

Automatic Software Verification for High-Confidence Cyber-Physical Systems

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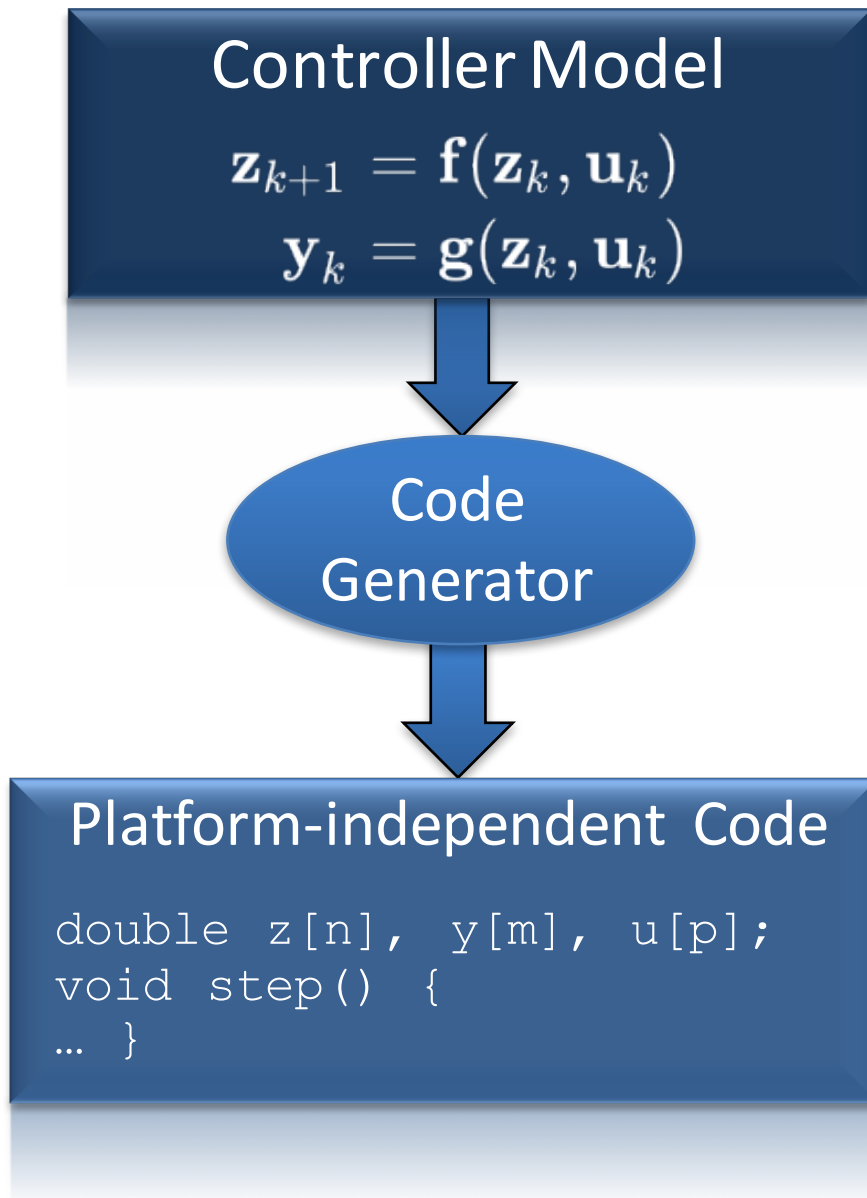
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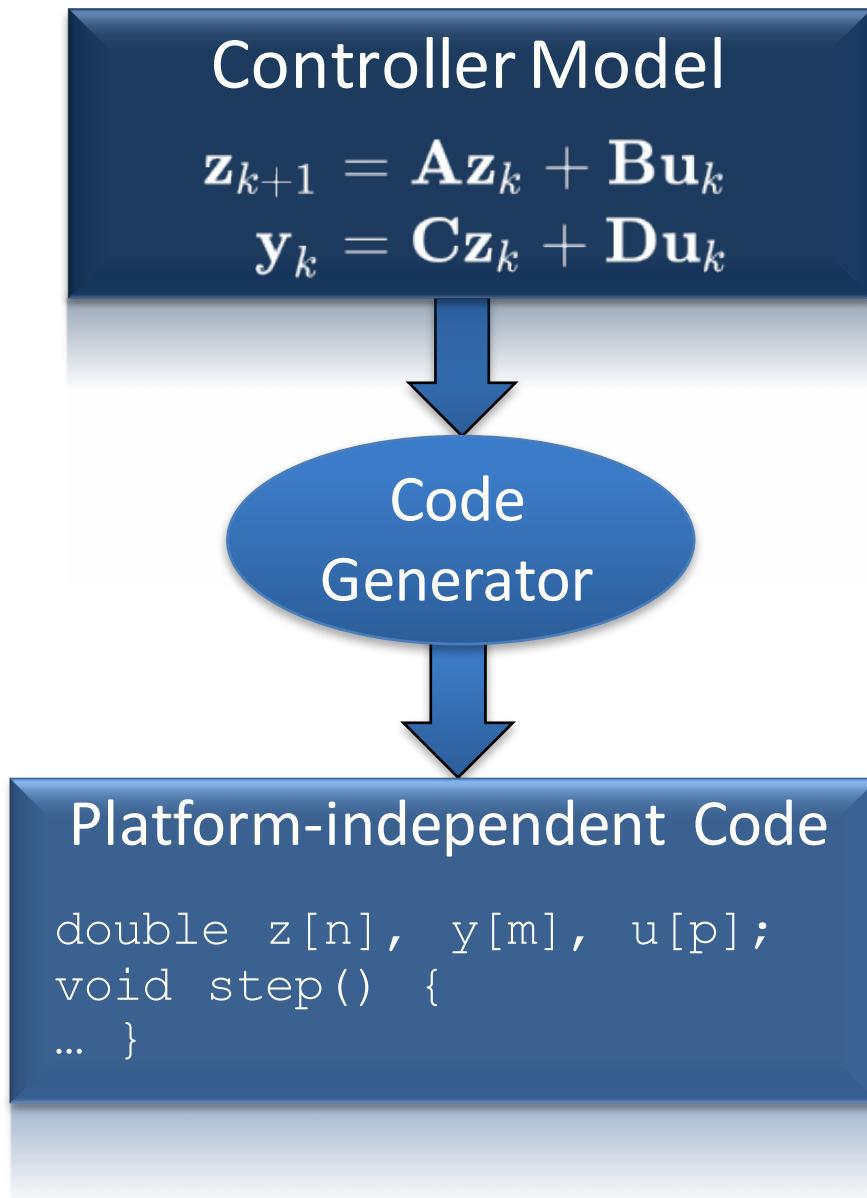
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Code Generation for Embedded Control



