

Deep Specifications and Certified Abstraction Layers

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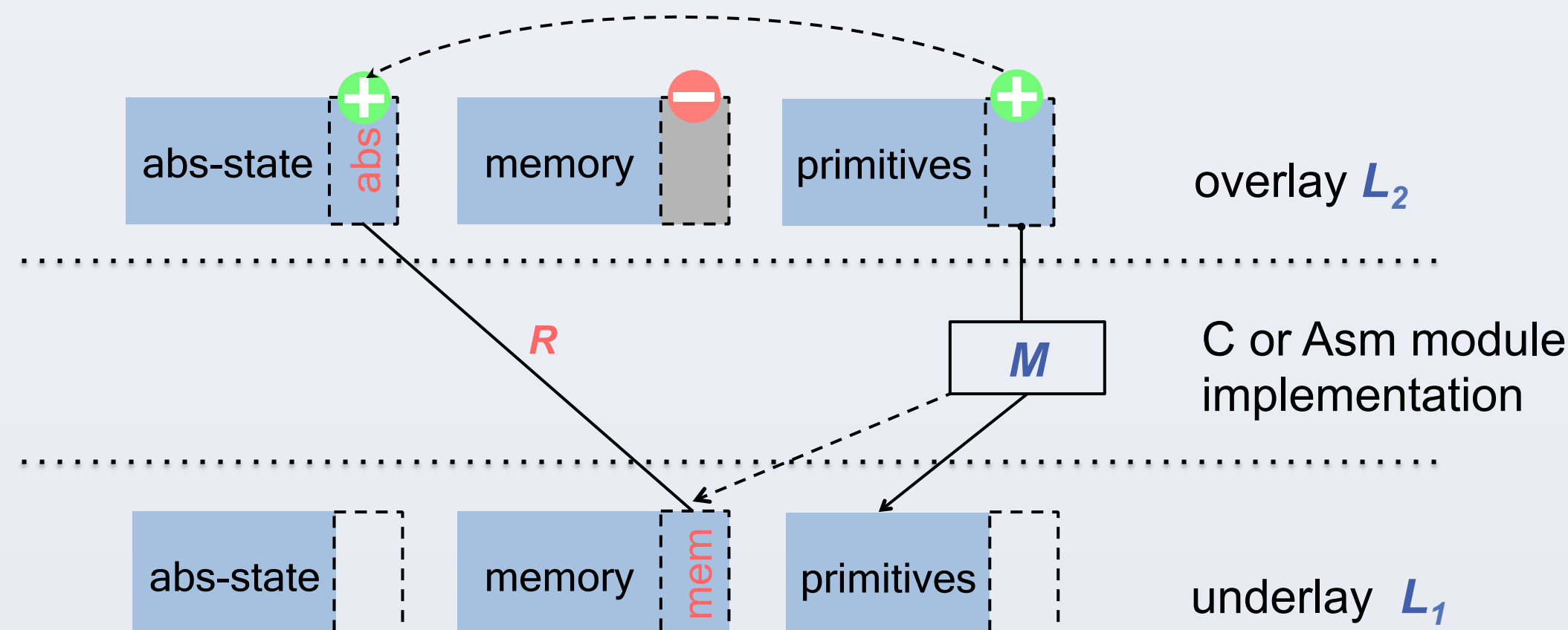
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Contributions

1. Present the first language-based account of certified abstraction layers and show how they correspond to a rigorous form of abstraction over deep specifications used widely in the system community.
2. Provide a layer calculus showing how to formally specify, program, verify, and compose certified abstraction layers.
3. Instantiate the layer calculus on top of two core languages: *ClightX*, a variant of the CompCert Clight language; and *LAsm*, an x86 assembly language.
4. Extend CompCert to build a new verified compiler, CompCertX, that can compile *ClightX* abstraction layers into *LAsm* layers.
5. Construct several feature-rich certified OS kernels in Coq. The hypervisor consists of 5500 lines of C and x86 assembly, and can boot a version of Linux as a guest.

