Pragmatic Approaches with COTS

Reverse Engineering for the Detection of Undesirable Functionality

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Wisdom Software, Ltd.

- Small Business- Incorporated in Virginia
- Business Interests
 - Formal Methods research and application development
 - Software Certification
 - System Safety Analysis
- Customer Base
 - NSA
 - US Army WRAIR
 - FDA

Research Objectives

- Evaluate the practicality of using reverse engineering (RE) to identify undesirable functionality (UF)
- Document and implement a practical (formal methods) process for RE
- Examine implications for addressing system security and safety

Undesirable Functionality

Definition

- permits non-secure access
- exports sensitive information
- damages computer or network operations
- modifies other unrelated systems or resources
- otherwise compromises system security

The UF Problem

- There is potentially an infinite set of unique instances of UF
- Each instance might be implemented in many ways
- Implications
 - A general solution is needed
 - Functionality must be examined in context

Finding Undesirable Functionality

- Pattern matching is insufficient
 - Only useful if you already know specifically what you are looking for
 - Does not address variations on implementation
 - Does not identify unknown undesirable functionality
- The alternative: Examining functional behavior in context

Examination of Behavior in Context

- Context is based on the scope of information abstracted from the software
 - Context of definition
 - Context of use
- Obvious behavior
 - Local function definition
 - Short stimulus sequence trigger
- Not so obvious behavior
 - Function definition is distributed
 - Long stimulus sequence trigger

Reverse Engineering Process

- Segment the software into functional blocks
- Abstract the program function for each segment
- Compose the abstracted segment functions
- Generate the Legal Sequence Table from the composed segment function

Segmentation

- Segmentation divides the software into functional blocks
- A segmentation graph shows the hierarchical relationship of the segments
- Each segment represents a context of definition

```
static void free_proc_chain(struct process_chain *procs)
    struct process_chain *p;
   int need_timeout = 0;
   int status;
   if (procs == ((void *)0) )
       return;
# 2666 "alloc.c"
   for (p = procs; p; p = p->next) {
        if (waitpid(p->pid, (int *) 0, 1 ) > 0) {
           p->kill_how = kill_never;
        }
    }
   for (p = procs; p; p = p->next) {
       if ((p->kill_how == kill_after_timeout)
            || (p->kill_how == kill_only_once)) {
           if (kill( p->pid , 15 ) != -1)
               need timeout = 1;
        }
        else if (p->kill_how == kill_always) {
           kill(p->pid, 9 );
        }
    }
   if (need_timeout)
       sleep(3);
   for (p = procs; p; p = p->next) {
        if (p->kill_how == kill_after_timeout)
           kill(p->pid, 9 );
        if (p->kill_how != kill_never)
           waitpid(p->pid, &status, 0);
    }
```

Function Abstraction

- Function abstraction is the process of deriving a state machine for a segment
- A state machine is abstracted for each segment in the system being analyzed
- The state machine is represented in a tabular format for each segment

Function Composition

- Abstracted state machines for related segments are combined by function composition
- A loop unrolling algorithm is employed to eliminate dynamic looping or recursion in the software
 - Resulting definition of behavior is an approximation
 - Accuracy increases as the number of cycles unrolled is increased

Segment 82			
Stimulus	Current Condition	State Update	Response
Invoke	procs == null		return terminate
Invoke	procs != null		waitpid(procs->pid, (int*) 0, 1)
waitpid(procs->pid, (int*) 0, 1) > 0	procs!= null	procs->kill_how = kill_never	terminate
waitpid(procs->pid, (int*) 0, 1) ≤ 0	procs!= null (procs->kill_how == kill_after_timeout OR procs->kill_how == kill_only_once)		kill(procs->pid, 15)
waitpid(procs->pid, (int*) 0, 1) ≤ 0	procs!= null procs->kill_how == kill_always		kill(procs->pid, 9) waitpid(procs->pid, &status, 0) terminate
waitpid(procs->pid, (int*) 0, 1) ≤ 0	procs!= null procs->kill_how != kill_after_timeout procs->kill_how != kill_never procs->kill_how != kill_only_once procs->kill_how != kill_always		waitpid(procs->pid, &status, 0) terminate
waitpid(procs->pid, (int*) 0, 1) ≤ 0	procs!= null procs->kill_how == kill_never		terminate
kill(procs->pid, 15) == -1	procs!= null procs->kill_how == kill_after_timeout		kill(procs->pid, 9) waitpid(procs->pid, &status, 0) terminate
kill(procs->pid, 15) == -1	procs!= null procs->kill_how == kill_only_once		waitpid(procs->pid, &status, 0) terminate
kill(procs->pid, 15) != -1	procs!= null procs->kill_how == kill_after_timeout		sleep(3) kill(procs->pid, 9) waitpid(procs->pid, &status, 0) terminate
kill(procs->pid, 15) != -1	procs!= null procs->kill_how == kill_only_once		sleep(3) waitpid(procs->pid, &status, 0) terminate