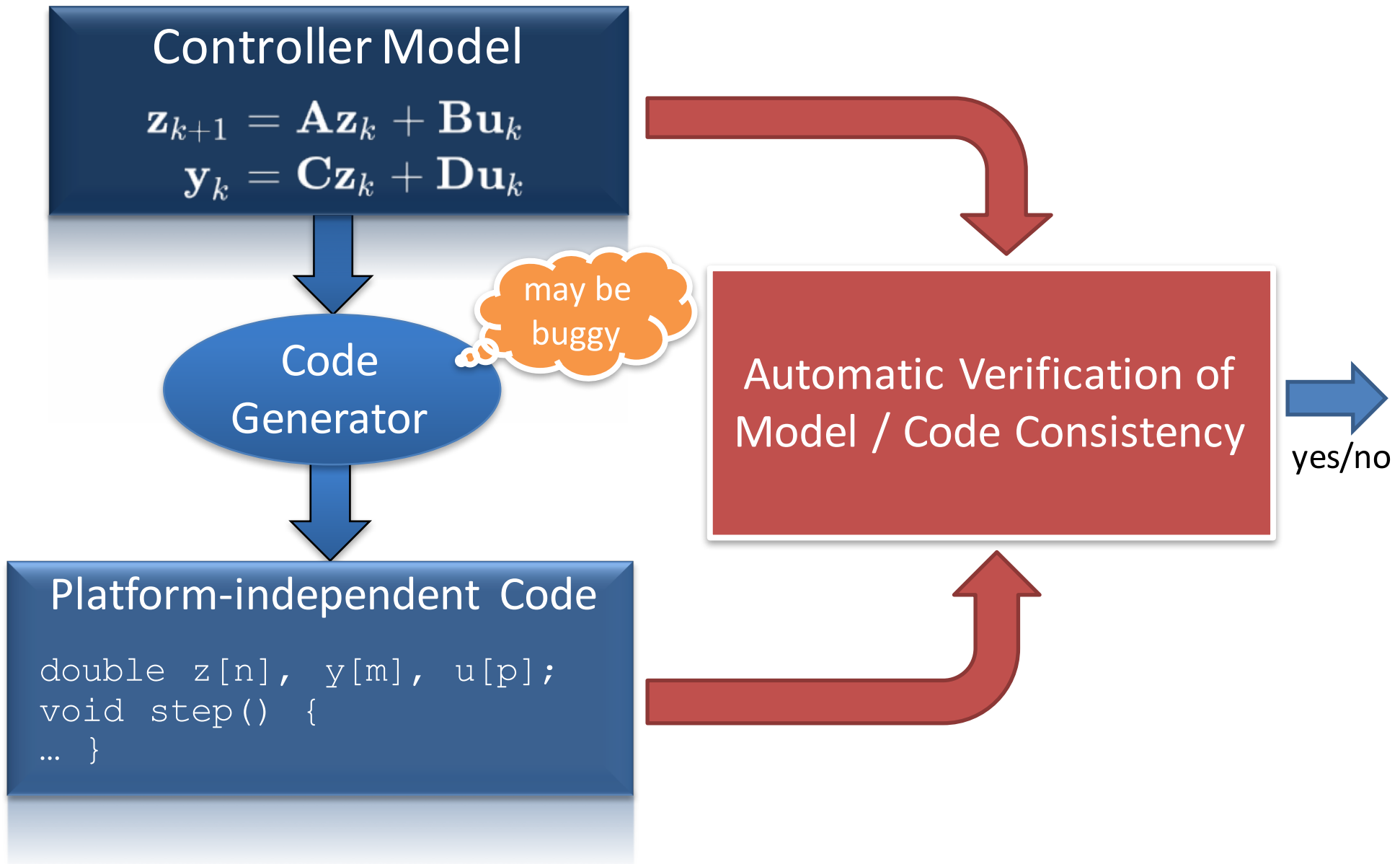
- From a closed-loop system model
 - Controller model $\Sigma(\mathbf{f}, \mathbf{g})$ (i.e., controller parameters)
- Code generator
 - *step* function
 - may employ optimization that affects the controller state

Code Generation for Embedded Control



- From a closed-loop system model
 - Controller model $\Sigma(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ (i.e., controller parameters)
- Code generator
 - *step* function
 - may employ optimization that affects the controller state
- Goal – **verification** of the generated code
 - Linear controllers - a very large class of embedded controllers

Code Generation for Embedded Control



A straightforward approach...

**Defining Invariants for Linear Controllers
based on input-output and state invariants**

Annotating Input-Output and State Invariants

- Exploit the ACSL's notion of the function contract
 - effectively a Hoare triple

Controller Model

$$\mathbf{z}_{k+1} = \mathbf{A}\mathbf{z}_k + \mathbf{B}\mathbf{u}_k$$
$$\mathbf{y}_k = \mathbf{C}\mathbf{z}_k + \mathbf{D}\mathbf{u}_k$$

- Running example:

$$\mathbf{A} = \begin{bmatrix} 0.8147 & 1.1534 \\ 2.6413 & 3.6411 \end{bmatrix},$$

$$\mathbf{B} = \begin{bmatrix} 3.1019 \\ 2.1432 \end{bmatrix},$$

$$\mathbf{C} = [1.7121 \quad 0.1351]$$

```
double x[2], u, v;
/*@ requires \valid(x+(0..1));
   @ ensures x[0] == 0.8147*\old(x[0]) +
   @       1.1534*\old(x[1]) + 3.1019*\old(u);
   @ ensures x[1] == 2.6413*\old(x[0]) +
   @       3.6411*\old(x[1]) + 2.1432*\old(u);
   @ ensures y == 1.7121*\old(x[0]) +
   @       0.1351*\old(x[1]) + 0*\old(u);
*/
void step() {
    double t1, t2;
    y = 1.7121*x[0] + 0.1351*x[1];
    t1 = 0.8147*x[0] + 1.1534*x[1] + 3.1019*u;
    t2 = 2.6413*x[0] + 3.6411*x[1] + 2.1432*u;
    x[0] = t1;
    x[1] = t2;
}
```

A straightforward approach...
... does not always work

A simple linear integrator

$$y_k = \sum_{i=0}^{k-1} \alpha u_i, \quad k \geq 1,$$

$$y_0 = 0$$

$$\Sigma(1, \alpha, 1, 0)$$

$$z_{k+1} = z_k + \alpha u_k,$$

$$y_k = z_k$$

$$\hat{\Sigma}(1, 1, \alpha, 0)$$

$$z_{k+1} = z_k + u_k,$$

$$y_k = \alpha z_k$$

- Both functionally correct but the maintained states are different
 - The latter could introduce a lower computational error when finite precision computations are taken into account

Example

MIMO control of a batch reactor

$$\mathbf{z}_{k+1} = \underbrace{\begin{bmatrix} 0.942 & 0.006888 & 0.04187 & -0.02319 \\ -0.01543 & 0.7965 & -0.03386 & 0.001563 \\ -0.1537 & 0.0137 & 0.7417 & 0.2006 \\ -0.03841 & 0.05637 & -0.02116 & 0.9949 \end{bmatrix}}_{\mathbf{A}} \mathbf{z}_k + \underbrace{\begin{bmatrix} 0.0774 & -0.0103 \\ -0.0022 & 0.0227 \\ 0.0267 & 0.0398 \\ 0.0356 & 0.0001 \end{bmatrix}}_{\mathbf{B}} \mathbf{y}_k$$

$$\mathbf{u}_k = \underbrace{\begin{bmatrix} 0.0583 & 0.9093 & 0.3258 & 0.08721 \\ -2.464 & -0.0504 & -1.71 & 1.165 \end{bmatrix}}_{\mathbf{C}} \mathbf{z}_k$$

$n(n+p)$ multiplications

$$\hat{\mathbf{z}}_{k+1} = \underbrace{\begin{bmatrix} 0.7636 & 0 & 0 & 0 \\ 0 & 0.8393 & 0 & 0 \\ 0 & 0 & 0.9595 & 0 \\ 0 & 0 & 0 & 0.9127 \end{bmatrix}}_{\hat{\mathbf{A}}} \hat{\mathbf{z}}_k + \underbrace{\begin{bmatrix} -0.2867 & -0.2581 \\ -0.3964 & -0.04506 \\ -0.07256 & 0.03278 \\ 0.5478 & -0.003331 \end{bmatrix}}_{\hat{\mathbf{B}}} \mathbf{y}_k,$$

$$\mathbf{u}_k = \underbrace{\begin{bmatrix} -0.1318 & 0.03834 & 0.02127 & -0.01226 \\ 0.147 & 0.08209 & -0.08674 & -0.2307 \end{bmatrix}}_{\hat{\mathbf{C}}} \hat{\mathbf{z}}_k$$

$n(1+p)$ multiplications

There exists a non-singular matrix \mathbf{T} : $\hat{\mathbf{A}} = \mathbf{T}^{-1} \mathbf{A} \mathbf{T}$, $\hat{\mathbf{B}} = \mathbf{T}^{-1} \mathbf{B}$, $\hat{\mathbf{C}} = \mathbf{C} \mathbf{T}$

If the same inputs then: $\forall k \geq 0, \hat{\mathbf{z}}_k = \mathbf{T}^{-1} \mathbf{z}_k \Leftrightarrow \hat{\mathbf{z}}_0 = \mathbf{T}^{-1} \mathbf{z}_0$

