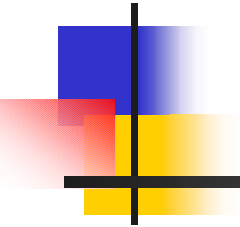


GALOISCONNECTIONS

purely functional



Cryptol Tutorial

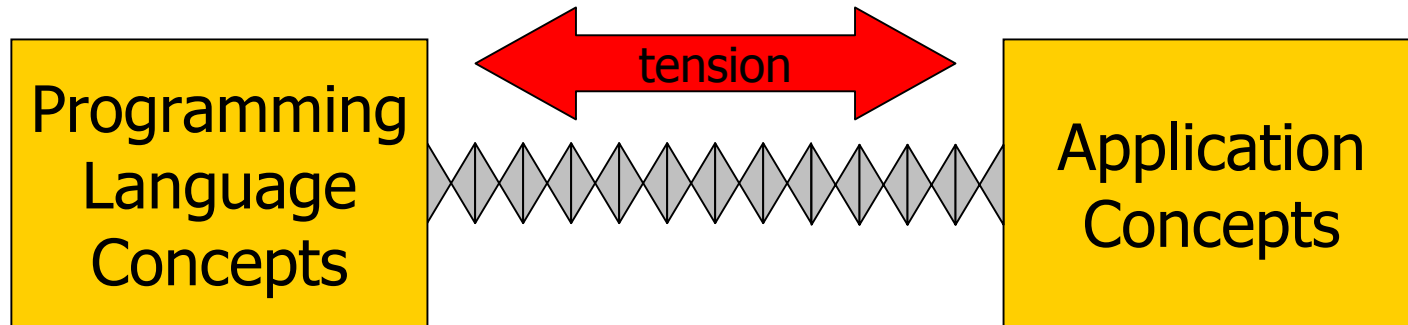
Overview and Elements



Purpose of the tutorial

- Overview of Cryptol
 - Domain-specific language for crypto-algorithms
 - Elements of Cryptol
- Cryptol in use for full crypto-algorithm
 - IDEA
 - Idioms
- Non-goal
 - Principles and technology underlying the Cryptol compiler — where does all the code come from?

Tension



- Domain-specific languages (DSLs) attempt to bridge this semantic gap
- Programs are written in domain-specific terms
- Programs “execute” as if a program had been written



Classic examples

- **Spreadsheets**: Accountancy concepts and notations
- **Parser generators**: LEX, YACC: use BNF grammars
- **SQL**: Relational database queries
- **PERL**: Text manipulation scripting
- **TeX** and **LaTeX**: Document layout
- **Postscript**: Low-level graphics
- **Mathematica** / **Maple**: Symbolic computation



Value of DSLs

- Design-level programming
 - Huge productivity increase
 - Major flexibility for code evolution
 - Natural maintenance of design documents
- Multiple use
 - Code
 - Test generation
 - Analysis

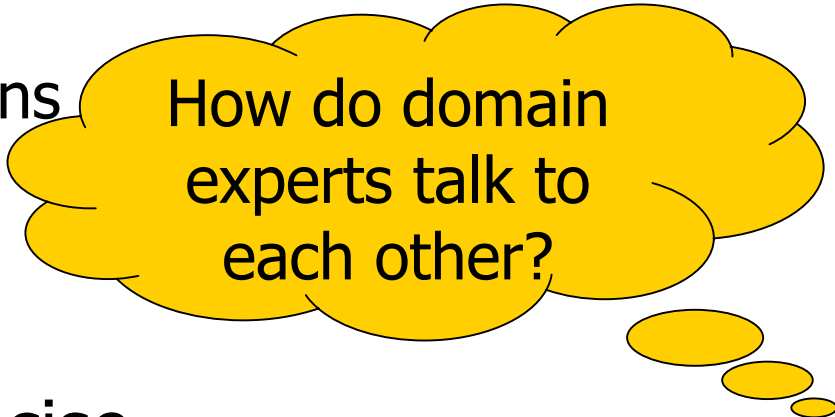


Where do DSLs come from?

- Domain analysis

- Existing domain notations

- Textual
 - Mathematical
 - Graphical



How do domain experts talk to each other?

