiliary tools.

PADS applications

- Facilitating Hancock stream descriptions
- Helping statisticians analyze telecom provisioning data
 Replacing brittle awk/perl scripts
- Auditing billing systems
 - Automatically converting Cobol data
- Analyzing internet packets for conformance to specification
- Loading data into database/stream management systems?

PADS architecture



PADS language

- Can describe ASCII, EBCDIC (Cobol), binary, and mixed data formats.
- Type-based model: each type indicates how to process associated data.
 - Provides rich and *extensible* set of base types.
 - Pa_int8, Pa_uint8, Pb_int8, Pb_uint8, Pint8, Puint8
 - Pstring(:term-char:), Pstring_FW(:size:), Pstring_ME(:reg_exp:)
 - Supports user-defined compound types to describe data source structure:
 Pstruct, Parray, Punion, Ptypedef, Penum
 - Allows arbitrary boolean constraint expressions to describe expected properties of data.

Simple example: CLF web log

• Common Log Format from *Web Protocols and Practice*.

207.136.97.50 - - [15/Oct/1997:18:46:51 -0700] "GET /turkey/amnty1.gif HTTP/1.0" 200 3013

- Fields:
 - IP address of remote host
 - Remote identity (usually '-' to indicate name not collected)
 - Authenticated user (usually '-' to indicate name not collected)
 - Time associated with request
 - Request (request method, request-uri, and protocol version)
 - Response code
 - Content length

Example: Parray

```
Parray host {
   Puint8[4]: Psep(`.') && Pterm(` ');
};
```

207.136.97.50 - - [15/Oct/1997:18:46:51 -0700] "GET /turkey/amnty1.gif HTTP/1.0" 200 3013

Array declarations allow the user to specify:

- Size (fixed, lower-bounded, upper-bounded, unbounded)
- **Psep**, **Pterm**, and termination predicates
- Constraints over sequence of array elements

Array terminates upon exhausting EOF/EOR, reaching terminator, reaching maximum size, or satisfying termination predicate.

Example: Pstruct

Precord Pstruct http weblog {

- ' '; auth id remoteID; /- Remote identity

};

- " ["; Pdate(:']':) date; /- Timestamp of request
- "] "; http request request; /- Request
- ' '; Puint16 FW(:3:) response; /- 3-digit response code
- ' '; Puint32 contentLength; /- Bytes in response

- host client; /- Client requesting service
- ' '; auth id auth; /- Name of authenticated user

207.136.97.50 - - [15/Oct/1997:18:46:51 -0700] "GET /turkey/amnty1.gif HTTP/1.0" 200 3013

PADS compiler

- Converts description to C header and implementation files.
- For each built-in/user-defined type:
 - Functions (read, write, initialize, cleanup, copy, ...)
 - In-memory representation
 - Mask (check constraints, set representation, suppress printing)
 - Parse descriptor
- Reading invariant: If mask is check and set and parse descriptor reports no errors, then in-memory representation satisfies all constraints in data description.

Example: Reading CLF web log

```
PDC t *pdc;
http weblog entry;
http weblog m mask;
http weblog pd pd;
P open(&pdc, 0 /* PADS disc */, 0 /* PADS IO disc */);
P IO fopen(pdc, fileName);
... call init functions ...
http weblog mask(&mask, PCheck & PSet);
while (!P IO at EOF(pdc)) {
   http weblog read(pdc, &mask, &pd, &entry);
   if (pd.nerr != 0) { ... Error handling ... }
   ... Process/query entry ...
};
... call cleanup functions ...
P IO fclose(pdc);
P close(pdc);
```

PADS: Value-added tools

- Accumulators collect "bird's eye" view of data source (per field percentage of errors, histogram of "Top N")
 - Billing audit: which feeds are interesting/changing/buggy?
 - CLF data: book specification is wrong.
- Interface with Galax implementation of XQuery
 - From PADS description, generate instance of Galax data API.
 - Provisioning data: questions expressible as XQueries (without translating data into XML).
- Canonical translation into XML, including XSchema.
- Web-based data selection programs/general data browser.
- In-memory representation "completion" functions.
- Sanitized test data generation.

PADS: To do list

- Finalize initial release: documentation and release process.
- Conduct careful performance study and tune library accordingly.
- Leverage semantic information to build value-added tools.
- Allow library generation to be customized with application-specific information:
 - Repair errors, ignore fields, customize in-memory representation, *etc*.

Related work

• ASN.1, ASDL

- Describe logical representation, generate physical.

- DataScript [Back: CGSE 2002] & PacketTypes [McCann & Chandra: SIGCOMM 2000]
 - Binary only
 - Stop on first error
- Database vendors have tools to load specific formats.
- YACC, etc.
- Hand-written parsers in C, perl, etc.