Example

MIMO control of a batch reactor

$$\mathbf{z}_{k+1} = \underbrace{\begin{bmatrix} 0.942 & 0.006888 & 0.04187 & -0.02319 \\ -0.01543 & 0.7965 & -0.03386 & 0.001563 \\ -0.1537 & 0.0137 & 0.7417 & 0.2006 \\ -0.03841 & 0.05637 & -0.02116 & 0.9949 \end{bmatrix}}_{\mathbf{A}} \mathbf{z}_k + \underbrace{\begin{bmatrix} 0.0774 & -0.0103 \\ -0.0022 & 0.0227 \\ 0.0267 & 0.0398 \\ 0.0356 & 0.0001 \end{bmatrix}}_{\mathbf{B}} \mathbf{y}_k$$
$$\mathbf{u}_k = \underbrace{\begin{bmatrix} 0.0583 & 0.9093 & 0.3258 & 0.08721 \\ -2.464 & -0.0504 & -1.71 & 1.165 \end{bmatrix}}_{\mathbf{C}} \mathbf{z}_k$$

$n(n+p)$ multiplications

$$\hat{\mathbf{z}}_{k+1} = \underbrace{\begin{bmatrix} 0.7636 & 0 & 0 & 0 \\ 0 & 0.8393 & 0 & 0 \\ 0 & 0 & 0.9595 & 0 \\ 0 & 0 & 0 & 0.9127 \end{bmatrix}}_{\hat{\mathbf{A}}} \hat{\mathbf{z}}_k + \underbrace{\begin{bmatrix} -0.2867 & -0.2581 \\ -0.3964 & -0.04506 \\ -0.07256 & 0.03278 \\ 0.5478 & -0.003331 \end{bmatrix}}_{\hat{\mathbf{B}}} \mathbf{y}_k,$$
$$\mathbf{u}_k = \underbrace{\begin{bmatrix} -0.1318 & 0.03834 & 0.02127 & -0.01226 \\ 0.147 & 0.08209 & -0.08674 & -0.2307 \end{bmatrix}}_{\hat{\mathbf{C}}} \hat{\mathbf{z}}_k$$

$n(1+p)$ multiplications

When same inputs are applied, the controllers' outputs will be identical!

- The controllers provide the same control functionality – **input-output conformance**

How to verify LTI controllers when the maintained state is now known?

We need a specification of the controller that is insensitive to the representation of control state

Invariant-Checking Approach (IC)

Original Model

$$z_{k+1} = \begin{bmatrix} -0.500311 & 0.16751 & 0.028029 & -0.395599 & -0.652079 \\ 0.850942 & 0.181639 & -0.29276 & 0.481277 & 0.638183 \\ -0.458583 & -0.002389 & -0.154281 & -0.578708 & -0.769495 \\ 1.01855 & 0.638926 & -0.668256 & -0.258506 & 0.119959 \\ 0.100383 & -0.432501 & 0.122727 & 0.82634 & 0.892296 \end{bmatrix} z_k + \begin{bmatrix} 1.1149 & 0.164423 \\ -1.56592 & 0.634384 \\ 1.04856 & -0.196914 \\ 1.96066 & 3.11571 \\ -3.02046 & -1.96087 \end{bmatrix} u_k$$

$$y_k = \begin{bmatrix} 0.283441 & 0.032612 & -0.75658 & 0.085468 & 0.161088 \\ -0.528786 & 0.050734 & -0.681773 & -0.432334 & -1.17988 \end{bmatrix} z_k$$

?

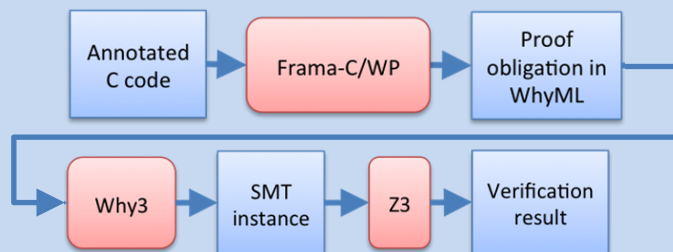
Code

```
void LTIS_step(void)
{
  {
    static const int_T colCidxRow0[5] = { 0, 1, 2, 3, 4 };
    const int_T *pCidx = &colCidxRow0[0];
    const real_T *pC0 = LTIS_ConstP.Internal_C;
    const real_T *xd = &LTIS_DW.Internal_DSTATE[0];
    real_T *y0 = &LTIS_Y.y[0];
    int_T numNonZero = 4;
    *y0 = (*pC0++) * xd[*pCidx++];
    while (numNonZero--) {
      *y0 += (*pC0++) * xd[*pCidx++];
    }
  }
  {
    static const int_T colCidxRow1[5] = { 0, 1, 2, 3, 4 };
    const int_T *pCidx = &colCidxRow1[0];
    const real_T *pC5 = &LTIS_ConstP.Internal_C[5];
    const real_T *xd = &LTIS_DW.Internal_DSTATE[0];
    real_T *y1 = &LTIS_Y.y[1];
    int_T numNonZero = 4;
    *y1 = (*pC5++) * xd[*pCidx++];
    while (numNonZero--) {
      *y1 += (*pC5++) * xd[*pCidx++];
    }
  }
  real_T xnew[5];
  int_T i;
  xnew[0] = (0.87224)*LTIS_DW.Internal_DSTATE[0];
  xnew[0] += (0.822174)*LTIS_U.u[0]+(-0.438008)*LTIS_U.u[1];
  xnew[1] = (0.366378)*LTIS_DW.Internal_DSTATE[1];
  xnew[1] += (-0.278536)*LTIS_U.u[0]+(-0.824313)*LTIS_U.u[1];
  xnew[2] = (-0.540795)*LTIS_DW.Internal_DSTATE[2];
  xnew[2] += (0.874484)*LTIS_U.u[0]+(0.858857)*LTIS_U.u[1];
  xnew[3] = (-0.332664)*LTIS_DW.Internal_DSTATE[3];
  xnew[3] += (-0.117628)*LTIS_U.u[0]+(-0.506362)*LTIS_U.u[1];
  xnew[4] = (-0.204322)*LTIS_DW.Internal_DSTATE[4];
  xnew[4] += (-0.955459)*LTIS_U.u[0]+(-0.622498)*LTIS_U.u[1];
  for(i=0; i<5; i++) LTIS_DW.Internal_DSTATE[i] = xnew[i];
}
```

$$y_k = \sum_{i=0}^n \beta_i u_{k-i} - \sum_{i=1}^n \alpha_i y_{k-i} \quad (2) \text{ Automatic Annotation}$$

@ assert \at (y, k_n) + \alpha_1\at (y, k_{n-1}) + ...
@ \alpha_n*\at (y, k_0) == \beta_0*\at (u, k_n) + ...
@ \beta_n*\at (u, k_0)

Frama/C-based Toolchain



(1) Loop unrolling



A More Scalable Approach (SC)

Original Model

$$z_{k+1} = \begin{bmatrix} -0.500311 & 0.16751 & 0.028029 & -0.395599 & -0.652079 \\ 0.850942 & 0.181639 & -0.29276 & 0.481277 & 0.638183 \\ -0.458583 & -0.002389 & -0.154281 & -0.578708 & -0.769495 \\ 1.01855 & 0.638926 & -0.668256 & -0.258506 & 0.119959 \\ 0.100383 & -0.432501 & 0.122727 & 0.82634 & 0.892296 \end{bmatrix} z_k + \begin{bmatrix} 1.1149 & 0.164423 \\ -1.56592 & 0.634384 \\ 1.04856 & -0.196914 \\ 1.96066 & 3.11571 \\ -3.02046 & -1.96087 \end{bmatrix} u_k$$

$$y_k = \begin{bmatrix} 0.283441 & 0.032612 & -0.75658 & 0.085468 & 0.161088 \\ -0.528786 & 0.050734 & -0.681773 & -0.432334 & -1.17988 \end{bmatrix} z_k$$

?

