### Logical Foundations of Cyber-Physical Systems

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#### DARPA

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Logical Foundations of Cyber-Physical Systems

# $\mathcal{R}$ Outline

### 1 CPS are Multi-Dynamical Systems

- Hybrid Systems
- Hybrid Games

### 2 Dynamic Logic for Multi-Dynamical Systems

- Syntax
- Semantics

### 3 Proofs for CPS

### 4 Theory of CPS

- Soundness and Completeness
- Differential Invariants
- Differential Radical Invariants
- 5 Applications
  - Ground Robots

### 6 Summary

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### 6 Summary

# Can you trust a computer to control physics?

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#### Rationale

- Safety guarantees require analytic foundations
- Poundations revolutionized digital computer science & society
- Need even stronger foundations when software reaches out into our physical world

## Can you trust a computer to control physics?

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#### CPS Core Question

How can we provide people with cyber-physical systems they can bet their lives on?

# $\mathcal{R}$ CPS Applications



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# $\mathcal{R}$ Correctness Questions in CPS Design

Safety The system must be safe under all circumstances Liveness The system must reach a given goal

#### How do we make cyber-physical systems safe?

Extensive testing? Code reviews? When are we done? How many test cases are enough? Did we cover all relevant tests?







# $oldsymbol{\mathcal{R}}$ Benefits of Logical Foundations for V & V

#### Proving

Safety Formalize system properties: What is "Safe"? "Reach goal"? Models Formalize system models

Assumptions Make assumptions explicit

Constraints Reveal invariants, switching conditions, starting conditions

Design Invariants guide safe controller design

Constructive Construct models along with their proof

#### Byproducts

Analyze Determine design trade-offs & feasibility early

Synthesize Turn high-level models into code & correctness monitors

Certify Proofs as artifacts for certification

#### Tools

#### KeYmaera Theorem prover for CPS

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### Diverse Application Domains

- Automotive
- Aircraft

Energy

Railway

- Robotics
- Surgery



### Various Levels

All the way from

- Engine idle control
- Adaptive cruise control
- Highway traffic control



# $\mathcal{R}$ Successful CPS Proofs



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All the way from

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## ℜ CPS are Multi-Dynamical Systems

#### **CPS** Dynamics

CPS are characterized by multiple facets of dynamical systems.



### **CPS** Compositions

CPS combine multiple simple dynamical effects.

#### Tame Parts

Exploiting compositionality tames complexity.

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### Challenge (Hybrid Systems)

Fixed rule describing state evolution with both

- Discrete dynamics (control decisions)
- Continuous dynamics (differential equations)





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# ${\mathscr R}$ CPS Analysis: Other Agents

### Challenge (Hybrid Games)

Game rules describing play evolution with

- Discrete dynamics (control decisions)
- Continuous dynamics (differential equations)
- Adversarial dynamics
   (Angel ◊ vs. Demon □),





# $\mathscr{R}$ CPS Analysis: Other Agents

### Challenge (Hybrid Games)

Game rules describing play evolution with

- Discrete dynamics (control decisions)
- Continuous dynamics (differential equations)
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# ℜ CPS are Multi-Dynamical Systems



HS = discrete + ODE



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# $\mathcal{R}$ Family of Differential Dynamic Logics



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# lpha Family of Differential Dynamic Logics



# ℛ Successful CPS Proofs



#### ICFEM'09,CAV'08,FM'09,HSCC'11

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# ℜ Successful CPS Proofs



#### FM'11,LMCS'12,ICCPS'12,ITSC'11,ITSC'13,IJCAR'12

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# ℜ Successful CPS Proofs



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# $\mathcal{R}$ Family of Differential Dynamic Logics



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# lpha Family of Differential Dynamic Logics



#### Definition (Hybrid program $\alpha$ )

$$\mathsf{x} := \theta \mid \mathsf{?}H \mid \mathsf{x}' = \mathsf{f}(\mathsf{x}) \, \& \, H \mid \alpha \cup \beta \mid \alpha; \beta \mid \alpha^*$$

