

Evolving Verified Cloud Authorization

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Formal verification enables faster evolution of critical systems







Policies

```
{
 "Version": "2012-10-17",
  "Statement":
    Ł
      "Effect": "Allow",
      "Principal": {
        "AWS": "1111111111111"
      },
      "Action": "s3:GetObject",
      "Resource": "arn:aws:s3:::my-photos-bucket/*",
      "Condition": {
        "Bool": {
          "aws:SecureTransport": "true"
        }
}
```

Requests

```
{
    "Principal": {
        "AWS": "1111111111"
    },
    "Action": "s3:GetObject",
    "Resource": "arn:aws:s3:::my-photos-bucket/cats.jpg",
    "aws:SecureTransport": true,
    "aws:CurrentTime": "2022-05-16T01:02:03Z",
    ...
}
```



The authorization problem

Policies + Request = Allowed



"Principal": { "AWS": "111111111111" },

{

} ...

"Action": "s3:GetObject", "Resource": "arn:aws:s3:::my-photos-bucket/cats.jpg", "aws:SecureTransport": true, "aws:CurrentTime": "2022-05-16T01:02:03Z",

Yes/No



AWS LAM Half a billion API calls/second

a=

H



```
{
  "Version": "2012-10-17",
  "Statement":
      "Effect": "Allow",
      "Principal": {
        "AWS": "1111111111111"
     },
      "Action": "s3:GetObject",
      "Resource": "arn:aws:s3:::my-photos-bucket/*",
      "Condition": {
        "Bool": {
          "aws:SecureTransport": "true"
        }
}
```



```
{
  "Version": "2012-10-17",
  "Statement":
      "Effect": "Allow",
      "Principal": {
        "AWS": "1111111111111"
      },
      "Action": "s3:GetObject",
      "Resource": "arn:aws:s3:::my-photos-bucket/*",
      "Condition": {
        "Bool": {
          "aws:SecureTransport": "true"
        },
        "StringEquals": {
          "aws:SourceVpc": "vpc-123"
```







```
{
  "Version": "2022-10-17",
  "Statement":
    Ł
      "Effect": "Allow",
      "Principal": {
        "AWS": "1111111111111"
      },
      "Action": "s3:GetObject",
      "Resource": "arn:aws:s3:::my-photos-bucket/*",
      "Condition": {
        "OR"
          { "Bool": { "aws:SecureTransport": "true" } },
          { "StringEquals": { "aws:SourceVpc": "vpc-123" } }
```



Adding "OR"

Pros

• More direct, succinct policies

Cons

- Availability risk
- Security risk





Testing

What we can do

- Fuzzing
- Historical data
- Side-by-side execution

We can't test exhaustively

- Language is complex (111 pages of guide)
- Single request can have dozens of policies and dozens of variables



Formal verification enables faster evolution of critical systems



Formal Verification Method



Verified Implementation

Idiomatic Compiler





Running Example

StringEqualsIgnoreDashes("ABC", "ABC")
StringEqualsIgnoreDashes("A-B-C", "ABC")
StringEqualsIgnoreDashes("ABC", "A-B-C")
StringEqualsIgnoreDashes("----", "")



Verified Implementation

Idiomatic Compiler

- 1. Specification
- 2. Theorems
- 3. Implementation



1. Specification

```
predicate AllDashes(s: string) {
1
2
      s == [] || (s[0] == '-' && AllDashes(s[1..]))
3
    }
4
5
    predicate EqualsIgnoreDashes(s: string, t: string) {
    if s == [] then AllDashes(t)
 6
7 else if t == [] then AllDashes(s)
      else if s[0] == '-' then EqualsIgnoreDashes(s[1..], t)
8
      else if t[0] == '-' then EqualsIgnoreDashes(s, t[1..])
9
      else s[0] == t[0] && EqualsIgnoreDashes(s[1..], t[1..])
10
11
    }
```

2. Theorems

lemma EqualsIgnoreDashesRefl(s: string) ensures EqualsIgnoreDashes(s, s)

lemma EqualsIgnoreDashesSym(s: string, t: string)
 requires EqualsIgnoreDashes(s, t)
 ensures EqualsIgnoreDashes(t, s)

lemma EqualsIgnoreDashesTrans(s: string, t: string, u: string)
 requires EqualsIgnoreDashes(s, t)
 requires EqualsIgnoreDashes(t, u)
 ensures EqualsIgnoreDashes(s, u)