Code

```
void LTIS_step(void)
{
  {
    static const int_T colCidxRow0[5] = { 0, 1, 2, 3, 4 };
    const int_T *pCidx = &colCidxRow0[0];
    const real_T *pC0 = LTIS_ConstP.Internal_C;
    const real_T *xd = &LTIS_DW.Internal_DSTATE[0];
    real_T *y0 = &LTIS_Y.y[0];
    int_T numNonZero = 4;
    *y0 = (*pC0++) * xd[*pCidx++];
    while (numNonZero--) {
      *y0 += (*pC0++) * xd[*pCidx++];
    }
  }
  {
    static const int_T colCidxRow1[5] = { 0, 1, 2, 3, 4 };
    const int_T *pCidx = &colCidxRow1[0];
    const real_T *pC5 = &LTIS_ConstP.Internal_C[5];
    const real_T *xd = &LTIS_DW.Internal_DSTATE[0];
    real_T *y1 = &LTIS_Y.y[1];
    int_T numNonZero = 4;
    *y1 = (*pC5++) * xd[*pCidx++];
    while (numNonZero--) {
      *y1 += (*pC5++) * xd[*pCidx++];
    }
  }
}

real_T xnew[5];
int_T i;
xnew[0] = (0.87224)*LTIS_DW.Internal_DSTATE[0];
xnew[0] += (0.822174)*LTIS_U.u[0]+(-0.438008)*LTIS_U.u[1];
xnew[1] = (0.366378)*LTIS_DW.Internal_DSTATE[1];
xnew[1] += (-0.278536)*LTIS_U.u[0]+(-0.824313)*LTIS_U.u[1];
xnew[2] = (-0.540795)*LTIS_DW.Internal_DSTATE[2];
xnew[2] += (0.874484)*LTIS_U.u[0]+(0.858857)*LTIS_U.u[1];
xnew[3] = (-0.332664)*LTIS_DW.Internal_DSTATE[3];
xnew[3] += (-0.117628)*LTIS_U.u[0]+(-0.506362)*LTIS_U.u[1];
xnew[4] = (-0.204322)*LTIS_DW.Internal_DSTATE[4];
xnew[4] += (-0.955459)*LTIS_U.u[0]+(-0.622498)*LTIS_U.u[1];
for(i=0; i<5; i++) LTIS_DW.Internal_DSTATE[i] = xnew[i];
}
```

(2) Input-Output Similarity
Checking

Extracted Model

$$\hat{z}_{k+1} = \begin{bmatrix} 0.87224 & 0 & 0 & 0 & 0 \\ 0 & 0.366378 & 0 & 0 & 0 \\ 0 & 0 & -0.540795 & 0 & 0 \\ 0 & 0 & 0 & -0.332664 & 0 \\ 0 & 0 & 0 & 0 & -0.204322 \end{bmatrix} \hat{z}_k + \begin{bmatrix} 0.822174 & -0.438008 \\ -0.278536 & -0.824313 \\ 0.874484 & 0.858857 \\ -0.117628 & -0.506362 \\ -0.955459 & -0.622498 \end{bmatrix} u_k$$

$$y_k = \begin{bmatrix} -0.793176 & 0.154365 & -0.377883 & -0.360608 & -0.142123 \\ 0.503767 & -0.573538 & 0.170245 & -0.583312 & -0.56603 \end{bmatrix} \hat{z}_k$$

(1) Model
Extraction

Defining Invariants for Linear Controllers

- Annotating input-output and state invariants
- Annotating input-output only invariants
- Inexact controller implementations
- Instantiation-based input-output invariants

Problem: How to check input-output conformance when state conformance is violated?

- Input-output invariants obtained from controllers *transfer functions*

$$\Sigma = (\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$$

$$\mathbf{G}(z) = \mathbf{C}(z\mathbf{I}_n - \mathbf{A})^{-1}\mathbf{B} + \mathbf{D}$$

- In the general case for Single-Input-Single-Output controllers

$$G(z) = \frac{\beta_0 + \beta_1 z^{-1} + \dots + \beta_n z^{-n}}{1 + \alpha_1 z^{-1} + \dots + \alpha_n z^{-n}},$$

and the controllers inputs and outputs satisfy

$$y_k = \sum_{i=0}^n \beta_i u_{k-i} - \sum_{i=1}^n \alpha_i y_{k-i}$$

with $y_k = 0, k < 0$ because $\mathbf{z}_0 = 0$ and $u_k = 0, k < 0$

Annotating Input-Output Only Invariants

- Cannot be specified using pre- and post-conditions for every execution of the `step` function
 - relates the last $n+1$ executions of the step function
- Perform execution **unrolling** of the `step` function
 - construct the `verif_driver` function invoking the `step` function exactly $n+1$ times

$$y_k = \sum_{i=0}^n \beta_i u_{k-i} - \sum_{i=1}^n \alpha_i y_{k-i}$$

```
\*@ assert \at (y, k_n) + \alpha_1 * \at (y, k_{n-1}) + ...  
  @ \alpha_n * \at (y, k_0) == \beta_0 * \at (u, k_n) + ...  
  @ \beta_n * \at (u, k_0)
```


Annotating Input-Output Only Invariants

- Cannot be specified using pre- and post-conditions for every execution of the `step` function
 - relates the last $n+1$ executions of the step function
- Perform execution unrolling of the `step` function
 - construct the `verif_driver` function invoking the `step` function exactly $n+1$ times

$$y_k = \sum_{i=0}^n \beta_i u_{k-i} - \sum_{i=1}^n \alpha_i y_{k-i}$$

```
\*@ assert \at(y, k_n) + \alpha_1*\at(y, k_{n-1}) + ..
  @ \alpha_n*\at(y, k_0) == \beta_0*\at(u, k_n) + ...
  @ \beta_n*\at(u, k_0)
```

$$\mathbf{A} = \begin{bmatrix} 0.8147 & 1.1534 \\ 2.6413 & 3.6411 \end{bmatrix}, \mathbf{B} = \begin{bmatrix} 3.1019 \\ 2.1432 \end{bmatrix}, \mathbf{C} = [1.7121 \quad 0.1351]$$

$$G(z) = \frac{5.60030931z^{-1} - 14.233777166248z^{-2}}{1 - 4.4558z^{-1} - 0.08007125z^{-2}}$$

```
extern double input();
```

```
void verif_driver() {
```

```
  u = input();  step();
  k0::
```

```
  u = input();  step();
  k1::
```

```
  u = input();  step();
  k2::
```

```
  /* @assert \at(y, k2) - 4.4558*\at(y, k1)
    @ - 0.08007125*\at(y, k0)
    @ == 5.60030931*\at(u, k1)
    @ - 14.233777166248*\at(u, k0);
    @ */
```

```
}
```

Errors from optimization in code generation

- Back to the running example
 - a more efficient controller obtained in Matlab using the function canon for the modal type of decomposition

$$\mathbf{A} = \begin{bmatrix} 0.8147 & 1.1534 \\ 2.6413 & 3.6411 \end{bmatrix}, \mathbf{B} = \begin{bmatrix} 3.1019 \\ 2.1432 \end{bmatrix}, \mathbf{C} = [1.7121 \quad 0.1351]$$
$$G(z) = \frac{5.60030931z^{-1} - 14.233777166248z^{-2}}{1 - 4.4558z^{-1} - 0.08007125z^{-2}}$$

$$\hat{\mathbf{A}} = \begin{bmatrix} -0.0179 & 0 \\ 0 & 4.474 \end{bmatrix}, \hat{\mathbf{B}} = \begin{bmatrix} -1.051 \\ -1.055 \end{bmatrix}, \hat{\mathbf{C}} = [-3.037 \quad -2.283]$$
$$\hat{G}(z) = \frac{5.600452z^{-1} - 14.2373891245z^{-2}}{1 - 4.4561z^{-1} - 0.0800846z^{-2}}$$

Defining Invariants for Linear Controllers

- Annotating input-output and state invariants
- Annotating input-output only invariants
- Inexact controller implementations
- Instantiation-based input-output invariants

Inexact Controller Implementations

- There is a need to extend our input-output invariants for the case with imprecise specification of the transfer functions
- Start by assuming that the transfer function could take the form

$$G(z) = \frac{\hat{\beta}_0 + \hat{\beta}_1 z^{-1} + \dots + \hat{\beta}_n z^{-n}}{1 + \hat{\alpha}_1 z^{-1} + \dots + \hat{\alpha}_n z^{-n}},$$

such that

$$i = 0, 1, \dots, n$$

$$\beta_i - \epsilon_\beta \leq \hat{\beta}_i \leq \beta_i + \epsilon_\beta, \quad \alpha_i - \epsilon_\alpha \leq \hat{\alpha}_i \leq \alpha_i + \epsilon_\alpha.$$