- Semantics must be precise

- Prototype interpretation must match compiled interpretation must match testing interpretation etc.
 - Source level reasoning
 - DSL programmers may not understand traditional programming



Crypto-algorithm domain analysis

- What we did

- Spoke with crypto-algorithm designers
 - What are the important elements of algorithm specification?
- Studied five AES finalists and DES
 - What do these algorithms have in common?
 - What differences occur between them?
- Examined previous attempts at crypto DSL

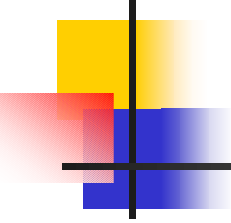
- Now

- Embodying the domain analysis within a language
- Obtaining feedback from crypto specialists



Requirements on Cryptol

- Focus on symmetric-key block algorithms
- Emphasize cryptographic concepts and abstractions
 - Rather than traditional computing abstractions
- Specifications expressed in Cryptol should lend themselves to mathematical analysis
 - Capture mathematical structure rather than detailed representation issues
 - Issues of space use and memory allocation should not cloud the specification of crypto algorithms



Consequently, Cryptol should provide:

- Crypto data-types plus operators
 - Fixed-width words, small matrices and vectors, Galois fields
 - Flexible views of data without excessive annotations by the user
- Ability to express parameterized algorithms
 - Work over varying key sizes and varying block sizes
- Appropriate control structures for crypto-algorithms
 - Including iteration and recurrence
- Predefined standard cryptographic modes
 - Plus facilities for describing new modes



Block Ciphers

- Logical pattern of use

encrypt : (Xkey, PT) \rightarrow CT

decrypt : (Xkey, CT) \rightarrow PT

keySchedule : Key \rightarrow Xkey

- In the context of different modes
 - ECB, CBC, etc...



Bit Vectors

- Sizes ranging from 4 bits to 128 bits
 - 8 and 32 most common
- All the usual boolean ops
- Simple modulo arithmetic (+, -, *, /)
- Permutations
 - Mostly just rotations of bit vectors
 - More general permutation used in DES
- Splitting and combining bit vectors
 - Big and little endian



Matrices

- Matrix/vector multiplication used in TwoFish and Rijndael

$$\begin{aligned}x_i &= \lfloor X/2^{8i} \rfloor \bmod 2^8 \quad i = 0, \dots, 3 \\y_i &= s_i[x_i] \quad i = 0, \dots, 3 \\ \begin{pmatrix} z_0 \\ z_1 \\ z_2 \\ z_3 \end{pmatrix} &= \begin{pmatrix} \cdot & \dots & \cdot \\ \vdots & \text{MDS} & \vdots \\ \cdot & \dots & \cdot \end{pmatrix} \cdot \begin{pmatrix} y_0 \\ y_1 \\ y_2 \\ y_3 \end{pmatrix}\end{aligned}$$

- Other arithmetics
 - Polynomials
 - Galois fields

$$Z = \sum_{i=0}^3 z_i \cdot 2^{8i}$$



Iteration

- For loops
 - Over fixed counts
- Feistel networks (TwoFish)
- SP-networks (MARS, RC6, ...)
- Recurrence relations



Common practice

- Use of arrays and in-place update
 - Not fundamental to the domain
 - Common implementation technique
 - Underlying model is recurrence relations

$$\begin{aligned}(F_{r,0}, F_{r,1}) &= F(R_{r,0}, R_{r,1}, r) \\ R_{r+1,0} &= \text{ROR}(R_{r,2} \oplus F_{r,0}, 1) \\ R_{r+1,1} &= \text{ROL}(R_{r,3}, 1) \oplus F_{r,1} \\ R_{r+1,2} &= R_{r,0} \\ R_{r+1,3} &= R_{r,1}\end{aligned}$$



Cryptol



Cryptol

Domain-specific
language for
cryptoalgorithms

- Now that we know the domain ...

... what should the domain-specific language be like?