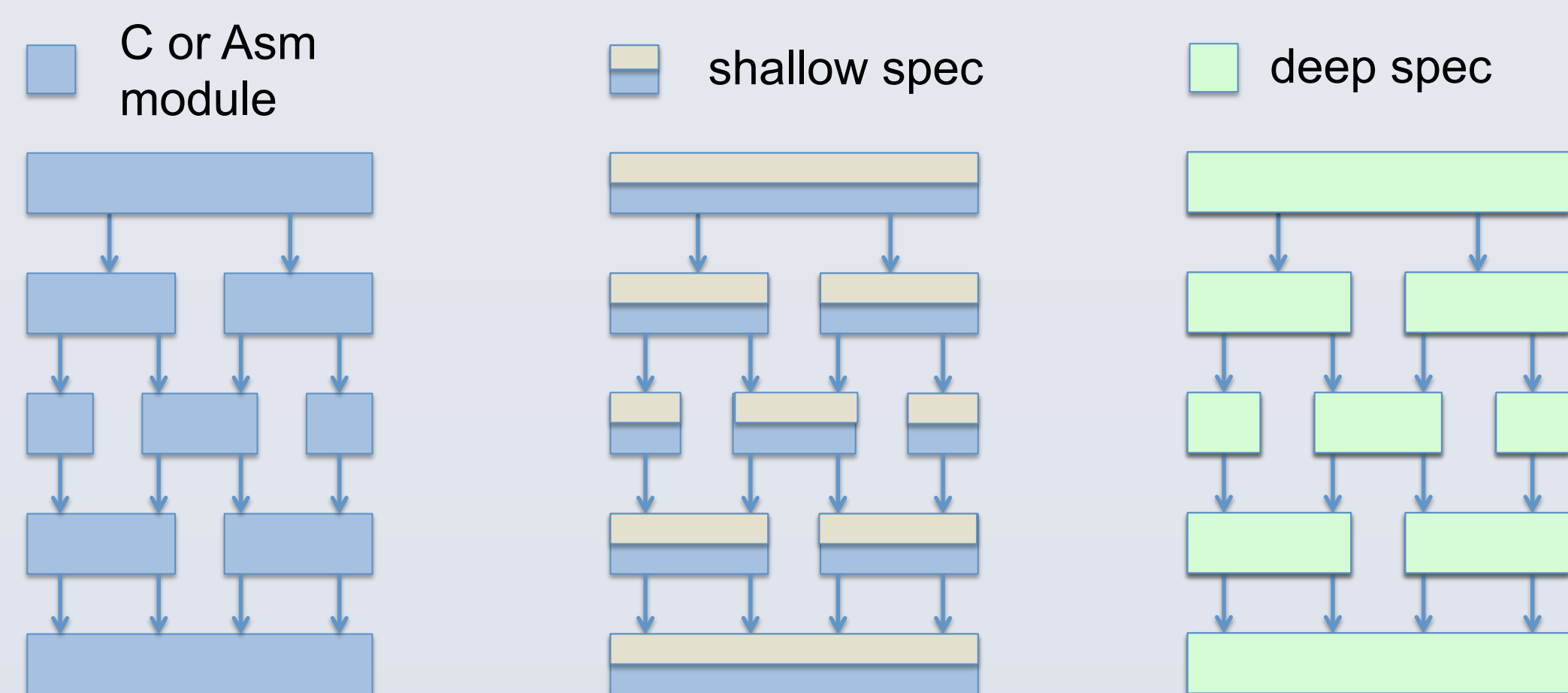
Abstraction Layer

- An abstraction layer is a triple (L_1, M, L_2) .
- The module M implements the overlay interface L_2 on top of underlay L_1 .