- 'Inexact' invariant

$$\exists \Delta\beta_i, \Delta\alpha_i \in \mathbb{R}, i = 0, \dots, n, \quad |\Delta\beta_i| \leq \epsilon_\beta \wedge |\Delta\alpha_i| \leq \epsilon_\alpha \wedge$$

$$y_k = \sum_{i=0}^n (\beta_i + \Delta\beta_i) u_{k-i} - \sum_{i=1}^n (\alpha_i + \Delta\alpha_i) y_{k-i}$$

nonlinear

Inexact Controller Implementations

- There is a need to extend our input-output invariants for the case with imprecise specification of the transfer functions
- Start by assuming that the transfer function could take the form

$$G(z) = \frac{\hat{\beta}_0 + \hat{\beta}_1 z^{-1} + \dots + \hat{\beta}_n z^{-n}}{1 + \hat{\alpha}_1 z^{-1} + \dots + \hat{\alpha}_n z^{-n}},$$

such that

$$i = 0, 1, \dots, n$$

$$\beta_i - \epsilon_\beta \leq \hat{\beta}_i \leq \beta_i + \epsilon_\beta, \quad \alpha_i - \epsilon_\alpha \leq \hat{\alpha}_i \leq \alpha_i + \epsilon_\alpha.$$

- 'Inexact' **linear** invariant

$$\exists \tilde{u}_{k-i}, \tilde{y}_{k-i} \in \mathbb{R}, i = 0, 1, \dots, n,$$

$$|\tilde{u}_{k-i}| \leq \epsilon_\beta |u_{k-i}| \wedge |\tilde{y}_{k-i}| \leq \epsilon_\alpha |y_{k-i}| \sim \wedge$$

$$\tilde{u}_{k-i} = \Delta\beta_i u_{k-i},$$

$$\tilde{y}_{k-i} = \Delta\alpha_i y_{k-i}$$

$$y_k = \sum_{i=0}^n (\beta_i u_{k-i} + \tilde{u}_{k-i}) - \sum_{i=1}^n (\alpha_i y_{k-i} + \tilde{y}_{k-i})$$

Inexact Controller Implementations

- Start by assuming that the transfer function could take the form

$$G(z) = \frac{\hat{\beta}_0 + \hat{\beta}_1 z^{-1} + \dots + \hat{\beta}_n z^{-n}}{1 + \hat{\alpha}_1 z^{-1} + \dots + \hat{\alpha}_n z^{-n}},$$

such that $i = 0, 1, \dots, n$

$$\beta_i - \epsilon_\beta \leq \hat{\beta}_i \leq \beta_i + \epsilon_\beta, \quad \alpha_i - \epsilon_\alpha \leq \hat{\alpha}_i \leq \alpha_i + \epsilon_\alpha.$$

- 'Inexact' **linear** invariant – *for all* u_{k-i} y_{k-i}

$$\exists \tilde{u}_{k-i}, \tilde{y}_{k-i} \in \mathbb{R}, i = 0, 1, \dots, n,$$

$$|\tilde{u}_{k-i}| \leq \epsilon_\beta |u_{k-i}| \wedge |\tilde{y}_{k-i}| \leq \epsilon_\alpha |y_{k-i}| \sim \wedge$$

$$\tilde{u}_{k-i} = \Delta\beta_i u_{k-i},$$

$$\tilde{y}_{k-i} = \Delta\alpha_i y_{k-i}$$

$$y_k = \sum_{i=0}^n (\beta_i u_{k-i} + \tilde{u}_{k-i}) - \sum_{i=1}^n (\alpha_i y_{k-i} + \tilde{y}_{k-i})$$

A mixture of both universal and existential quantifiers

Defining Invariants for Linear Controllers

- Annotating input-output and state invariants
- Annotating input-output only invariants
- Inexact controller implementations
- Instantiation-based input-output invariants

Instantiation-based Input-Output Invariants

$$\mathbf{D}_N = \begin{bmatrix} y_n & y_{n-1} & \dots & y_1 & y_0 & u_n & u_{n-1} & \dots & u_1 & u_0 \\ y_{n+1} & y_n & \dots & y_2 & y_1 & u_{n+1} & u_n & \dots & u_2 & u_1 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ y_{n+N-1} & y_{n+N-1} & \dots & y_N & y_{N-1} & u_{n+N-1} & u_{n+N-2} & \dots & u_N & u_{N-1} \end{bmatrix}$$

$$\mathbf{D}_N \cdot \theta = \mathbf{0}$$

$$\theta = [1 \quad \alpha_1 \quad \dots \quad \alpha_n \quad \beta_0 \quad \beta_1 \quad \dots \quad \beta_n]^T$$

$$y_k = \sum_{i=0}^n \beta_i u_{k-i} - \sum_{i=1}^n \alpha_i y_{k-i}$$

PROPOSITION 1. Consider LTI controller Σ of size n . Then the rank of any matrix \mathbf{D}_N cannot be larger than $2n + 1$. Furthermore, when the rank of \mathbf{D}_N is $2n + 1$, then linear conditions $\mathbf{D}_N \cdot \theta = \mathbf{0}$ are satisfied if and only if the condition (8) is satisfied for all k .

$$N = 2n + 1$$

Instantiation-based Input-Output Invariants

$$\mathbf{D}_N = \begin{bmatrix} y_n & y_{n-1} & \dots & y_1 & y_0 & u_n & u_{n-1} & \dots & u_1 & u_0 \\ y_{n+1} & y_n & \dots & y_2 & y_1 & u_{n+1} & u_n & \dots & u_2 & u_1 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ y_{n+N-1} & y_{n+N-1} & \dots & y_N & y_{N-1} & u_{n+N-1} & u_{n+N-2} & \dots & u_N & u_{N-1} \end{bmatrix}$$

$$\mathbf{D}_N \cdot \theta = \mathbf{0}$$

$$\theta = [1 \quad \alpha_1 \quad \dots \quad \alpha_n \quad \beta_0 \quad \beta_1 \quad \dots \quad \beta_n]^T$$

$$y_k = \sum_{i=0}^n \beta_i u_{k-i} - \sum_{i=1}^n \alpha_i y_{k-i}$$

Code annotations

```
\*@ assert ((\at(y, k_0) == y_0) && ... && (\at(y, k_{n-1}) == y_{n-1}) && (\at(u, k_0) == u_0) && ... && (\at(u, k_{3n}) == u_{3n}))  
@ => ((\at(y, k_n) == y_n) && ... && (\at(y, k_{3n}) == y_{3n}))
```

Allows us to specify a set of $2n + 1$ linear invariants

Instantiation-based Input-Output Invariants

$$\mathbf{D}_N = \begin{bmatrix} y_n & y_{n-1} & \dots & y_1 & y_0 & u_n & u_{n-1} & \dots & u_1 & u_0 \\ y_{n+1} & y_n & \dots & y_2 & y_1 & u_{n+1} & u_n & \dots & u_2 & u_1 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ y_{n+N-1} & y_{n+N-1} & \dots & y_N & y_{N-1} & u_{n+N-1} & u_{n+N-2} & \dots & u_N & u_{N-1} \end{bmatrix}$$

$$\mathbf{D}_N \cdot \theta = \mathbf{0}$$

$$\theta = [1 \quad \alpha_1 \quad \dots \quad \alpha_n \quad \beta_0 \quad \beta_1 \quad \dots \quad \beta_n]^T$$

$$y_k = \sum_{i=0}^n \beta_i u_{k-i} - \sum_{i=1}^n \alpha_i y_{k-i}$$

Code annotations for inexact controllers

$$\exists \Delta\beta_i, \Delta\alpha_i \in \mathbb{R}, i = 0, \dots, n, |\Delta\beta_i| \leq \epsilon_\beta \wedge |\Delta\alpha_i| \leq \epsilon_\alpha \wedge$$

$$\mathbf{D}_{2n+1}^y \Delta\alpha + \mathbf{D}_{2n+1}^u \Delta\beta = \mathbf{v} \wedge$$

$$y_{n+i} = \mathbf{D}_{2n+1}^y(n+i), \quad i = 0, \dots, 2n,$$

```
\*@ assert \exists real a_0, ..., a_{n-1}, b_0, ..., b_n
```