#### Definition (d $\mathcal{L}$ Formula $\phi$ )

$$\theta_1 \ge \theta_2 \mid \neg \phi \mid \phi \land \psi \mid \forall x \phi \mid \exists x \phi \mid [\alpha] \phi \mid \langle \alpha \rangle \phi$$

# ℜ Differential Dynamic Logic dL: Syntax



### Definition (d $\mathcal{L}$ Formula $\phi$ )

$$\theta_1 \geq \theta_2 \mid \neg \phi \mid \phi \land \psi \mid \forall x \phi \mid \exists x \phi \mid [\alpha] \phi \mid \langle \alpha \rangle \phi$$



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# $\mathcal{R}$ Differential Dynamic Logic d $\mathcal{L}$ : Semantics

#### Definition (Hybrid program $\alpha$ )

$$\rho(x := \theta) = \{(v, w) : w = v \text{ except } \llbracket x \rrbracket_w = \llbracket \theta \rrbracket_v \}$$
  

$$\rho(?H) = \{(v, v) : v \models H \}$$
  

$$\rho(x' = f(x)) = \{(\varphi(0), \varphi(r)) : \varphi \models x' = f(x) \text{ for some duration } r \}$$
  

$$\rho(\alpha \cup \beta) = \rho(\alpha) \cup \rho(\beta)$$
  

$$\rho(\alpha; \beta) = \rho(\beta) \circ \rho(\alpha)$$
  

$$\rho(\alpha^*) = \bigcup_{n \in \mathbb{N}} \rho(\alpha^n)$$

#### Definition (d $\mathcal{L}$ Formula $\phi$ )

v	$\models \theta_1 \geq \theta_2$	iff	$\llbracket  heta_1  rbracket_{v} \geq \llbracket  heta_2  rbracket_{v}$
V	$\models [\alpha]\phi$	iff	$w \models \phi$ for all $w$ with $v \rho(\alpha) w$
v	$\models \langle \alpha \rangle \phi$	iff	$w \models \phi$ for some $w$ with $v \rho(\alpha) w$
v	$\models \forall x \phi$	iff	$w \models \phi$ for all $w$ that agree with $v$ except for $x$
v	$\models \exists x  \phi$	iff	$w \models \phi$ for some $w$ that agrees with $v$ except for $x$
v	$\models \phi \land \psi$	iff	$\mathbf{v} \models \phi$ and $\mathbf{v} \models \psi$
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# R Differential Dynamic Logic: Axiomatization

- $[:=] \quad [x := \theta] \phi(x) \leftrightarrow \phi(\theta)$ 
  - $[?] \quad [?H]\phi \leftrightarrow (H \to \phi)$
  - $['] \quad [x' = f(x)]\phi \leftrightarrow \forall t \ge 0 \, [x := y(t)]\phi$

 $\left(y'(t)=f(y)\right)$ 

- $[\cup] \quad [\alpha \cup \beta]\phi \leftrightarrow [\alpha]\phi \land [\beta]\phi$
- $[;] \quad [\alpha;\beta]\phi \leftrightarrow [\alpha][\beta]\phi$
- $[*] \quad [\alpha^*]\phi \leftrightarrow \phi \land [\alpha][\alpha^*]\phi$
- $\mathsf{K} \quad [\alpha](\phi \to \psi) \to ([\alpha]\phi \to [\alpha]\psi)$
- $\mathsf{I} \quad [\alpha^*](\phi \to [\alpha]\phi) \to (\phi \to [\alpha^*]\phi)$
- $\mathsf{C} \quad [\alpha^*] \forall \mathsf{v} > \mathsf{0} \left( \varphi(\mathsf{v}) \to \langle \alpha \rangle \varphi(\mathsf{v}-1) \right) \to \forall \mathsf{v} \left( \varphi(\mathsf{v}) \to \langle \alpha^* \rangle \exists \mathsf{v} \leq \mathsf{0} \varphi(\mathsf{v}) \right)$