2. Theorems

lemma EqualsIgnoreDashesAppend(s1: string, t1: string, s2: string, t2: string)
 requires EqualsIgnoreDashes(s1, s2)
 requires EqualsIgnoreDashes(t1, t2)
 ensures EqualsIgnoreDashes(s1 + t1, s2 + t2)



2. Theorems

```
lemma EqualsIgnoreDashesAppend(s1: string, t1: string, s2: string, t2: string)
 requires EqualsIgnoreDashes(s1, s2)
 requires EqualsIgnoreDashes(t1, t2)
 ensures EqualsIgnoreDashes(s1 + t1, s2 + t2)
 if s1 == [] {
   assert s1 + t1 == t1;
   assert AllDashes(s2);
   if s2 == [] {
     assert s^2 + t^2 == t^2;
   } else {
     EqualsIgnoreDashesAppend(s1, t1, s2[1..], t2);
     EqualsIgnoreDashesRightAdd(s1 + t1, s2[1..] + t2);
     assert s2 + t2 == ['-'] + (s2[1..] + t2);
   }
 } else {
   if s1[0] == '-' {
     assert EqualsIgnoreDashes(s1[1..] + t1, s2 + t2);
     assert s1 + t1 = ['-'] + (s1[1..] + t1);
   } else {
     if s2 == [7] {
     } else {
       if s2[0] == '-' {
         assert s2 + t2 == ['-'] + (s2[1..] + t2);
       } else {
         assert s1 + t1 = [s1[0]] + (s1[1..] + t1);
         assert s^2 + t^2 = [s^2[0]] + (s^2[1..] + t^2);
```

3. Implementation

ensures res == EqualsIgnoreDashes(s: string32, t: string32) returns (res: bool)

var i := 0 as nat32; var j := 0 as nat32;

var slength := strLen(s); var tlength := strLen(t):

i.< i < slength && j < tlength invariant i <= slength && j <= tlength invariant EqualsIgnoreDashes(s[..i], t[..j]) invariant EqualsIgnoreDashes(s, t) <==> EqualsIgnoreDashes(s[i..], t[j..]); decreases |s| + |t| - i as nat - j as nat

t if sfir

```
cqualsIgnoreDashesAppend(s[...], [s[...], t[...j], []);
assert s[..i] + [s[i]] == s[..i+1];
assert t[..j] + [] == t[..j];
```

```
assert s[..i] + [] == s[..i];
assert t[..j] + [t[j]] == t[..j+1];
```

```
} i + 1;
} else if sur
```

else if sum

return false;

```
[....signoreDashesAppend(s[...], [[i]], t[..j], [t[j]]);
assert s[..i] + [s[i]] == s[..i+1];
assert t[..j] + [t[j]] == t[..j+1];
i = i + 1;
j := j + 1;
} else {
```





}

3. Implementation

```
newtype nat32 = x | 0 <= x <= 0x7fff_fff</pre>
```

```
type string32 = x: string | 0 <= |x| <= 0x7fff_fff</pre>
```

```
function method {:javainline "$s.length()"} strLen(s: string32): (res: nat32)
    ensures res == |s| as nat32
```



