

Inconsistencies in Specification of Intel TDX Remote Attestation

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Promise of talk

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- Complexity is the **best friend** of Intel!

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- 1 Introduction
- 2 Formal Security Analysis Approach
- 3 TDX
 - Discrepancies Identified
 - Formal Specification
 - Automated Verification
- 4 Summary


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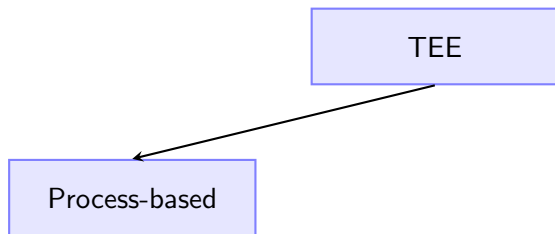
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- Without attestation, no better than conventional computing for possible threat models
 - Remote user cannot distinguish a malicious platform and a genuine one
 - Even with alternative of attestation: authentication
- “Any attack that could compromise the attestation of a TEE instance could lead to a workload or data being compromised in turn.”⁴

³Confidential Computing Consortium, *A Technical Analysis of Confidential Computing*, v1.1, 2021

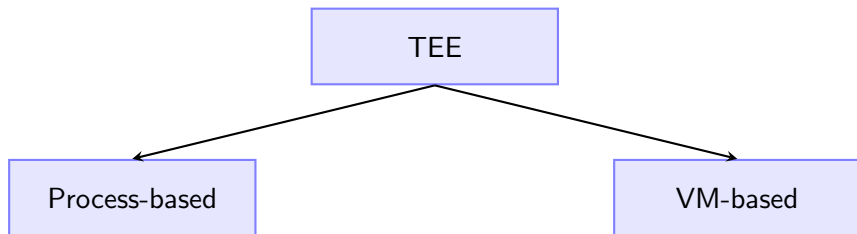
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TEEs Granularity (Public cloud commercial solutions)



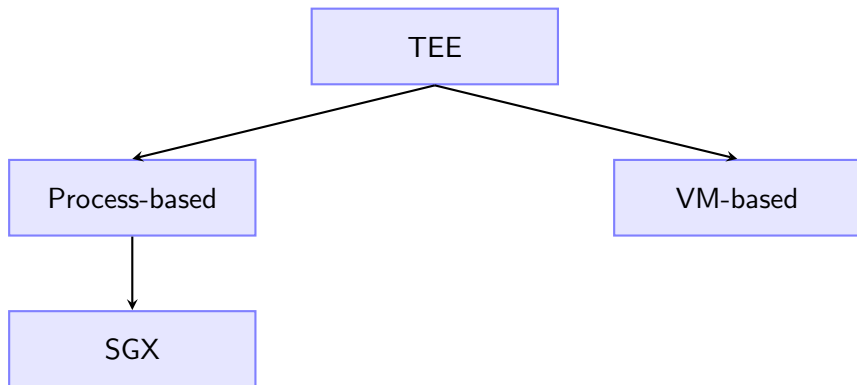
- Smaller TCB

TEEs Granularity (Public cloud commercial solutions)