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equations of truth

# R Differential Dynamic Logic: Axiomatization

equations of truth





# $\mathcal{R}$ Proofs for Hybrid Systems





# $\mathscr{R}$ Proofs for Hybrid Systems



w

w

# $\mathscr{R}$ Proofs for Hybrid Systems



 $\frac{\forall t \ge 0 \, [x := y_x(t)]\phi}{[x' = f(x)]\phi}$ 


#### compositional semantics $\Rightarrow$ compositional rules!









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### ℜ Complete Proof Theory of Hybrid Systems

### Theorem (Sound & Complete)

(J.Autom.Reas. 2008, LICS'12)

d*L* calculus is a sound & complete axiomatization of hybrid systems relative to either differential equations **or** discrete dynamics.

Corollary (Complete Proof-theoretical Alignment & Bridging)

proving continuous = proving hybrid = proving discrete

#### JAutomReas'08,LICS'12

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#### JAutomReas'08,LICS'12







Verification?

looks correct



## $\mathcal{R}$ Air Traffic Control



vermeation:

looks correct NO!

### $\mathcal{R}$ Air Traffic Control



 $x_{1}(t) = \frac{1}{\omega \varpi} (x_{1} \omega \varpi \cos t \omega - v_{2} \omega \cos t \omega \sin \vartheta + v_{2} \omega \cos t \omega \cos t \varpi \sin \vartheta - v_{1} \varpi \sin t \omega$  $+ x_{2} \omega \varpi \sin t \omega - v_{2} \omega \cos \vartheta \cos t \varpi \sin t \omega - v_{2} \omega \sqrt{1 - \sin \vartheta^{2}} \sin t \omega$ 

 $+ v_2\omega\cosartheta\cos t\omega\sin t\omega + v_2\omega\sinartheta\sin t\omega\sin t\omega)\dots$ 