PADS summary

- Data analysts vary widely in programming ability.
 PADS supports declarative programming, automatic tool generation.
- Data arrives "as is."
 - Format determined by data source, not consumers.
 - PADS language allows consumers to describe data as it is.
 - Documentation is often out-of-date or nonexistent.
 - PADS description can serve as documentation for data source.
 - Some percentage of data is "buggy."
 - Constraints allow consumers to express checked expectations about data.
- Often streams have high volume.
 - Data may not fit into main memory.
 - Multiple entry-points allow different levels of granularity.
 - Processing must detect *relevant* errors (without necessarily halting program)
 - Masks specify relevancy; returned descriptors characterize errors.

Domain Specific Languages

- Facilitate data processing
- Domain-specific abstractions
 - Enable broader class of users to manipulate data.
 - Shorten and simplify user code, improving maintainability.
 - Facilitate error detection.
 - Enable many useful tools.
- Different paradigms useful
 - Procedural vs. declarative
 - Application vs. library
- Embedded DSLs
 - Allow fast prototyping.
 - Avoid duplicating existing functionality.
 - May impede analysis longer term.
 - Can facilitate acceptance by users.

Users are key

- Involve users from the beginning.
 - They understand the domain, its constraints, and what functionality they require.
 - Invaluable input when evaluating the many trade-offs that arise in designing a language.
 - Language will only be successful *if they use it*, so getting them to "buy-in" early is crucial.

Summary

- Hancock (Available online from <u>www.research.att.com/projects/hancock</u>):
 - Domain-specific language for processing arbitrary streams of fixedwidth data.
- PADS (In progress. Some information: <u>www.research.att.com/projects/pads</u>):
 - Declarative description of data source, including both layout information and semantic constraints.
 - Compiler generates data-manipulation library.
 - In progress: Suite of tools to leverage declarative specification.

Why not use C / Perl / Shell scripts...?

- Writing hand-coded parsers is time consuming & error prone.
- Reading them a few months later is difficult.
- Maintaining them in the face of even small format changes can be difficult.
- Programs break in subtle and machine-specific ways (endien-ness, word-sizes).
- Such programs are often incomplete, particularly with respect to errors.

Why not use traditional parsers?

- Specifying a lexer and parser separately can be a barrier.
- Need to handle data-dependent parsing.
- Need more flexible error processing.
- Need support for multiple-entry points.



PADS will be available shortly for download with a noncommercial-use license.

http://www.research.att.com/projects/pads

Example: arrays and unions

```
Parray nIP {
 Puint8 [4] : Psep('.');
};
Parray sIP {
  Pstring SE(:"[. ]":) [] : Psep('.') && Pterm ('.');
}
Punion host {
 nIP resolved; /- 135.207.23.32
  sIP symbolic; /- www.research.att.com
};
Punion auth id {
  Pchar unauthorized : unauthorized == '-';
                 /- non-authenticated http session
  Pstring(:' ':) id;
                 /- login supplied during authentication
};
```

207.136.97.50 - - [15/Oct/1997:18:46:51 -0700] "GET /turkey/amnty1.gif HTTP/1.0" 200 3013

Generated type declarations

```
typedef struct {
 host client; /* Client requesting service */
  auth id remoteID; /* Remote identity */
} http weblog;
typedef struct {
 host m client;
  auth id m remoteID;
} http weblog m;
typedef struct {
  int nerr;
  int errCode;
 PDC loc loc;
  int panic;
 host pd client;
  auth id pd remoteID;
 ...;
} http weblog pd;
```

Generated accumulator representation (acc)

```
pstruct http_request {
    '\"'; http_method meth; /- Request method
    ''; Pstring(:'':) req_uri; /- Requested uri.
    ''; http_version version : check(version, meth);
    '\"';
};
```

typedef struct {

	PDC_uint64_acc	nerr;
	http_method_acc	meth;
	PDC_string_acc	req_uri;
	http_version_acc	version;
}	<pre>http request acc;</pre>	

Generated read function





We also generate initialization and cleanup functions for representations and error descriptors for variable width data.

Example: Punion

```
Punion id {
   Pchar unavailable : unavailable == '-';
   Pstring(:' ':) id;
};
```

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- Union declarations allow the user to describe variations.
- Implementation tries branches in order.
- Stops when it finds a branch whose constraints are all true.
- Switched version branches on supplied tag.