Data in Cryptol

- The smallest elements: Bits
- Everything else is a matrix (a parameterized collection)

7 single bits

`[False True False True False False True]`

`0x4A`

7 (or more) bits

4 elements, each
8 (or more) bits

`[0x3F 0x02 0x41 0xD8]`

2 elements, each
having 4 elements,
each 4 (or more) bits

`[[1 2 3 4] [5 6 7 8]]`

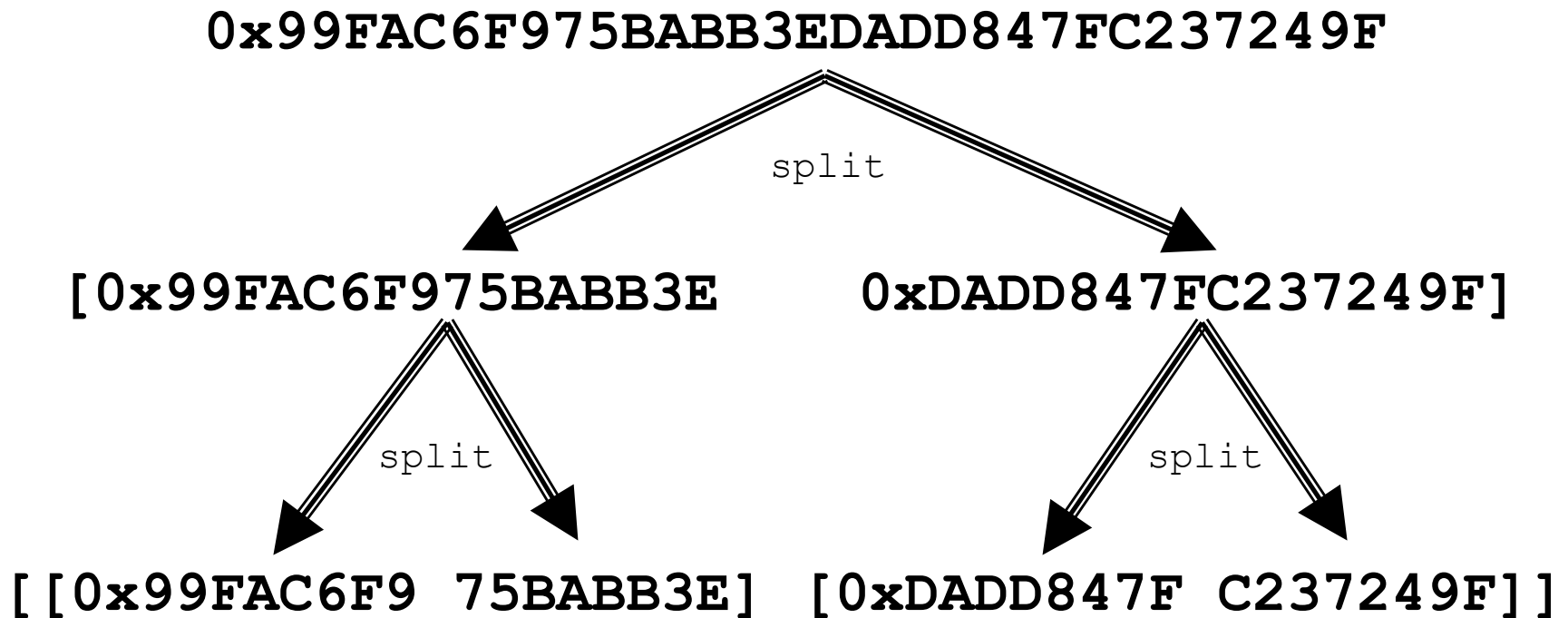
`[1 .. 10]`

10 elements, each
of 4 (or more) bits



Hierarchical Views of Data

- Data can be split into n subparts





Primitive Operations

- Arithmetic operators
 - Result is modulo the word size of the arguments
 - $+$ $-$ $*$ $/$ $\%$
- Boolean operators
 - From bits, to arbitrarily nested matrices of the same shape
 - $\&$ $|$ $^$ \sim
- Shift and rotate operators
 - \ll \gg \lll \ggg
- Comparison operators
 - Equality, order
 - $==$ $<$ $<=$ $>$ $>=$ $!=$
- Conditional operator
 - Expression-level *if-then-else*
 - Like C's $a?b:c$
 - Booleans are just bits



Matrices

- Matrix operators
 - Concatenation, indexing, size
 - # @ @@ width

`[1..5] # [3 6 8] = [1 2 3 4 5 3 6 8]`

- Zero-based indexing from the left

`[50 .. 99] @ 10 = 60`



Matrix Operations

- Logical operations lift to matrices of any size
- Arithmetic operations lift to matrices of words
 - Word = matrix of bits