### $\mathcal{R}$ Air Traffic Control



```
\forall B ts2
   ( 0 <= ts2 & ts2 <= t2 0
    -> ( (om 1)^-1
              * (omb 1)^-1
              * ( om 1 * omb 1 * x1 * Cos(om 1 * ts2)
                  + om 1 * v2 * \cos(\text{om 1 * ts2}) * (1 + -1 * (\cos(u))^2)^(1 / 2)
                 + -1 * omb_1 * v1 * Sin(om_1 * ts2)
                 + om 1 * omb 1 * x2 * Sin(om 1 * ts2)
                 + \text{ om } 1 * v2 * Cos(u) * Sin(om 1 * ts2)
                 + -1 * om_1 * v2 * Cos(omb_1 * ts2) * Cos(u) * Sin(om_1 * ts2)
                 + om 1 * v2 * Cos(om 1 * ts2) * Cos(u) * Sin(omb 1 * ts2)
                 + om_1 * v2 * Cos(om_1 * ts2) * Cos(omb_1 * ts2) * Sin(u)
                 + om 1 * v2 * Sin(om 1 * ts2) * Sin(omb 1 * ts2) * Sin(u)))
          ^2
        + ( (om 1)^-1
              * (omb 1)^-1
              *(-1 * omb 1 * v1 * Cos(om 1 * ts2))
                 + om 1 * omb 1 * x2 * Cos(om 1 * ts2)
                 + omb 1 * v1 * (Cos(om_1 * ts2))^2
                 + om_1 * v2 * Cos(om_1 * ts2) * Cos(u)
                 + -1 * \text{om } 1 * \text{v2} * \text{Cos(om } 1 * \text{ts2}) * \text{Cos(omb } 1 * \text{ts2}) * \text{Cos(u)}
                 + -1 * om 1 * omb 1 * x1 * Sin(om 1 * ts2)
                 + -1
                    * om 1
                    * v2
                    * (1 + -1 * (Cos(u))^2)^{(1 / 2)}
                    * Sin(om 1 * ts2)
                 + omb 1 * v1 * (Sin(om 1 * ts2))^2
                 + -1 * \text{om } 1 * v2 * \cos(u) * \sin(\text{om } 1 * \text{ts}2) * \sin(\text{om } 1 * \text{ts}2)
                 + -1 * \text{om } 1 * v2 * \text{Cos(omb } 1 * \text{ts2}) * \text{Sin(om } 1 * \text{ts2}) * \text{Sin(u)}
                 + om 1 * v2 * Cos(om 1 * ts2) * Sin(omb 1 * ts2) * Sin(u)))
         ^2
       >= (p)^2).
t2 0 >= 0.
x1^2 + x2^2 >= (p)^2
==>
```

```
\forall R t7.
  ( t7 >= 0
   ->
          ( (om_3)^-1
            * ( om 3
                 * ( (om 1)^-1
                    * (omb_1)^-1
                    * ( om_1 * omb_1 * x1 * Cos(om_1 * t2_0)
                      + om 1
                         * v2
                         * Cos(om_1 * t2_0)
                         * (1 + -1 * (Cos(u))^2)^{(1 / 2)}
                      + -1 * omb 1 * v1 * Sin(om 1 * t2 0)
                      + om_1 * omb_1 * x2 * Sin(om_1 * t2_0)
                      + om 1 * v2 * Cos(u) * Sin(om 1 * t2 0)
                      + -1
                         * om_1
                         * v2
                         * Cos(omb 1 * t2 0)
                        * Cos(11)
                         * Sin(om_1 * t2_0)
                      + om 1
                         * v2
                         * Cos(om_1 * t2_0)
                         * Cos(u)
                         * Sin(omb 1 * t2 0)
                      + om_1
                         * v2
                        * Cos(om_1 * t2_0)
                        * Cos(omb_1 * t2_0)
                         * Sin(u)
                      + om 1
                         * v2
                         * Sin(om_1 * t2_0)
                         * Sin(omb 1 * t2 0)
                         * Sin(u)))
```

```
* Cos(om_3 * t5)
+ v2
  * Cos(om 3 * t5)
  * (1
       + -1
         * (\cos(-1 * \text{ om } 1 * t2 0 + \text{ omb } 1 * t2 0 + u + \text{Pi} / 4))^2)
   ^(1 / 2)
+ -1 * v1 * Sin(om 3 * t5)
+ om 3
  * ( (om 1)^-1
     * (omb_1)^-1
     * ( -1 * omb_1 * v1 * Cos(om_1 * t2_0)
        + om_1 * omb_1 * x2 * Cos(om_1 * t2_0)
        + omb_1 * v1 * (Cos(om_1 * t2_0))^2