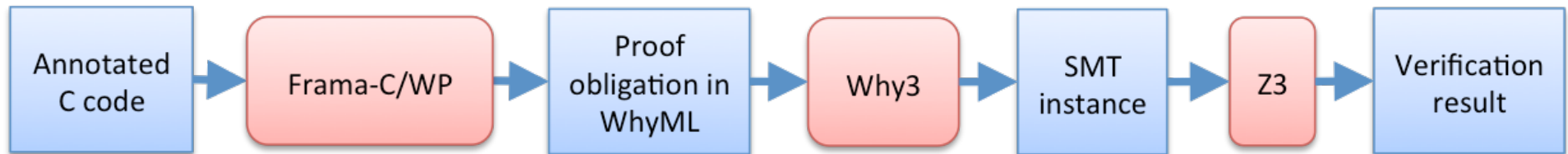
```
  @ (a_0 ≤ ε_α) && (a_0 ≥ -ε_α) && ... && (a_{n-1} ≤ ε_α) && (a_{n-1} ≥ -ε_α) && (b_0 ≤ ε_β) && (b_0 ≥ -ε_β) && ... && (b_n ≤ ε_β) && (b_n ≥ -ε_β)
```

```
\*@ ((\at(y, k_0) == y_0) && ... && (\at(y, k_{n-1}) == y_{n-1}) && (\at(u, k_0) == u_0) && ... && (\at(u, k_{3n}) == u_{3n}))
```

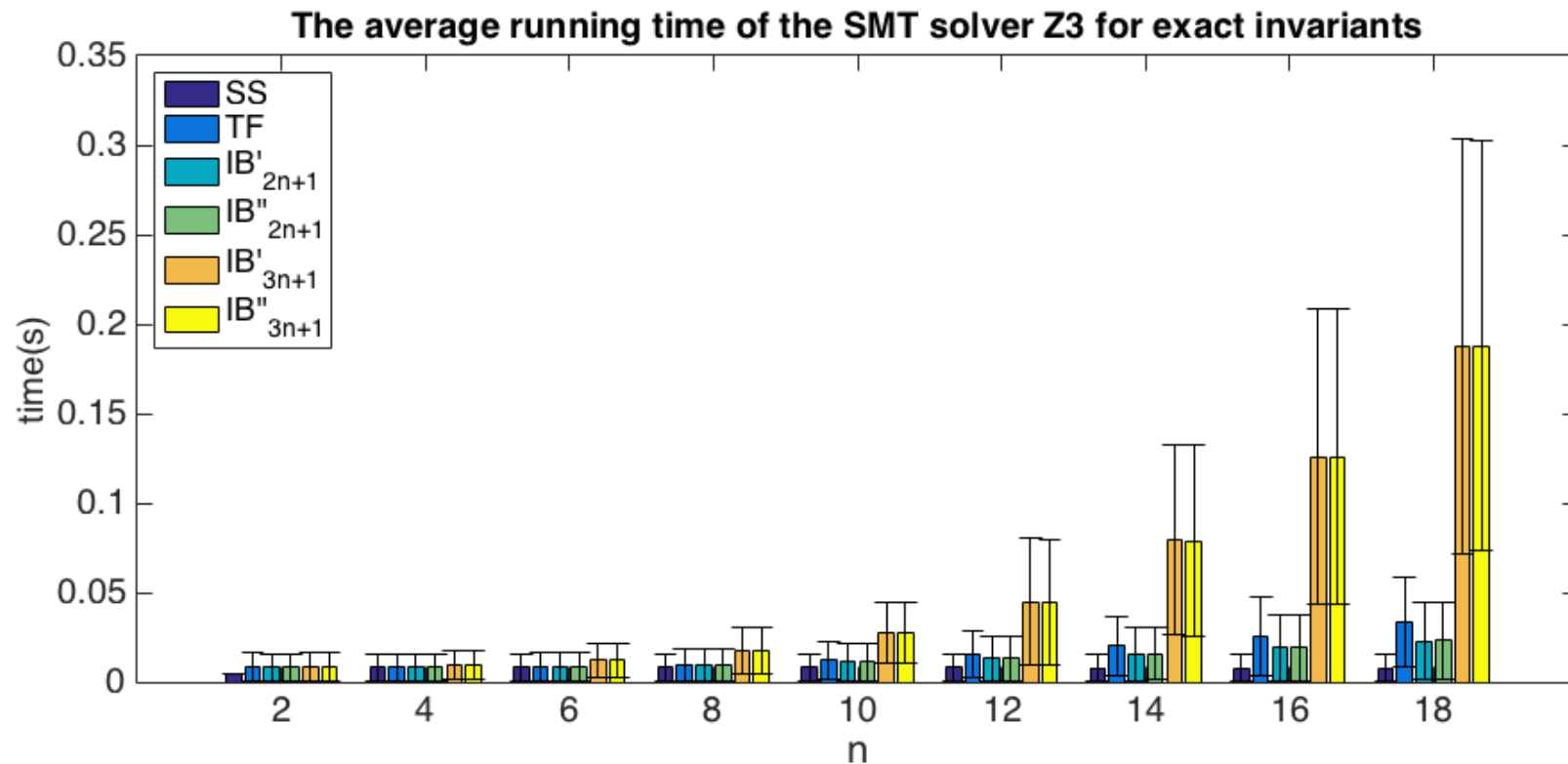
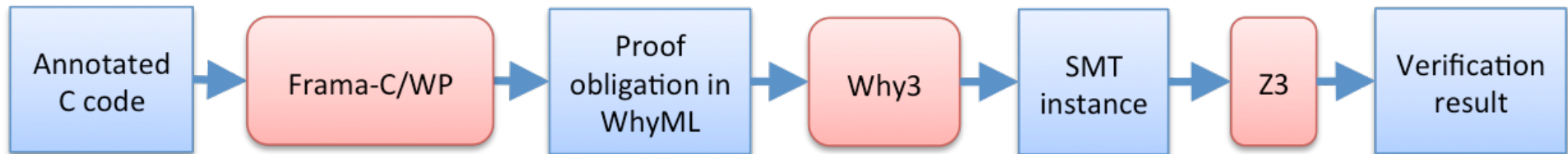
```
  @ ⇒ ((\at(y, k_n) == y_n) && ... && (\at(y, k_{3n}) == y_{3n}) &&
```

```
\*@ vector_equal((lin_comb(Dy, 1, a_0, ..., a_{n-1}) + lin_comb(Du, b_0, ..., b_n)), v)
```