- Bulk indexing

$$[50 \dots 99] @@ [10 \dots 20] = [60 \dots 70]$$

- Permutations

$$[1 \dots 4] @@ [1 \ 2 \ 3 \ 0] = [2 \ 3 \ 4 \ 1]$$

$$[1 \dots 4] @@ [3 \ 2 \ \dots \ 0] = [4 \ 3 \ 2 \ 1]$$



Cryptol Definitions

- Definitions of values or of functions

```
x = 13;
```

```
incr x = x + 1;
```

```
f (x, y) = 2 * x + 3 * y + 1;
```

- Pattern Matching on Matrices

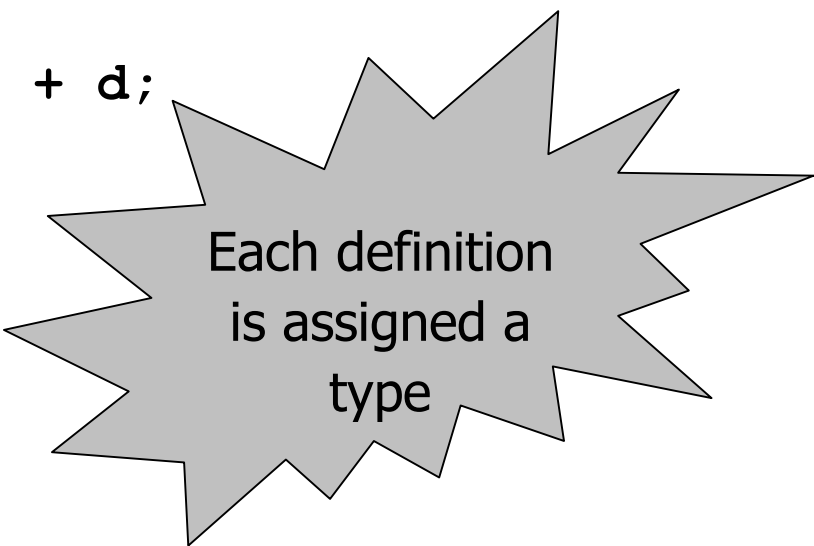
```
sum4 [a b c d] = a + b + c + d;
```

- Nested definitions

```
f x = [y z]
```

```
  where {y = x + 1;
```

```
         z = not x};
```



Each definition
is assigned a
type



Simple Cryptol Size Types

- Data sizes

<code>Bit</code>	single bit
<code>[32]</code>	32-bit word (same as <code>[32]Bit</code>)
<code>[16] [48]</code>	sixteen 48-bit words