        + om 1 * v2 * Cos(om 1 * t2 0) * Cos(u)
        + -1
          * om_1
          * v2
         * Cos(om 1 * t2 0)
         * Cos(omb_1 * t2_0)
          * Cos(11)
        + -1 * om 1 * omb 1 * x1 * Sin(om 1 * t2 0)
        + -1
          * om_1
          * v2
          * (1 + -1 * (Cos(u))^2)^{(1 / 2)}
          * Sin(om_1 * t2_0)
        + omb 1 * v1 * (Sin(om 1 * t2 0))^2
        + -1
          * om_1
          * v2
          * Cos(u)
          * Sin(om_1 * t2_0)
          * Sin(omb_1 * t2_0)
```

```
+ -1
                  * om 1
                  * v2
                  * Cos(omb 1 * t2 0)
                  * Sin(om 1 * t2 0)
                  * Sin(u)
               + om_1
                  * v2
                  * Cos(om_1 * t2_0)
                  * Sin(omb_1 * t2_0)
                  * Sin(u)))
          * Sin(om 3 * t5)
       + v2
          * Cos(-1 * om 1 * t2 0 + omb 1 * t2 0 + u + Pi / 4)
          * Sin(om 3 * t5)
       + v2
          * (Cos(om_3 * t5))^2
          * Sin(-1 * om 1 * t2 0 + omb 1 * t2 0 + u + Pi / 4)
       + v2
          * (Sin(om_3 * t5))^2
          * Sin(-1 * om 1 * t2 0 + omb 1 * t2 0 + u + Pi / 4)))
 ^2
+ ( (om_3)^-1
     * ( -1 * v1 * Cos(om 3 * t5)
       + om 3
          * ( (om_1)^-1
            * (omb_1)^-1
             * ( -1 * omb 1 * v1 * Cos(om 1 * t2 0)
               + om_1 * omb_1 * x2 * Cos(om_1 * t2_0)
               + omb_1 * v1 * (Cos(om_1 * t2_0))^2
               + om 1 * v2 * Cos(om 1 * t2 0) * Cos(u)
               + -1
                  * om_1
                  * v2
                  * Cos(om 1 * t2 0)
                  * Cos(omb_1 * t2_0)
```

```
+ -1 * \text{om } 1 * \text{omb } 1 * x1 * \text{Sin(om } 1 * t2 0)
        + -1
          * om 1
          * v2
          * (1 + -1 * (Cos(u))^2)^(1 / 2)
          * Sin(om 1 * t2 0)
        + omb 1 * v1 * (Sin(om 1 * t2 0))^2
        + -1
          * om_1
          * v2
          * Cos(11)
          * Sin(om_1 * t2_0)
          * Sin(omb 1 * t2 0)
        + -1
          * om_1
          * v2
          * Cos(omb 1 * t2 0)
          * Sin(om_1 * t2_0)
          * Sin(u)
        + om 1
          * v2
          * Cos(om_1 * t2_0)
          * Sin(omb 1 * t2 0)
          * Sin(u)))
  * Cos(om_3 * t5)
+ v1 * (Cos(om 3 * t5))^2
+ v2
  * Cos(om_3 * t5)
  * Cos(-1 * om 1 * t2 0 + omb 1 * t2 0 + u + Pi / 4)
+ -1
  * v2
  * (Cos(om_3 * t5))^2
  * Cos(-1 * om 1 * t2 0 + omb 1 * t2 0 + u + Pi / 4)
```

```
+ -1
  * om 3
  * ( (om_1)^-1
     * (omb_1)^-1
     * ( om 1 * omb 1 * x1 * Cos(om 1 * t2 0)
        + om 1
          * v2
          * Cos(om 1 * t2 0)
          * (1 + -1 * (Cos(u))^2)^{(1 / 2)}
        + -1 * omb_1 * v1 * Sin(om_1 * t2_0)
        + om_1 * omb_1 * x2 * Sin(om_1 * t2_0)
        + om_1 * v2 * Cos(u) * Sin(om_1 * t2_0)
        + -1
          * om_1
          * v2
          * Cos(omb 1 * t2 0)
          * Cos(u)
          * Sin(om 1 * t2 0)
        + om 1
          * v2
          * Cos(om_1 * t2_0)
          * Cos(u)
          * Sin(omb_1 * t2_0)
        + om_1
          * v2
          * Cos(om 1 * t2 0)
          * Cos(omb_1 * t2_0)
          * Sin(u)
        + om 1
          * v2
          * Sin(om_1 * t2_0)
          * Sin(omb_1 * t2_0)
          * Sin(u)))
  * Sin(om_3 * t5)
```

```
+ -1
* v2
* (1
+ -1
* (Cos(-1 * om_1 * t2_0 + omb_1 * t2_0 + u + Pi / 4))^2)
^(1 / 2)
* Sin(om_3 * t5)
+ v1 * (Sin(om_3 * t5))^2
+ -1
* v2
* Cos(-1 * om_1 * t2_0 + omb_1 * t2_0 + u + Pi / 4)
* (Sin(om_3 * t5))^2))
2
>= (p)^2)
```