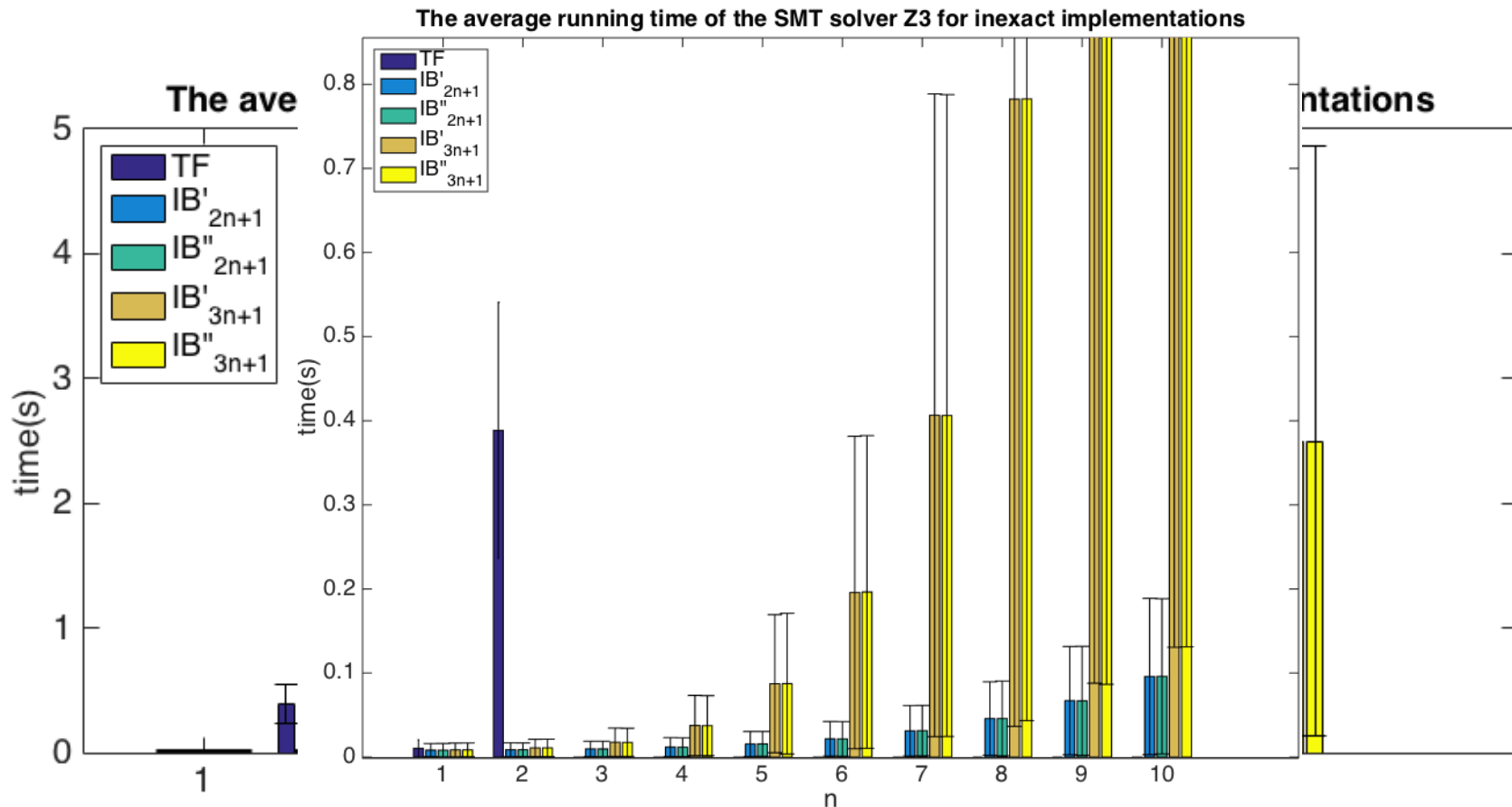
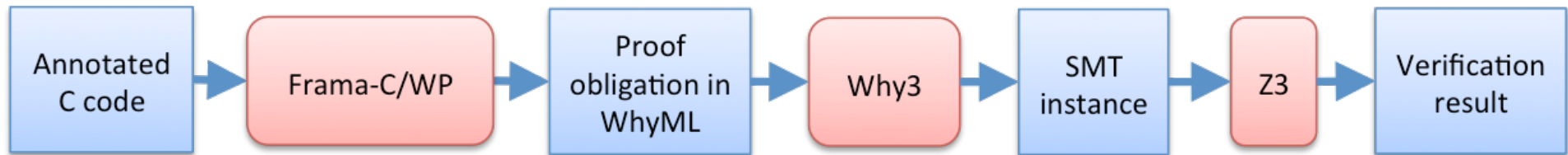
Framework For Automatic Verification



Automatic Verification for Exact Invariants



Automatic Verification – Inexact Invariants



A More Scalable Approach (SC)

Original Model

$$\mathbf{z}_{k+1} = \begin{bmatrix} -0.500311 & 0.16751 & 0.028029 & -0.395599 & -0.652079 \\ 0.850942 & 0.181639 & -0.29276 & 0.481277 & 0.638183 \\ -0.458583 & -0.002389 & -0.154281 & -0.578708 & -0.769495 \\ 1.01855 & 0.638926 & -0.668256 & -0.258506 & 0.119959 \\ 0.100383 & -0.432501 & 0.122727 & 0.82634 & 0.892296 \end{bmatrix} \mathbf{z}_k + \begin{bmatrix} 1.1149 & 0.164423 \\ -1.56592 & 0.634384 \\ 1.04856 & -0.196914 \\ 1.96066 & 3.11571 \\ -3.02046 & -1.96087 \end{bmatrix} \mathbf{u}_k$$

$$\mathbf{y}_k = \begin{bmatrix} 0.283441 & 0.032612 & -0.75658 & 0.085468 & 0.161088 \\ -0.528786 & 0.050734 & -0.681773 & -0.432334 & -1.17988 \end{bmatrix} \mathbf{z}_k$$

?

Code

```
void LTIS_step(void)
{
  {
    static const int_T colCidxRow0[5] = { 0, 1, 2, 3, 4 };
    const int_T *pCidx = &colCidxRow0[0];
    const real_T *pC0 = LTIS_ConstP.Internal_C;
    const real_T *xd = &LTIS_DW.Internal_DSTATE[0];
    real_T *y0 = &LTIS_Y.y[0];
    int_T numNonZero = 4;
    *y0 = (*pC0++) * xd[*pCidx++];
    while (numNonZero--) {
      *y0 += (*pC0++) * xd[*pCidx++];
    }
  }
  {
    static const int_T colCidxRow1[5] = { 0, 1, 2, 3, 4 };
    const int_T *pCidx = &colCidxRow1[0];
    const real_T *pC5 = &LTIS_ConstP.Internal_C[5];
    const real_T *xd = &LTIS_DW.Internal_DSTATE[0];
    real_T *y1 = &LTIS_Y.y[1];
    int_T numNonZero = 4;
    *y1 = (*pC5++) * xd[*pCidx++];
    while (numNonZero--) {
      *y1 += (*pC5++) * xd[*pCidx++];
    }
  }
}

real_T xnew[5];
int_T i;
xnew[0] = (0.87224)*LTIS_DW.Internal_DSTATE[0];
xnew[0] += (0.822174)*LTIS_U.u[0]+(-0.438008)*LTIS_U.u[1];
xnew[1] = (0.366378)*LTIS_DW.Internal_DSTATE[1];
xnew[1] += (-0.278536)*LTIS_U.u[0]+(-0.824313)*LTIS_U.u[1];
xnew[2] = (-0.540795)*LTIS_DW.Internal_DSTATE[2];
xnew[2] += (0.874484)*LTIS_U.u[0]+(0.858857)*LTIS_U.u[1];
xnew[3] = (-0.332664)*LTIS_DW.Internal_DSTATE[3];
xnew[3] += (-0.117628)*LTIS_U.u[0]+(-0.506362)*LTIS_U.u[1];
xnew[4] = (-0.204322)*LTIS_DW.Internal_DSTATE[4];
xnew[4] += (-0.955459)*LTIS_U.u[0]+(-0.622498)*LTIS_U.u[1];
for(i=0; i<5; i++) LTIS_DW.Internal_DSTATE[i] = xnew[i];
}
```

(2) Input-Output Similarity
Checking

Extracted Model

$$\hat{\mathbf{z}}_{k+1} = \begin{bmatrix} 0.87224 & 0 & 0 & 0 & 0 \\ 0 & 0.366378 & 0 & 0 & 0 \\ 0 & 0 & -0.540795 & 0 & 0 \\ 0 & 0 & 0 & -0.332664 & 0 \\ 0 & 0 & 0 & 0 & -0.204322 \end{bmatrix} \hat{\mathbf{z}}_k + \begin{bmatrix} 0.822174 & -0.438008 \\ -0.278536 & -0.824313 \\ 0.874484 & 0.858857 \\ -0.117628 & -0.506362 \\ -0.955459 & -0.622498 \end{bmatrix} \mathbf{u}_k$$

$$\mathbf{y}_k = \begin{bmatrix} -0.793176 & 0.154365 & -0.377883 & -0.360608 & -0.142123 \\ 0.503767 & -0.573538 & 0.170245 & -0.583312 & -0.56603 \end{bmatrix} \hat{\mathbf{z}}_k$$

(1) Model
Extraction

Model Extraction

- Use symbolic execution to identify transition relation

fragment of
step function

```

const ConstP_LTIS_T LTIS_ConstP = {
    { -0.793176, 0.154365, -0.377883, -0.360608, -0.142123,
      0.503767, -0.573538, 0.170245, -0.583312, -0.56603 } };
static const int_T colCidxRow0[5] = { 0, 1, 2, 3, 4 };
const int_T *pCidx = &colCidxRow0[0];
const real_T *pC0 = LTIS_ConstP.Internal_C;
const real_T *xd = &LTIS_DW.Internal_DSTATE[0];
real_T *y0 = &LTIS_Y.y[0];
int_T numNonZero = 4;
*y0 = (*pC0++) * xd[*pCidx++];
while (numNonZero--) {
    *y0 += (*pC0++) * xd[*pCidx++];
}

```



symbolic execution

big-step
transition
relation

$$y[0]^{(new)} = (((((-0.793176 \cdot x[0]) + (0.154365 \cdot x[1])) + (-0.377883 \cdot x[2])) + (-0.360608 \cdot x[3])) + (-0.142123 \cdot x[4]))$$

y stands for **LTIS_Y.y**, and **x** stands for **LTIS_DW.Internal_DSTATE**

Model Extraction (cont.)