- Functions

- Form is: *argument -> result*



Examples

`x : [4];`

`x = 13;`

`incr : [32] -> [32];`

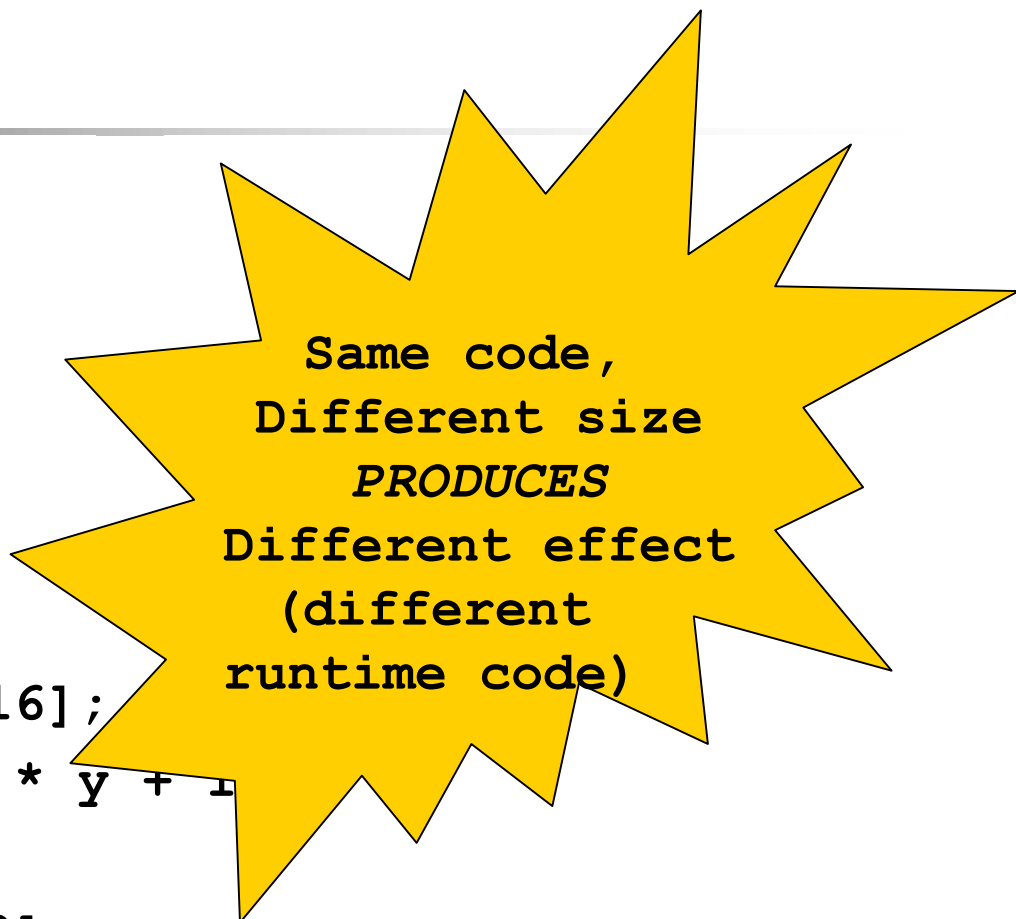
`incr x = x + 1;`

`f : ([16], [16]) -> [16];`

`f (x, y) = 2 * x + 3 * y + 1`

`sum4 : [4][32] -> [32];`

`sum4 [a b c d] = a + b + c + d;`



Same code,
Different size
PRODUCES
Different effect
(different
runtime code)



Widths

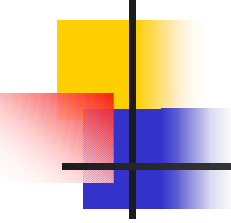
- Often it is important to know the size of a structure at the top level
 - Number of bits in a word
 - Number of rows in a matrix

`width 0x5f = 7`

`width [[2 3] [4 5] [6 7]] = 3`

- Width is approximately the \log_2 of a number

Equational Correspondence


$$\begin{aligned}(F_{r,0}, F_{r,1}) &= F(R_{r,0}, R_{r,1}, r) \\ R_{r+1,0} &= \text{ROR}(R_{r,2} \oplus F_{r,0}, 1) \\ R_{r+1,1} &= \text{ROL}(R_{r,3}, 1) \oplus F_{r,1} \\ R_{r+1,2} &= R_{r,0} \\ R_{r+1,3} &= R_{r,1}\end{aligned}$$

round [R0 R1 R2 R3] r = [S0 S1 R0 R1]

where { [F0 F1] = F (R0, R1, r) ;

S0 = (R2 ^ F0) >>> 1;

S1 = (R3 <<< 1) ^ F1;

} ;



Bounded Iteration

- Borrowed the comprehension notion from set theory
 - $\{ \mathbf{a} + \mathbf{b} \mid \mathbf{a} \in \mathbf{A}, \mathbf{b} \in \mathbf{B} \}$
 - Adapted to matrices (i.e. sequences)
- Applying an operation to each element

```
[ | 2*x + 3 | | x <- [1 2 3 4] | ]  
= [5 7 9 11]
```



Traversals

- Cartesian traversal

$$\begin{aligned} & [| \ [x \ y] \ || \ x \leftarrow [0..2], \ y \leftarrow [3..4] \ |] \\ &= \ [[0 \ 3] \ [0 \ 4] \ [1 \ 3] \ [1 \ 4] \ [2 \ 3] \ [2 \ 4]] \end{aligned}$$

- Parallel traversal

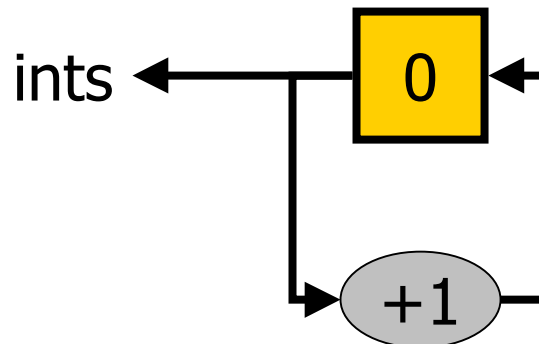
$$\begin{aligned} & [| \ x + y \ || \ x \leftarrow [1..3] \\ & \qquad \qquad \qquad || \ y \leftarrow [3..7] \ |] \\ &= \ [4 \ 6 \ 8] \end{aligned}$$

