

# Zappa for Correctly Implementing CPSA Analyzed Protocols

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# Problem Statement

**Given** Security protocol specification that meets its goals

- Eg, verified using analysis tool

**Want** Protocol implementation meeting specification

- With assurance that it does

# Problem Statement

**Given** Security protocol specification that meets its goals

- Eg, verified using analysis tool
- CPSA, for instance

**Want** Protocol implementation meeting specification

- With assurance that it does
- Zappa compiles the spec
- Coq script validates source/target pairs for simplified core compiler

**Same** source file for both CPSA and Zappa

# How does Zappac work?

- CPSA uses the Dolev-Yao adversary model.
- Receptions are destructured using Dolev-Yao elimination rules.
  - ▶ Examples: projection from a tuple and decryption.
- Transmissions are synthesized using Dolev-Yao introduction rules.
  - ▶ Examples: Tuple construction and encryption.

# Corrected Blanchet Simple Example Protocol

$A \rightarrow B: \{\{\{S, B\}_{K_A^{-1}}\}_{K_B} \quad [S \text{ generated}]$   
 $B \rightarrow A: \{D\}_S \quad [D \text{ generated}]$

## CPSA Style Roles

Initiator (*init* role)

$init \rightarrow \{\{\{S, B\}_{K_A^{-1}}\}_{K_B}$



$\bullet \leftarrow \{D\}_S$

Responder (*resp* role)

$\{\{\{S, B\}_{K_A^{-1}}\}_{K_B} \rightarrow resp$



$\{D\}_S \leftarrow \bullet$

# CPSA Analysis From the Responder Point-of-View

Query

*resp*



---

Analysis

*resp*  $\leftarrow$   $\{\{S, B\}_{K_A^{-1}}\}_{K_B}$   $\succ$   $\leftarrow$   $\{\{S, B\}_{K_A^{-1}}\}_{K_B}$  *init*



# CPSA Summary

## Cryptographic Protocol Shapes Analyzer

- Query is a protocol and points-of-views.
- Analysis is concise descriptions of allowed behaviors.
- Deduce the goals achieved from the allowed behaviors.

# Modifications to CPSA to Support Code Generation

**Channels:** Specify an endpoint for message reception and transmission.

- Channels, as added to CPSA, have a distinct sort.
- Channels cannot be included in messages.

**Proc Info:** Specify generated procedure parameters and return values.

- All channels must be included in the parameters.

- resp:
- 1 Parameters  $C$ ,  $B$ ,  $K_A$ , and  $K_B^{-1}$ .
  - 2 Receive on  $C$ :  $\{\{\{S, B\}_{K_A^{-1}}\}_{K_B}\}$ .
  - 3 Generate  $D$ .
  - 4 Send on  $C$ :  $\{D\}_S$ .
  - 5 Return  $D$  and  $S$ .



# Zappac Output for the Responder Role

## Parameters

- 1 *Parameters  $C$ ,  $B$ ,  $K_A$ , and  $K_B^{-1}$ .*

```
/// Role: resp (blanchet.scm:16:3)
pub fn resp<M, C, Z: Zil<M, C>>(
  z: &Z, v0: &mut C, v1: &M, v2: &M, v3: &M
) -> Result<(M, M)> {
  z.chkck(b"name", v1)?;
  z.chkck(b"akey", v2)?;
  z.chkck(b"~akey", v3)?;
```

# Zappac Output for the Responder Role

## Reception

② *Receive on C:  $\{\{S, B\}_{K_A^{-1}}\}_{K_B}$ .*

```
// Recv (blanchet.scm:19:6)
let v4 = z.recv(b"enc", v0)?;
let v5 = z.adec(b"enc", &v4, v3)?;
let v6 = z.vrfy(b"enc", &v5, v2)?;
let v7 = z.proj(b"cat", 2, &v6, 0)?;
z.chck(b"skey", &v7)?;
let v8 = z.proj(b"cat", 2, &v6, 1)?;
z.chck(b"name", &v8)?;
z.same(&v8, v1)?;
```

# Zappac Output for the Responder Role

## Transmission and Return Values

- 3 *Generate D.*
- 4 *Send on C:  $\{D\}_S$ .*
- 5 *Return D and S.*

```
// Send (blanchet.scm:20:6)
let v9 = z.frsh(b"data")?;
let v10 = z.senc(b"enc", &v9, &v7)?;
z.send(v0, &v10)?;
Ok((v9, v7))
}
```

# Zappa Summary

- Source is a CPSA query.
- Target contains a procedure for each role in the protocol.
- Implement the Zil trait to create an executable.
  - ▶ Use the extensive Zappa libraries.
  - ▶ Provide your own implementation of the Zil trait.

# Why is Zappac Output Correct?

- Zappac implements the Roletran algorithm.
  - ▶ Roletran's source file is a simplification of Zappac's.
  - ▶ Roletran's target file is easily translated into Zappac's.
- A Coq scrip automatically proves Roletran source/target pairs have the same semantics, and therefore achieve the same goals.
  - ▶ Script proves:
    - ★ A run of the protocol, implies an execution for its procedure (Liveness).
    - ★ An execution of the procedure, implies a run of the protocol (Safety).
  - ▶ The procedure semantics is specified using a small step semantics.
  - ▶ Zil method implementations must honor the small step semantics.

# Zappa Runtime Systems

- Two distinct runtime system libraries are provided.
- Tools automatically generate parts of a runtime system.
- ASN.1 specified communication formats are supported.
- Zappa system supports protocols with state.

# Conclusion

## Assured Protocol Implementation Scheme

- Use CPSA to ensure your protocol achieves desired goals.
- Use Zappa to ensure the generated code achieves the same goals.
- CPSA and Roletran are available [here](#).  
Roletran paper is [here](#).