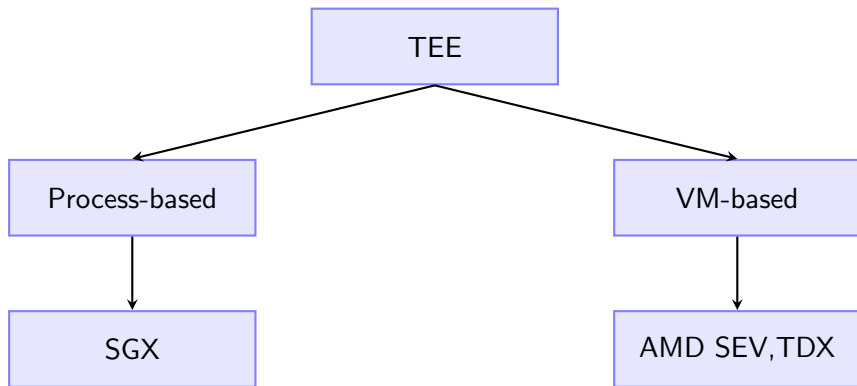


- Ease of use

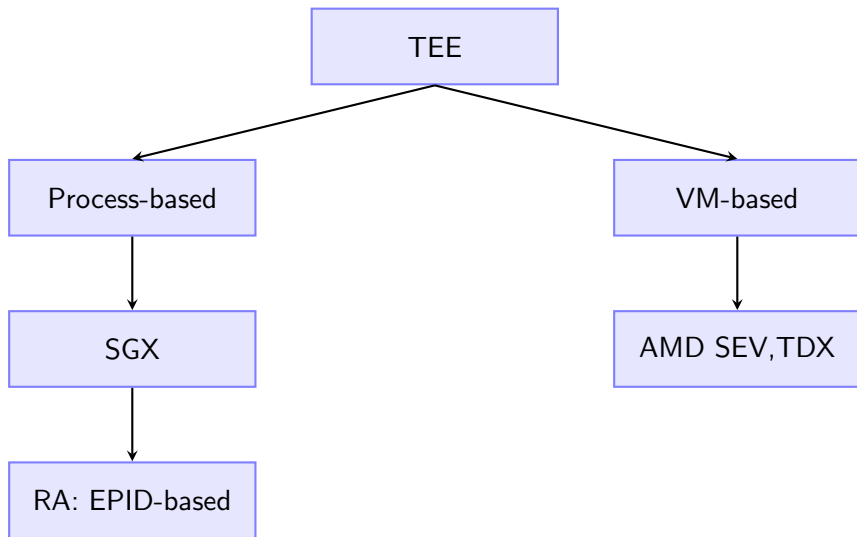
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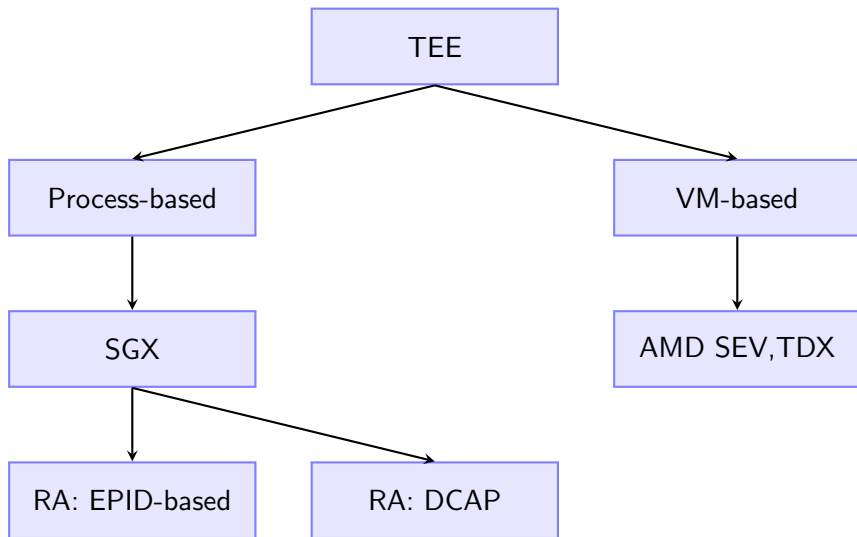
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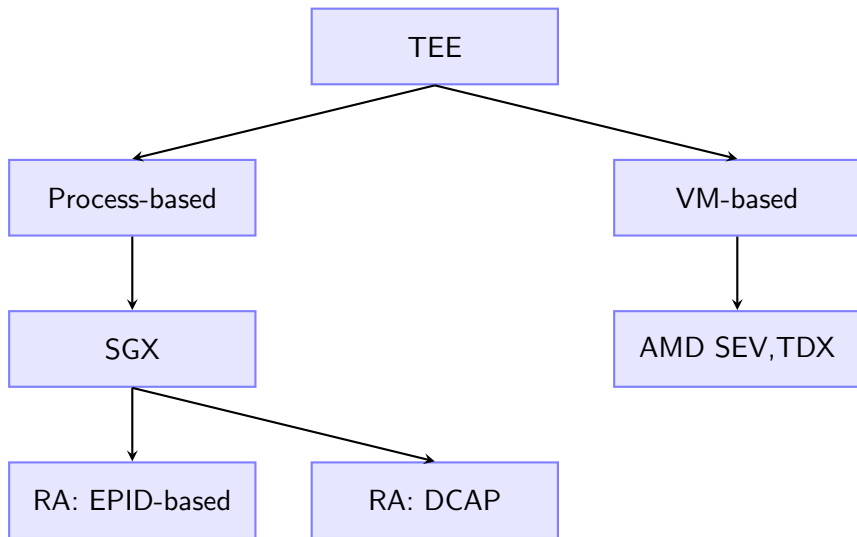
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- Different report generation mechanism
- Runtime TD measurements