Recurrence

- Textual description of shift circuits
 - Traditionally use a language of commands
 - Arrays, updates, and command-loops
 - Alternatively, use stream-equations
 - Stream-definitions can be *recursive*

```
ints : [inf][256]
```

```
ints = [0] # [| y+1 || y <- ints |];
```





Sizes of Recurrences

- Stream is unbounded
 - No finite size
 - But generated by a finite state machine
- Use the `inf` type

```
ints : [inf][8];  
ints = [0] # [| y+1 || y <- ints |];
```



Parallel Traversal

- Define an unbounded sequence of factorials

```
facts : [inf] [32];  
facts = [1] # [| x * n || x <- facts  
                || n <- [1..] ||];
```



Cryptol “Execution”

Execution by calculation

- Example

- `facts @@ [0..10]`

is a mathematical expression whose result will be computed

- Storage (and other issues) are left to the Cryptol system



Multiple Access Points

- Define an unbounded sequence of fibonacci numbers, parameterized on starting pair

```
fibs : ([32],[32]) -> [inf][32];  
fibs (p,q) = xs  
  where  
    xs = [p q]  
          # [| x + y || x <- xs  
              || y <- drop (1,xs) |];
```

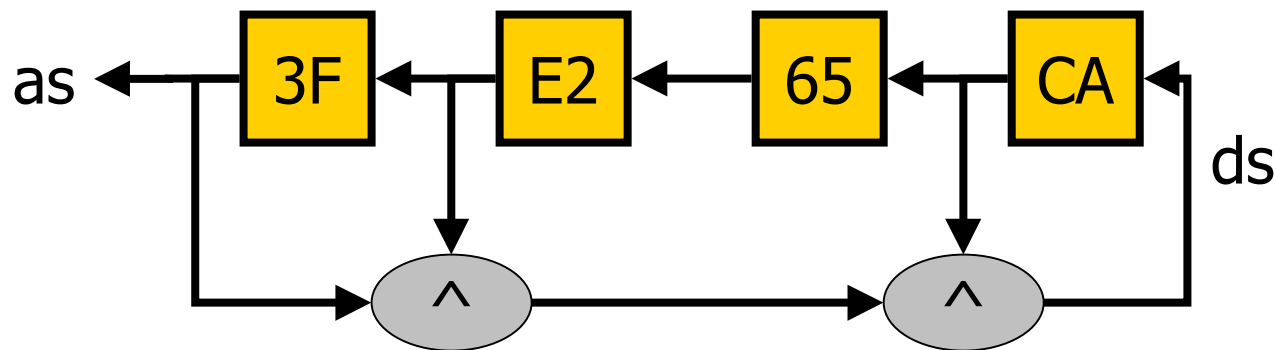

More Complex Stream Equations

```
as = [0x3F 0xE2 0x65 0xCA] # ds;
```

```
ds = [| a ^ b ^ c || a <- as
```

```
    || b <- drop(1, as)
```

```
    || c <- drop(3, as) ||];
```