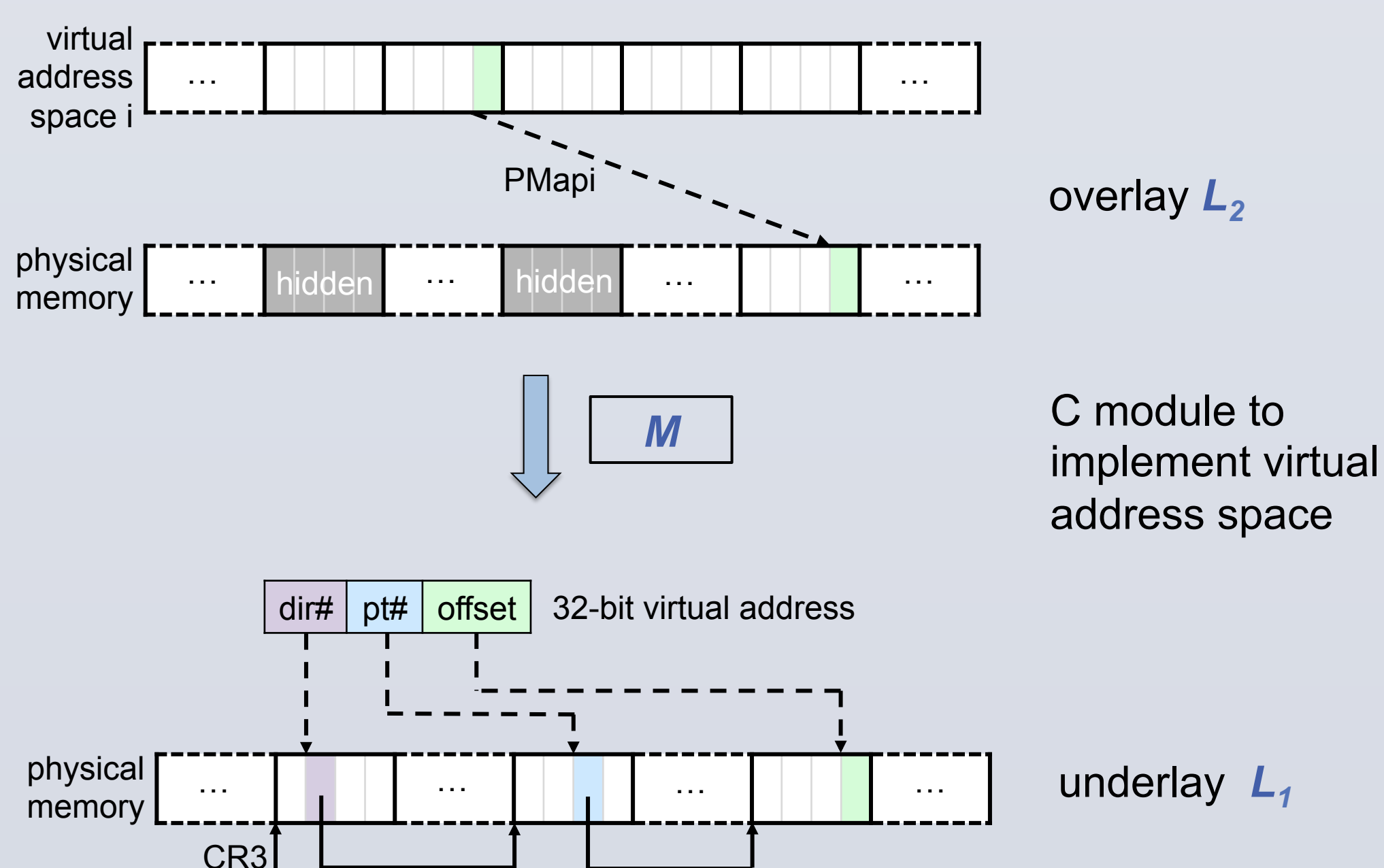


Deep Specification

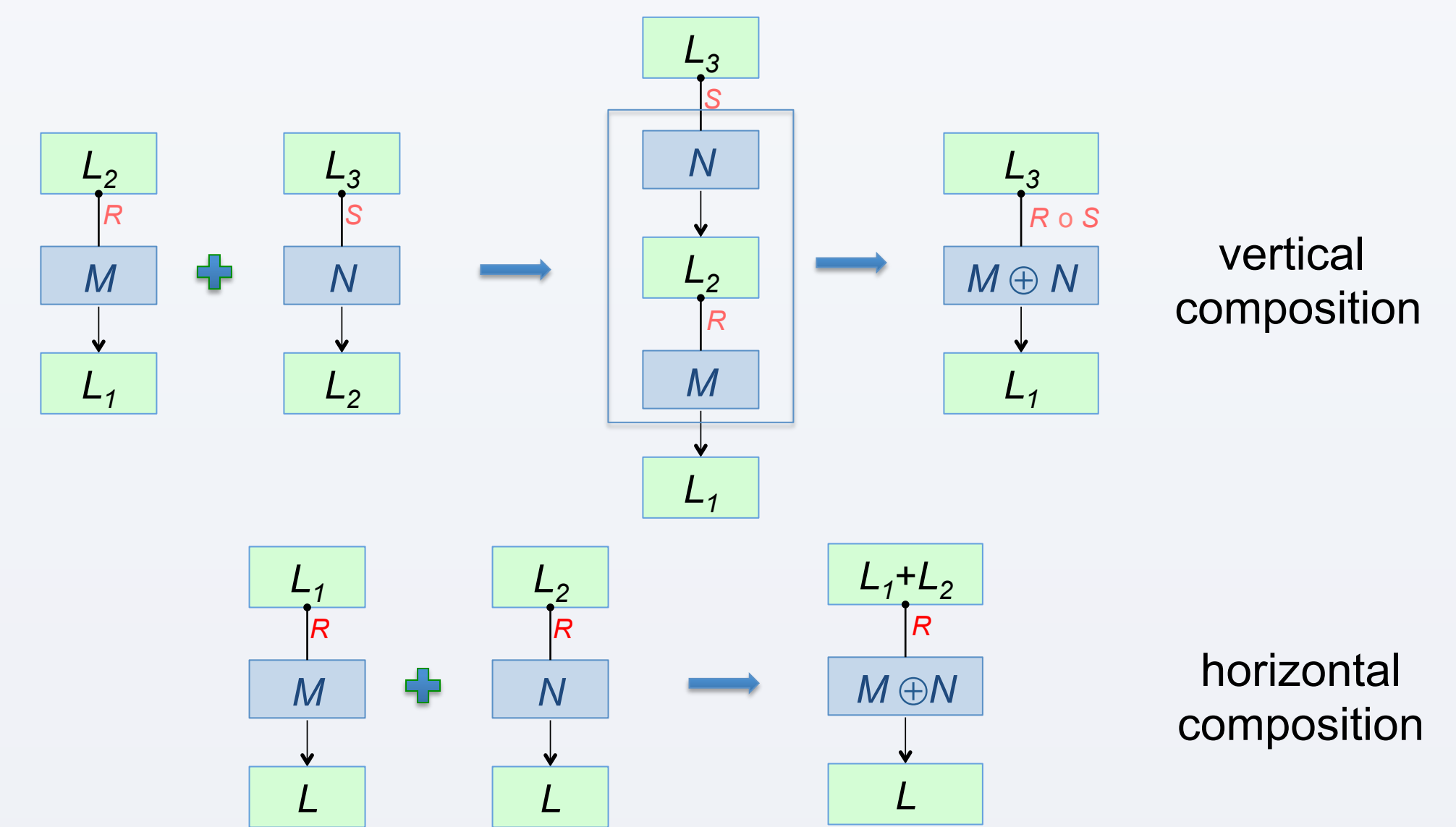
- L_2 is a deep specification of M over L_1 if under any valid program context P of L_2 , the whole-program semantics $\llbracket P \oplus M \rrbracket (L_1)$ and $\llbracket P \rrbracket (L_2)$ are observationally equivalent.
- Deep specification captures all we need to know about a module M .
- Any two implementations of the same deep spec are contextually equivalent.