This is just one branch to prove for aircraft ....



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### Verification

Simple proof distinguishes safe from unsafe flight maneuver
# R Air Traffic Control



### Verification

Simple proof distinguishes safe from unsafe flight maneuver

Theorem (Differential radical invariant characterization)

$$\frac{h=0\rightarrow \bigwedge_{i=0}^{N-1}\mathcal{L}_p^{(i)}(h)=0}{h=0\rightarrow [x'=p]h=0}$$

characterizes all algebraic invariants, where  $N = \text{ord } \sqrt[\mathcal{L}]{(h)}$ , i.e.

$$\mathcal{L}_{p}^{(N)}(h) = \sum_{i=0}^{N-1} g_{i}\mathcal{L}_{p}^{(i)}(h)$$

### Corollary (Algebraic Invariants Decidable)

Invariance decidable for real algebraic h = 0.

TACAS'14

## ℜ Case Study: Longitudinal Dynamics of an Airplane

### 6th Order Longitudinal Equations

$$u' = \frac{X}{m} - g\sin(\theta) - qw \qquad u : \text{ axial velocity}$$

$$w' = \frac{Z}{m} + g\cos(\theta) + qu \qquad w : \text{ vertical velocity}$$

$$x' = \cos(\theta)u + \sin(\theta)w \qquad x : \text{ range}$$

$$z' = -\sin(\theta)u + \cos(\theta)w \qquad z : \text{ altitude}$$

$$\theta' = q \qquad \theta : \text{ pitch angle}$$

$$q' = \frac{M}{l_{yy}} \qquad q : \text{ pitch rate}$$

X : thrust along u, Z : thrust along w, M : thrust moment for wg : gravity, m : mass,  $I_{yy}$  : inertia second diagonal

# ℜ Case Study: Automatically Generated Invariants

### **Automatically Generated Invariant Functions**

$$\frac{Mz}{I_{yy}} + g\theta + \left(\frac{X}{m} - qw\right)\cos(\theta) + \left(\frac{Z}{m} + qu\right)\sin(\theta)$$
$$\frac{Mx}{I_{yy}} - \left(\frac{Z}{m} + qu\right)\cos(\theta) + \left(\frac{X}{m} - qw\right)\sin(\theta)$$
$$- q^2 + \frac{2M\theta}{I_{yy}}$$

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# ℜ Successful CPS Proofs



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# ℜ Successful CPS Proofs

## By Undergrads



students in 15-424/624 Foundations of Cyber-Physical Systems course

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ℜ What is Safe?



Robot and obstacle shape

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ℜ What is Safe?