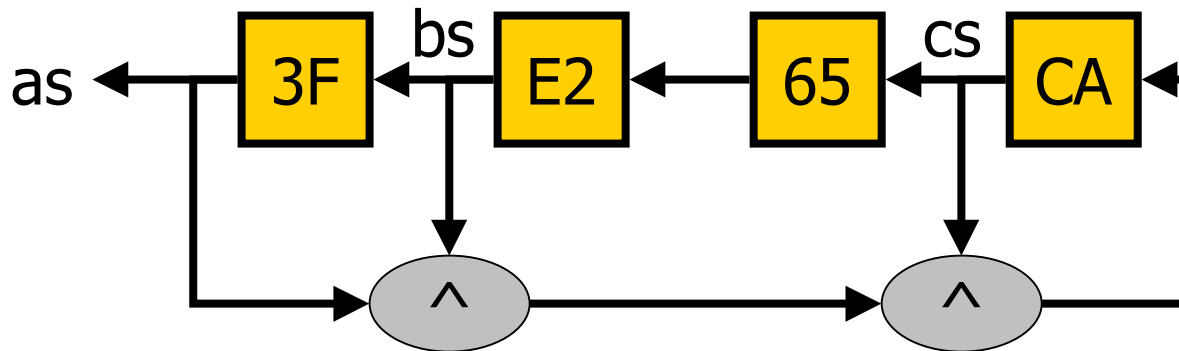


Alternative Description

`as = [0x3F] # bs;`

`bs = [0xE2 0x65] # cs;`

`cs = [0xCA] # [| a ^ b ^ c || a<-as
|| b<-bs
|| c<-cs |];`



Additional Complexity

`as = [0x3F 0xE2 0x65]`

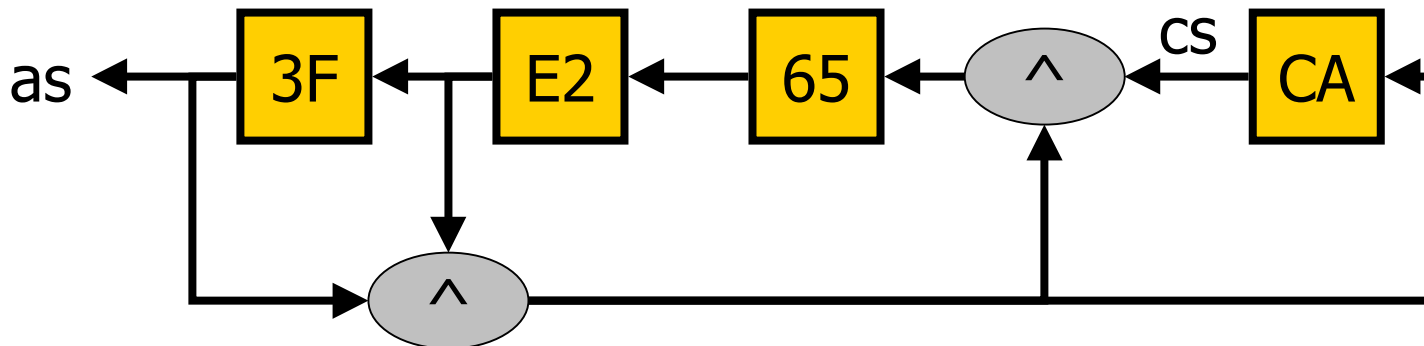
`# [| c^c' || c <- cs`

`|| c' <- drop(1,cs) ||];`

`cs = [0xCA] # [| a^a'`

`|| a <- as`

`|| a' <- drop(1,as) ||];`





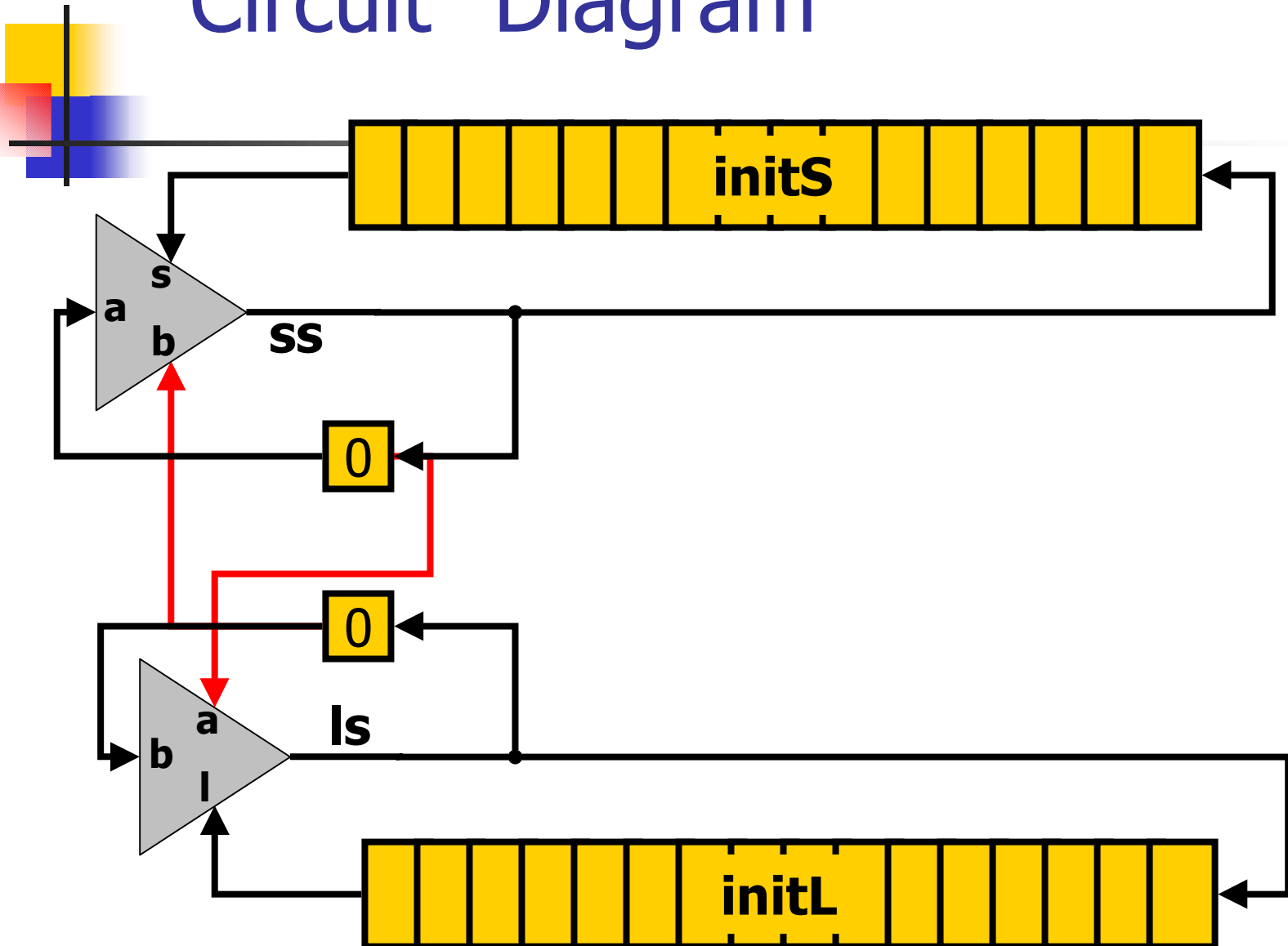
RC6 Key Expansion

- Published specification is in terms of array updates
 - Imperative code for key expansion appears entirely symmetrical — Cryptol exposes non-symmetry