- Identify the set of state variables V_{state}

$$V_{state} = (V_{updated} \setminus V_{output}) \cup (V_{used} \setminus V_{input})$$

- Transform into matrix form

$$y[0]^{(new)} = (((((-0.793176 \cdot x[0]) + (0.154365 \cdot x[1])) + (-0.377883 \cdot x[2])) + (-0.360608 \cdot x[3])) + (-0.142123 \cdot x[4])) \quad \text{transition relation}$$



$$= -0.793176 \cdot x[0] + 0.154365 \cdot x[1] + -0.377883 \cdot x[2] + -0.360608 \cdot x[3] + -0.142123 \cdot x[4] \quad \text{canonical form}$$



$$= [-0.793176, 0.154365, -0.377883, -0.360608, -0.142123] \cdot \mathbf{x} + [0, 0] \cdot \mathbf{u} \quad \text{vector form}$$



$$\begin{aligned} \mathbf{x}^{(new)} &= \hat{\mathbf{A}}\mathbf{x} + \hat{\mathbf{B}}\mathbf{u} \\ \mathbf{y}^{(new)} &= \hat{\mathbf{C}}\mathbf{x} + \hat{\mathbf{D}}\mathbf{u} \end{aligned} \quad \text{matrix form}$$

Input-output equivalence checking

- Check similarity between two models

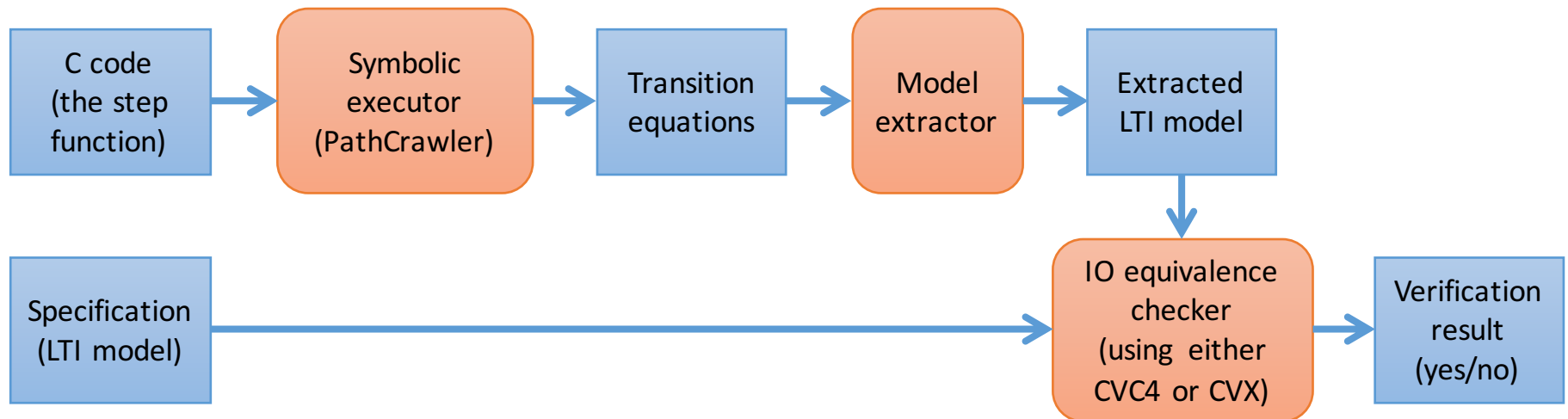
Two minimal LTI models $\Sigma(A, B, C, D)$ and $\hat{\Sigma}(\hat{A}, \hat{B}, \hat{C}, \hat{D})$ are input-output equivalent iff there exists a non-singular matrix T such that

$$\hat{A} = \mathbf{TAT}^{-1}, \quad \hat{B} = \mathbf{TB}, \quad \hat{C} = \mathbf{CT}^{-1}, \quad \text{and} \quad \hat{D} = \mathbf{D}$$

- Find the existence of similarity transformation matrix using
 - SMT formulation approach
 - Convex optimization formulation approach
- Need to tolerate the numerical errors on the model parameters

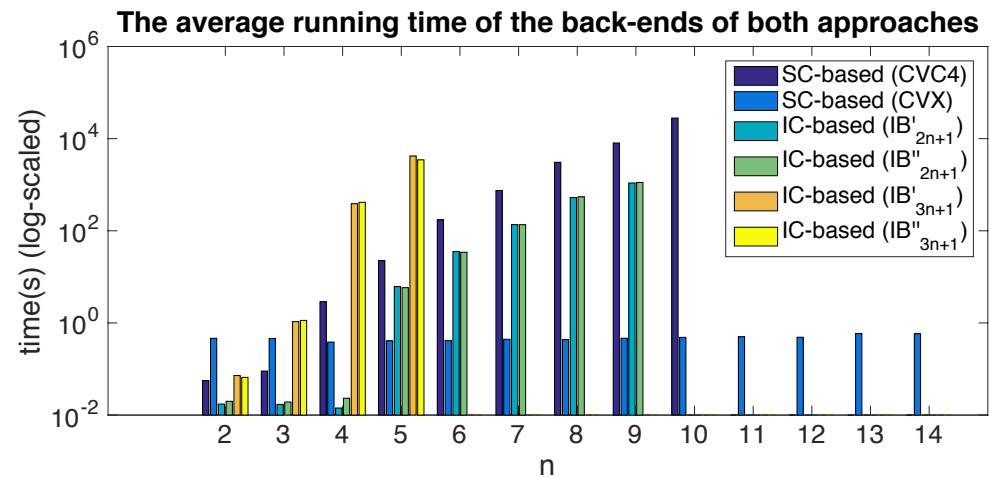
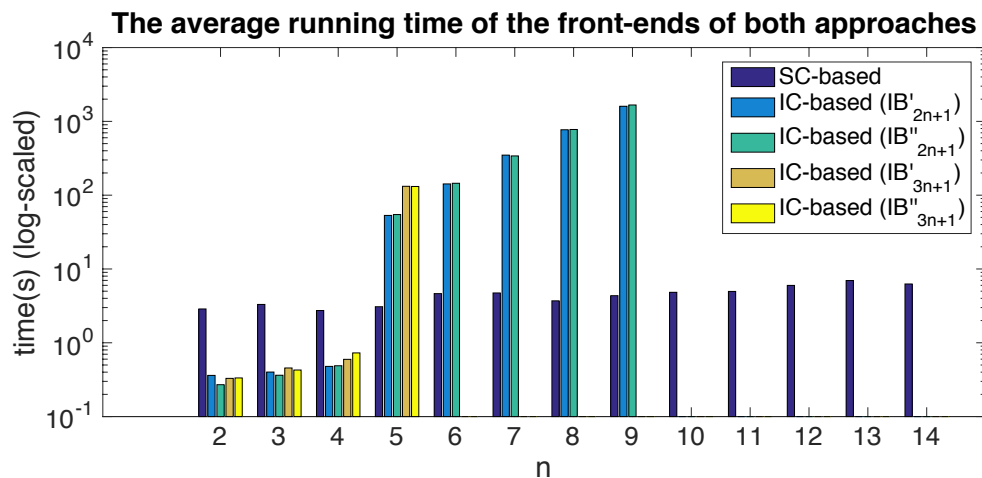
Verification Toolchain

- Similarity Checking (SC)-based approach



Evaluation

- Compare scalability of the two approaches
 - Random LTI models with a range of state sizes
 - Code obtained by Simulink Coder
- Similarity-checking approach (SC) dramatically outperforms invariant-checking approach (IC)



Current work

- Focus on more complex controllers
 - Convex optimization-based controllers

Thank You