Legal Sequence Table

The composed segment table is processed to derive the legal stimulus sequences

- This identifies behavior in context of use
 - Series of responses generated for each sequence
 - Series of state updates generated for each sequence
 - For lower level segments possible conditions under which a sequence may occur
- Gives the analyst a clearer picture of the actions being performed by the segment

Segment 82 Legal sequences			
Sequence	Conditions	State Updates	Responses
Invoke	procs == null		return terminate
Invoke waitpid(procs->pid, (int*) 0, 1) > 0	procs != null	procs->kill_how = kill_never	waitpid(procs->pid, (int*) 0, 1) terminate
Invoke waitpid(procs->pid, (int*) 0, 1) ≤ 0 kill(procs->pid, 15) == -1	procs!= null procs->kill_how == kill_after_timeout		<pre>waitpid(procs->pid, (int*) 0, 1) kill(procs->pid, 15) kill(procs->pid, 9) waitpid(procs->pid, &status, 0) terminate</pre>
Invoke waitpid(procs->pid, (int*) 0, 1) ≤ 0 kill(procs->pid, 15) == -1	procs!= null procs->kill_how == kill_only_once		waitpid(procs->pid, (int*) 0, 1) kill(procs->pid, 15) waitpid(procs->pid, &status, 0) terminate
Invoke waitpid(procs->pid, (int*) 0, 1) ≤ 0 kill(procs->pid, 15) != -1	procs!= null procs->kill_how == kill_after_timeout		<pre>waitpid(procs->pid, (int*) 0, 1) kill(procs->pid, 15) sleep(3) kill(procs->pid, 9) waitpid(procs->pid, &status, 0) terminate</pre>
Invoke waitpid(procs->pid, (int*) 0, 1) ≤ 0 kill(procs->pid, 15) != -1	procs!= null procs->kill_how == kill_only_once		waitpid(procs->pid, (int*) 0, 1) kill(procs->pid, 15) sleep(3) waitpid(procs->pid, &status, 0) terminate
Invoke waitpid(procs->pid, (int*) 0, 1) ≤ 0	procs != null procs->kill_how == kill_always		waitpid(procs->pid, (int*) 0, 1 kill(procs->pid, 9) waitpid(procs->pid, &status, 0) terminate
Invoke waitpid(procs->pid, (int*) 0, 1) ≤ 0	procs!= null procs->kill_how != kill_after_timeout procs->kill_how != kill_never procs->kill_how != kill_only_once procs->kill_how != kill_always		waitpid(procs->pid, (int*) 0, 1) waitpid(procs->pid, &status, 0) terminate
Invoke waitpid(procs->pid, (int*) 0, 1) ≤ 0	procs!= null procs->kill_how == kill_never		waitpid(procs->pid, (int*) 0, 1) terminate

Results Analysis

- Behavioral description of the program function
 - State machine definition
 - Legal Sequence Table
- Potentially large volume of data to address
 - Prioritization current research
 - Assertion tests proposed research
 - Semantic analysis proposed research

Prioritization

- Prioritization is being investigated to direct analyst attention to segments that have a higher likelihood of containing UF
- Based on a cursory examination of attributes
 - Segment level
 - State machine level
 - Legal sequence level
- Generates a weighted score for each segment that reflects a potential for or sensitivity to UF

Ongoing Research

- Refinement of the prioritization weighting scheme
 - Optimal weighting values for measured attributes
 - Additional attributes to examine
- Refinement of Abstraction process
 - Improved automation support
- Results Analysis and Presentation