```
ss = [ (s+a+b) <<< 3 || s <- initS # ss
      || a <- [0] # ss
      || b <- [0] # ls ];
```

```
ls = [ (l+a+b)<<<(a+b) || l <- initL # ls
      || a <- ss
      || b <- [0] # ls ];
```

"Circuit" Diagram

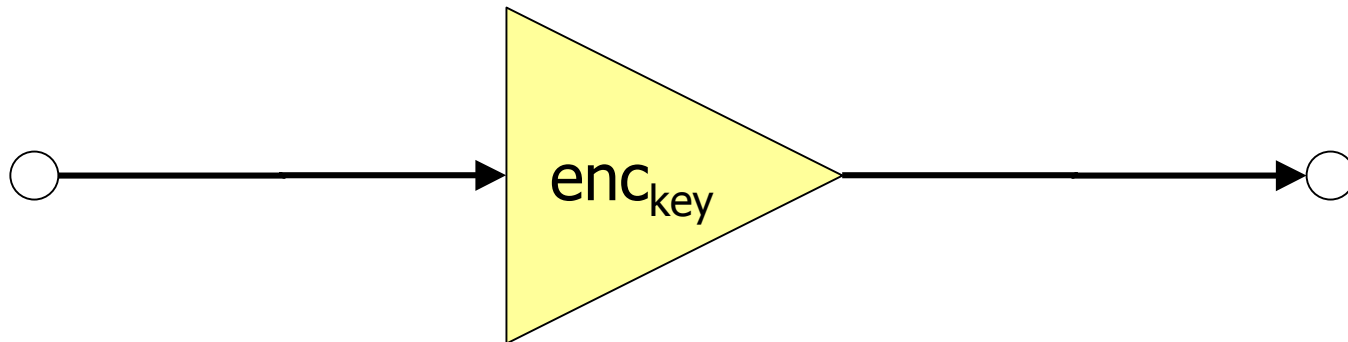


Electronic code book

$\text{ecb}(\text{pt}, \text{key}) = \text{ct}$

where

$\text{ct} = [\mid \text{encrypt}(\text{x}, \text{key}) \mid \mid \text{x} \leftarrow \text{pt} \mid]$



Cipher Block Chaining

`cbc(iv, pt, key) = ct`

where

```
ct = [iv] # [| encrypt (x^y, key) || x <- pt  
           || y <- ct |]
```