Advanced features: User constraints

```
int checkVersion(http v version, method t meth) {
 if ((version.major == 1) && (version.minor == 0)) return 1;
 if ((meth == LINK) || (meth == UNLINK)) return 0;
 return 1;
}
Pstruct http request {
  '\"'; method t meth; /- Request method
  ' '; Pstring(:' ':) req uri; /- Requested uri.
 ''; http v version : checkVersion(version, meth);
                         /- HTTP version number of request
  '\"';
};
```

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Advanced features: Sharing information

- "Early" data often affects parsing of later data:
 - Lengths of sequences
 - Branches of switched unions
- To accommodate this usage, we allow PADS types to be parameterized:

```
Punion packets_t (: Puint8 which, Puint8 length:) {
    Pswitch (which) {
        Pcase 1: header_t header;
        Pcase 2: body_t body;
        Pcase 3: trailer_t trailer;
        Pdefault: Pstring_FW(: length :) unknown;
};
```

Generated representation

Pstruct http request { '\"'; http method meth; /- Request method ' '; Pstring(:' ':) req uri; /- Requested uri. ' '; http version version : check(version, meth); '\"'; };

typedef struct {

- http method meth; /* Request method */
 - Pstring req uri; /* Requested uri */
- } http request;

- http version version; /* check(version, meth) */



typedef struct {

- P base m structLevel;
- http method m meth;
- P base m req uri; /* Check, Set, Print,... */
- http version m version;
- } http request m;

Generated parse descriptor (pd)

```
Pstruct http_request {
    '\"'; http_method meth; /- Request method
    ''; Pstring(:'':) req_uri; /- Requested uri.
    ''; http_version version : check(version, meth);
    '\"';
};
```

typedef struct {

int	nerr;
P_errCode_t	errCode;
P_loc_t	loc;
int	panic;
http_method_pd	meth;
P_base_pd	req_ uri;
http_version_pd	version;
<pre>http_request_pd;</pre>	

/* Structural error */

Expanding our horizons

- Design to this point primarily based on our experiences with Hancock data streams.
- To get feedback on the expressiveness of our language and the utility of the generated libraries, we started collecting other users:
 - Data analysts (Chris Volinksy, David Poole)
 - Cobol gurus (Andrew Hume, Bethany Robinson, ...)
 - Internet miner (Trevor Jim)

Data analysts: The domain

• Data: ASCII files, several gigabytes in size, sequence of provisioning records, each of which is a sequence of state, time-stamp pairs.

customer_id | order# | state1 | ts1 | ... | staten | tsn

• Desired application: aggregation queries.

. . .

- How many records that go through state 3 end in state 5?
- What is the average length of time spent in state 4?

• Original applications written in brittle awk and perl code.

Data analysts: Lessons learned

- PADS expressive enough to describe data sources.
- PADS descriptions much less brittle than originals.
- Generated code faster than original implementations.
- Support for declarative querying could be a big win, as analysts can then generate desired aggregates with only declarative programming.
 - Data and desired queries "semi-structured" rather than relational.

Developing support for declarative querying

- Wanted to leverage existing query language.
- XQuery, standardized query language for XML, is appropriately expressive for our queries.
- Modified Galax, open-source implementation of XQuery, to create data API, allowing Galax to "read" data not in XML format. (Mary Fernandez, Jérôme Simeon).
- Extended PADSC to generate data API.
- Currently, can run queries over PADS data if data fits in memory.
- Next: extend Galax API and generated library to support streaming interface.

Cobol gurus: The domain

- Data: EBCDIC encoded files with Cobol copy-book descriptions. Thousands of files arriving on daily basis, hundreds of relevant copy-books, with more formats arriving regularly.
- Desired application: developing a "bird's eye" view of data.
- Original applications: hand-written summary programs on an *ad hoc* basis.

Cobol gurus: Lessons learned

- Wrote a translator that converts from Cobol copy books to PADSL.
 - Added Palternates, which parse a block of data multiple times, making all parses available in memory.
 - Generated description uses **Palternates**, **Parrays**, **Pstructs**, and **Punions**.
 - Uses parameters to control data-dependent array lengths.
 - Successfully translated available copybooks.
- Added *accumulators* to generated library.
 - Aggregate parsed data, including errors.
 - Generate reports, providing bird's eye view.

Internet miner: The domain

- Data: data formats described by RFCs. Some binary (dns, for example), some ASCII (http, for example).
- Desired applications: detecting security-related protocol violations, data mining, semi-automatic generation of reference implementations.
- Original applications: being developed with PADS.

Internet miners: Lessons learned

- Constraints are a big win, because they allow semantic conditions to be checked.
- Parameterization used a lot, particularly to check buffer lengths in binary formats (dns).
- Additional features needed:
 - Predicates over parsed prefixes to express array termination.
 - Regular expression literals and user-defined character classes.
 - Positional information in constraint language.
 - Recursive declarations.
 - In-line declarations.
- Developing script to semi-automate translation of RFC EBNF (Trevor Jim).
 - URI spec successfully translated, HTTP close.