Verified Implementation

Idiomatic Compiler



```
...c static boolean equalsIgnoreDashes(
dafny.DafnySequence<? extends Character> s,
dafny.DafnySequence<? extends Character> t
boolean res = talse,
int _277_i;
277_i = 0;
int _278_j;
278_j = 0;
int _279_slength
_279_slength = ___default.strLen(s);
int _280_tlengt
_280_tlength / __default.strLen(t);
 hile ((Int or
                compareUnsigned(_277_i, _279_slength) < 0)</pre>
      a& (intege, compareUnsign a, 278_j, _280_tlength) < 0)) {</pre>
  if (((s) select(.277_i)) == ( -'))
    277_i = (int) ((277_i) (1));
  } else f ((((t).s lect(_278_j)) ==
                                      '-')) {
    278_ = (int) ((278_j) + (1));
  } else f (((s).s(lect(_277_i)) == (t).select(_278_j))) {
    _277_ = (int) ((_277_i) + (1));
    _278_j = (int) ((_278_j) (1));
  } else {
    res = f_{\alpha} se;
    return res;
  3
7
while (Integer.compareUnsigned(_277_i, _279_slength) < 0) {</pre>
  if (((s).select(_277_i)) != ('-')) {
    res = false;
    return res;
  }
  _277_i = (int) ((_277_i) + (1));
3
while (Integer.compareUnsigned(_278_j, _280_tlength) < 0) {</pre>
 if (((t).select(_278_j)) != ('-')) {
    res = false;
    return res;
  278_j = (int) ((278_j) + (1));
3
res = true;
return res;
```



```
static boolean equalsIgnoreDashes(String s, String t) {
  int i = 0;
  int j = 0;
  int slength = s.length();
  int tlength = t.length();
  while (i < slength && j < tlength) {</pre>
    if (s.charAt(i) == '-') {
     i = i + 1;
    } else if (t.charAt(j) == '-') {
      j = j + 1;
    } else if (s.charAt(i) == t.charAt(j)) {
     i = i + 1;
      j = j + 1;
    } else {
      return false;
    }
  }
  while (i < slength) {</pre>
    if (s.charAt(i) != '-') {
      return false;
    }
    i = i + 1;
  }
  while (j < tlength) {</pre>
    if (t.charAt(j) != '-') {
      return false;
    }
    j = j + 1;
  }
  return true;
}
```

Pros

- Lowers the risk of experimentation
- Generated code is reviewed as usual
- Can be directly changed during operational event

Cons

- Language restricted (Dafny-Lite)
- Exceptions are hard



Other languages?

- Rust
 - Exception headaches go away entirely
 - Modeling lifetimes in Dafny is challenging



Verified Implementation

Idiomatic Compiler



Specification

```
predicate Allbashes(s: string) {
1
  s == [] | (s[0] == '-' & AllDashes(s[1..]))
2
3
    }
4
5
    predicate EqualsIgnoreDashes(s: string, t: string) {
    if s == [] then AllDashes(t)
 6
7
   else if t _____ then AllDashes(s)
      else ir s[0] == '-' then EqualsIgnoreDashes(s[1..], t)
8
      else if t[0] == '-' then EqualsIgnoreDashes(s, t[1..])
9
      else s[0] == t[v] && EqualsIgnoreDashes(s[1..], t[1..])
10
11
    }
```

Specification

```
predicate method isDash(c: char) {
 2
      c == '-' || c == '_'
 4
 5
    predicate AllPasnes(s: string) {
      s == [] [ (isDash(s[0]) & AllDashes(s[1..]))
 6
 7
     }
 8
    predicate EqualsIgnoreDashes(s: string, t: string) {
 9
      if s == [] then AllDashes(t)
10
      else if t _____ then AllDashes(s)
11
      else i isDash(s[0]) then EqualsIgnoreDashes(s[1..], t)
12
      else if isDash(t[0]) then EqualsIgnoreDashes(s, t[1..])
13
      else s[0] + [0] a EqualsIgnoreDashes(s[1..], t[1..])
14
15 }
```

Implementation

method equalsIgnoreDashes(s: string32, t: string32) returns (res: bool) ensures res == EqualsIgnoreDashes(s, t) var i := 0 as nat32; var j := 0 as nat32; var slength := strLen(s); var tlength := strLen(t); while i < slength && j < tlength</pre> invariant i <= slength && j <= tlength</pre> invariant EqualsIgnoreDashes(s[..i], t[..j]) invariant EqualsIgnoreDashes(s, t) <==> EqualsIgnoreDashes(s[i..], t[j..]); decreases let _ ltl - i as nat - j as nat if s[i] == '-' { 1 - Tanona Dash ...ppend(s[...i], [s[i]], t[...j], []); assert s[..i] + [s[i]] == s[..i+1]; assert t[_____ == t[..j]; i := . + 1; } else if t[j] == '-' { Equals19. ______rend(s[..i], [], t[..j], [t[j]]); assert s[..i] + [] == s[..i]; assert t[..j] + [t[j]] == t[..j+1]; j := j + 1; } else if s[i] == t[j] { EqualsIgnoreDashesAppend(s[..i], [s[i]], t[..j], [t[j]]); assert s[..i] + [s[i]] == s[..i+1]; assert t[..j] + [t[j]] == t[..j+1]; i := i + 1;j := j + 1; } else { return false; }