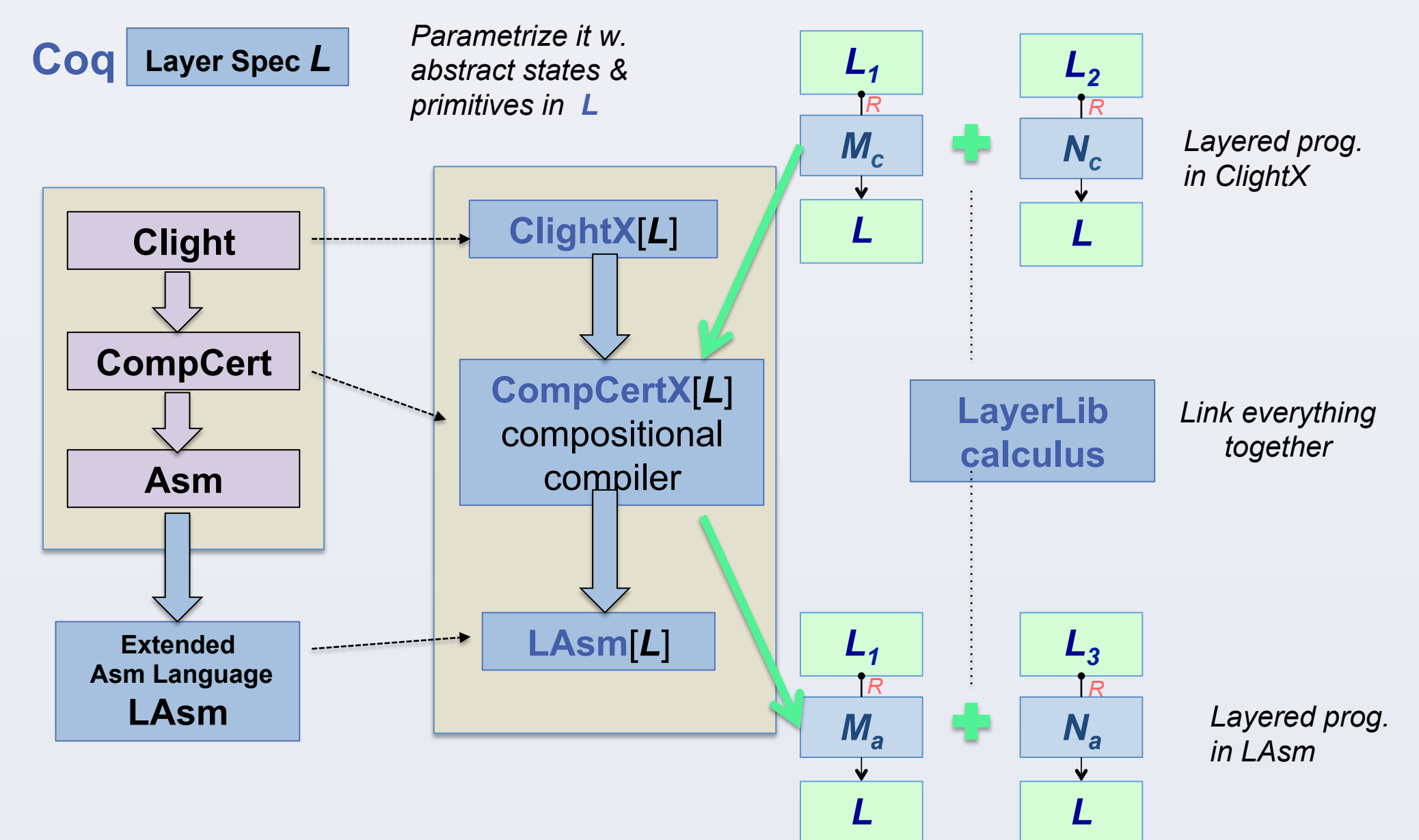
Example: Page Map and MemoryMode



Layer Calculus



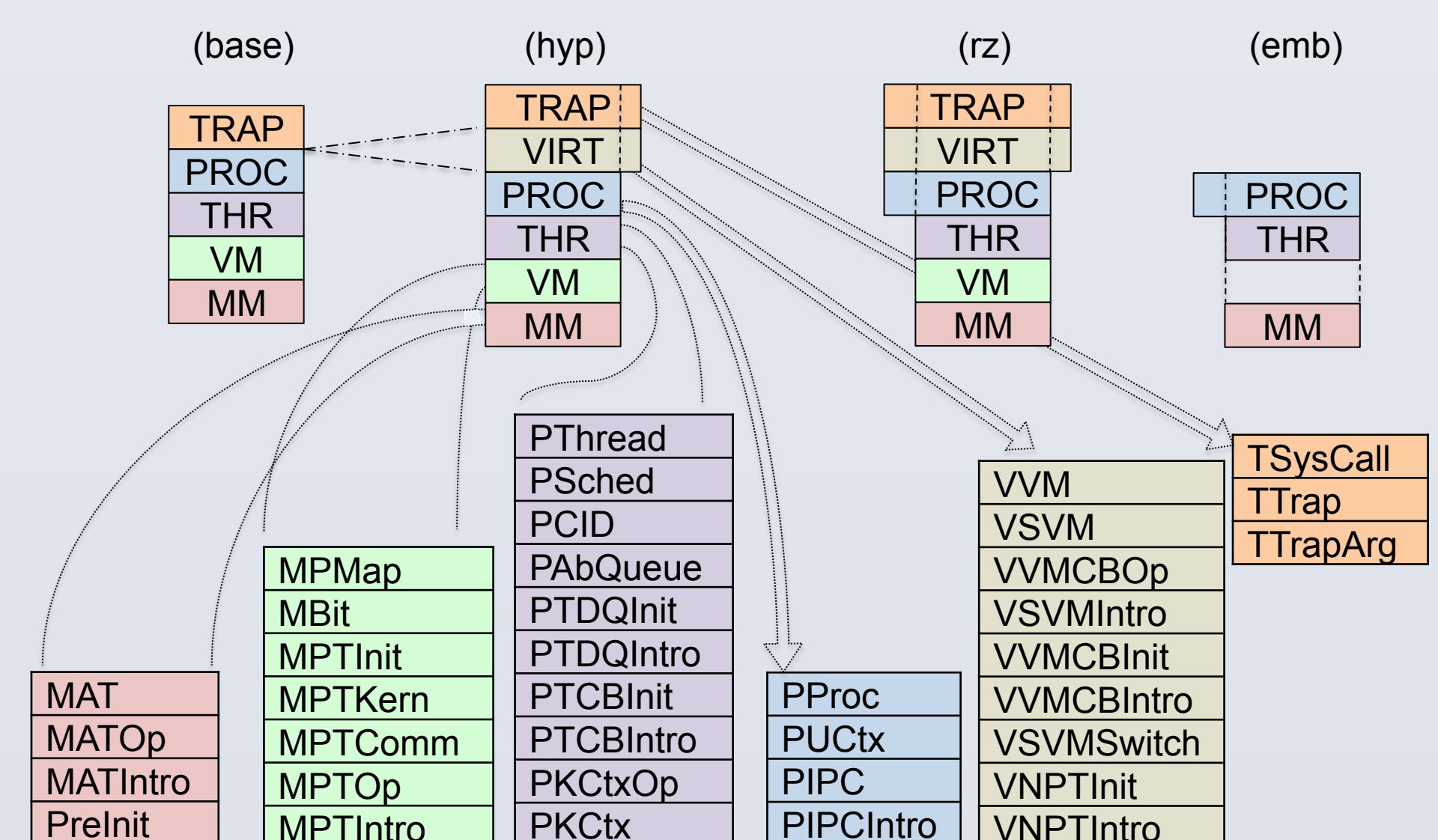
Programming & Compiling Layers



Variants of mCertiKOS Kernels

➤ Final theorem for mCertiKOS_hyp:

$$\forall P, \llbracket P \oplus \text{CompCertX}(\text{mCertiKOS_hyp}) \rrbracket (\text{Prelit}) \leq \llbracket P \rrbracket (\text{TSysCall})$$



Performance