Robot and obstacle shape

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 ✓ Verified with KeYmaera



 ✓ Verified with KeYmaera



KeYmaera



KeYmaera



 ✓ Verified with KeYmaera



 ✓ Verified with KeYmaera

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Safety	Invariant	+ Safe Control	(RSS'13)
static	$\ p_r - p_o\ _\infty > rac{v_r^2}{2b}$	$+\left(\frac{A}{b}+1\right)\left(\frac{A}{2}\varepsilon^{2}+\varepsilon\mathbf{v}_{r}\right)$	
passive	$v_r = 0 \vee \ p_r - p_o\ _{\infty} > \frac{v_r^2}{2b}$	$+Vrac{v_r}{b}+\Big(rac{A}{b}+1\Big)\Big(rac{A}{2}arepsilon^2+arepsilon($	$v_r + V) \Big)$
+ sensor	$\ \hat{p}_r-p_o\ _\infty>rac{v_r^2}{2b}+Vrac{v_r}{b}$	$+\left(\frac{A}{b}+1\right)\left(\frac{A}{2}\varepsilon^{2}+\varepsilon(v_{r}+V)\right)$	$\Big) + U_p$
+ disturb	$\ p_r - p_o\ _{\infty} > \frac{v_r^2}{2bU_m} + V \frac{v_r}{bU_m}$	$+\Big(rac{A}{bU_m}+1\Big)\Big(rac{A}{2}arepsilon^2+arepsilon(v_r+v_r)\Big)$	V))
+ failure	$\ \hat{p}_r-p_o\ _\infty>rac{v_r^2}{2b}+Vrac{v_r}{b}$	$+\left(\frac{A}{b}+1\right)\left(\frac{A}{2}\varepsilon^{2}+\varepsilon(v+V)\right)$	$\Big) + U_{ ho} + g\Delta$
friendly $\ p_r - p_o\ $	$\ _{\infty} > \frac{v_r^2}{2b} + \frac{V^2}{2b_o} + V\Big(\frac{v_r}{b} + \tau\Big)$	$+\Big(rac{A}{b}+1\Big)\Big(rac{A}{2}arepsilon^2+arepsilon(v_r+V)\Big)$	)

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**RSS'13** 

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Safety	Invariant	+ Safe Control	(RSS'13)
static	$\ p_r-p_o\ _\infty>rac{v_r^2}{2b}$	$+\left(\frac{A}{b}+1\right)\left(\frac{A}{2}\varepsilon^{2}-\frac{A}{b}\right)$	$+ \varepsilon v_r \Big)$
passive	$v_r = 0 \vee \ p_r - p_o\ _{\infty} > \frac{v_r^2}{2b}$	$+V\frac{v_r}{b}+\left(\frac{A}{b}+1\right)$	$\left(\frac{A}{2}\varepsilon^2+\varepsilon(v_r+V)\right)$
+ sensor	Question How to find and justify co	nstraints? Proof!	$+ \varepsilon (v_r + V) \Big) + U_p$
+ disturb    pr	$-p_o\ _{\infty} > \frac{v_r^2}{2bU_m} + V \frac{v_r}{bU_m}$	$+\left(\frac{A}{bU_m}+1\right)\left(\frac{A}{2}\right)$	$\varepsilon^2 + \varepsilon(v_r + V) \Big)$
+ failure	$\ \hat{p}_r - p_o\ _\infty > rac{v_r^2}{2b} + Vrac{v_r}{b}$	$+\Big(rac{A}{b}+1\Big)\Big(rac{A}{2}arepsilon^2+$	$+\varepsilon(v+V)\Big)+U_p+gZ$
friendly $\ p_r - p_o\ _{\infty}$	$>rac{v_r^2}{2b}+rac{V^2}{2b_o}+V\Big(rac{v_r}{b}+ au\Big)$	$+\left(\frac{A}{b}+1\right)\left(\frac{A}{2}\varepsilon^{2}-\frac{A}{b}\right)$	$+\varepsilon(v_r+V)\Big)$

**RSS'13** 

# $\mathcal{R}$ Outline

### 1 CPS are Multi-Dynamical Systems

- Hybrid Systems
- Hybrid Games

### 2 Dynamic Logic for Multi-Dynamical Systems

- Syntax
- Semantics

### 3 Proofs for CPS

## 4 Theory of CPS

- Soundness and Completeness
- Differential Invariants
- Differential Radical Invariants
- 5 Applications
  - Ground Robots

### Summary

## $oldsymbol{\mathcal{R}}$ Benefits of Logical Foundations for V & V

#### Proving

Safety Formalize system properties: What is "Safe"? "Reach goal"? Models Formalize system models

Assumptions Make assumptions explicit

Constraints Reveal invariants, switching conditions, starting conditions

Design Invariants guide safe controller design

Constructive Construct models along with their proof

#### Byproducts

Analyze Determine design trade-offs & feasibility early

Synthesize Turn high-level models into code & correctness monitors

Certify Proofs as artifacts for certification

#### Tools

#### KeYmaera Theorem prover for CPS

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