Implementation

```
method equalsIgnoreDashes(s: string, t: string) returns (res: bool)
 ensures res == EqualsIgnoreDashes(s, t) {
 var i := 0;
 var j := 0;
 while i < |s| && j < |t|
   invariant i <= |s| && j <= |t|
    invariant EqualsIgnoreDashes(s[..i], t[..j])
   invariant EqualsIgnoreDashes(s, t) <==> EqualsIgnoreDashes(s[i..], t[j..]);
   decreases Isi . Itl - i - j {
   if isDash(s[i])
     LyunisignoreDashesAppend(s[..i], [s[i]], t[..j], []);
      assert s[..i] + [s[i]] == s[..i+1];
      assert t[..j] + [] == t[..j];
     i :- i - 1;
    } else if isDash(t[j]) {
      Equals_____Perbook_____end(s[..i], [], t[..j], [t[j]]);
      assert s[..i] + [] == s[..i];
      assert t[..j] + [t[j]] == t[..j+1];
     j := j + 1;
   } else if s[i] == t[j] {
      EqualsIgnoreDashesAppend(s[..i], [s[i]], t[..j], [t[j]]);
      assert s[..i] + [s[i]] == s[..i+1];
      assert t[..j] + [t[j]] == t[..j+1];
     i := i + 1;
     j := j + 1;
   } else {
      return false;
    }
 }
```

```
while i < lsl
    invariant i <= |s|
    inv_mant Equals__noreDashes(s, t) <==> EqualsIgnoreDashes(s[i..], t[j..]) {
   if lisDash(s[i])
     i := i + 1;
 }
 while j < ltl
   invariant j <= |t|
    ip:___cant tqual_IgnoreDashes(s, t) <==> EqualsIgnoreDashes(s[i..], t[j..]) {
     { !isDash(t[j]) [
        unn fals
    j := j + 1;
  3
  return true:
}
```





```
module 0ld {
  predicate EqualsIgnoreDashes(...)
}
module New {
  predicate EqualsIgnoreDashes(...)
}
module Evolve {
  lemma Ok(s: string, t: string)
    requires Old.EqualsIgnoreDashes(s, t)
    ensures New.EqualsIgnoreDashes(s, t)
 { ... }
```



```
...ma OldNewOk(s: string, t: string)
 requires Old.EqualsIgnoreDashes(s, t)
 ensures New.EqualsIgnoreDashes(s, t) {
 if s == [7]
     AllDasnesoLanenoKCCJ;
   } else {
     if s[0] == '-' {
       assert New.EqualsIgnoreDashes(s[1..], t);
     }
     else {
       if t[0] == '-' {
         New.EqualsIgnoreDashesAppend([], s, [t[0]], t[1..]);
         assert [t[0]] + t[1..] == t;
         assert [] + s == s;
      } else {
         New.EqualsIgnoreDashesRefl([s[0]]);
         New.EqualsIgnoreDashesAppend([s[0]], s[1..], [t[0]], t[1..]);
         assert [s[0]] + s[1..] == s;
         assert [t[0]] + t[1..] == t;
       }
```

```
lemma NewoldOk(s: string, t: string)
 requires New.EqualsIgnoreDashes(s, t)
 requires '_' !in s
 requires '_' !in t
 ensures Old.EqualsIgnoreDashes(s, t)
ł
 if s == [] {
   AllDashesNewOldOk(t);
 } else {
   if New.isDash(s[0]) {
      assert Old.EqualsIgnoreDashes(s[1..], t);
```





Conclusion



```
{
  "Version": "2022-10-17",
  "Statement":
    Ł
      "Effect": "Allow",
      "Principal": {
        "AWS": "1111111111111"
      },
      "Action": "s3:GetObject",
      "Resource": "arn:aws:s3:::my-photos-bucket/*",
      "Condition": {
        "OR"
          { "Bool": { "aws:SecureTransport": "true" } },
          { "StringEquals": { "aws:SourceVpc": "vpc-123" } }
```



Adding "OR"

Pros

• More direct, succinct policies

Cons

- Availability risk
- Security risk





Adding "OR"

Pros

• More direct, succinct policies

Cons

• Need to prove backward compatibility





Formal verification enables faster evolution of critical systems



Customer-facing auth coming soon...





Thank you!

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