Outline

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ProVerif vs. Tamarin prover

- More automation vs. user interaction

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- Faster⁶

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⁶Lafourcade and Puits, "Performance Evaluations of Cryptographic Protocols Verification Tools Dealing with Algebraic Properties", 2016

Workflow of the Analysis Approach

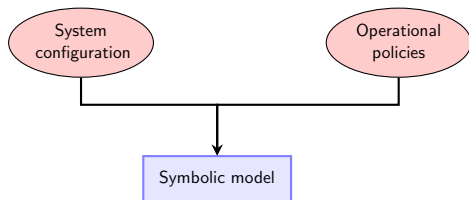
System
configuration

Workflow of the Analysis Approach

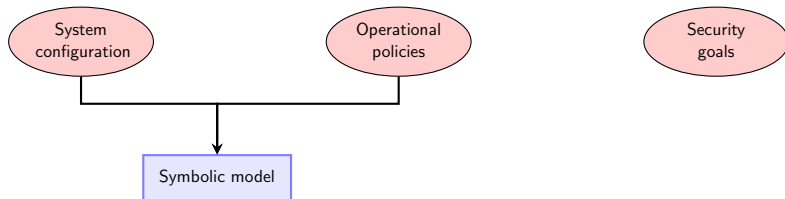
System
configuration

Operational
policies

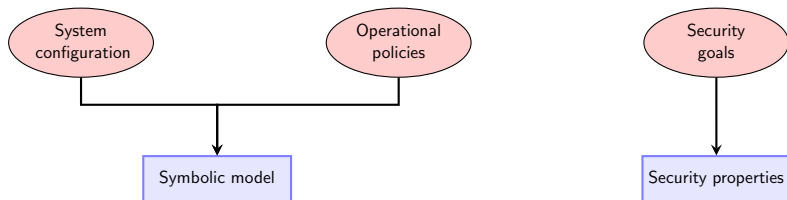
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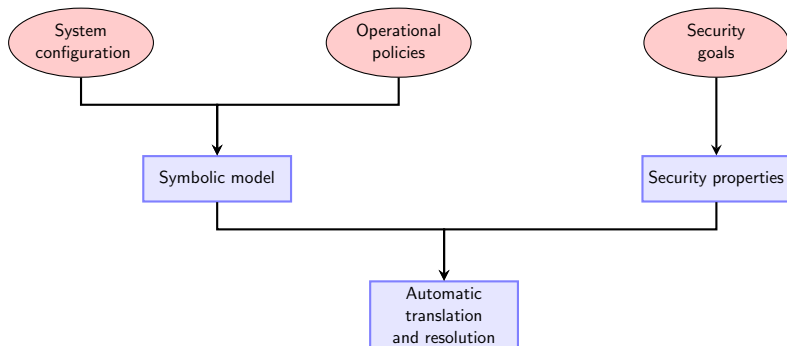
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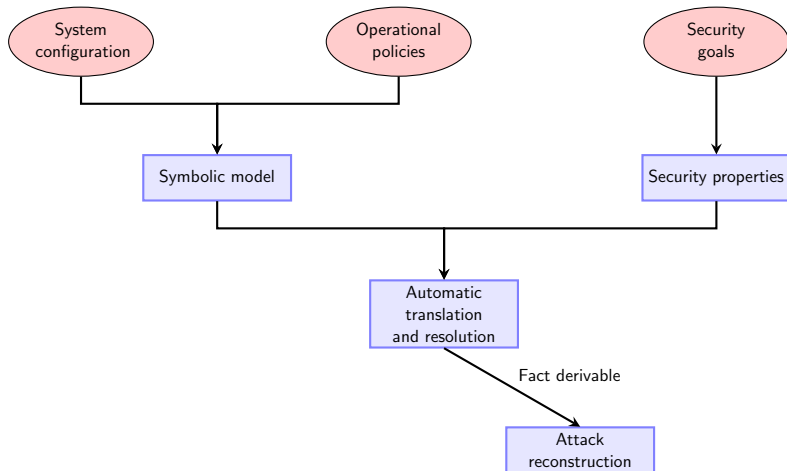
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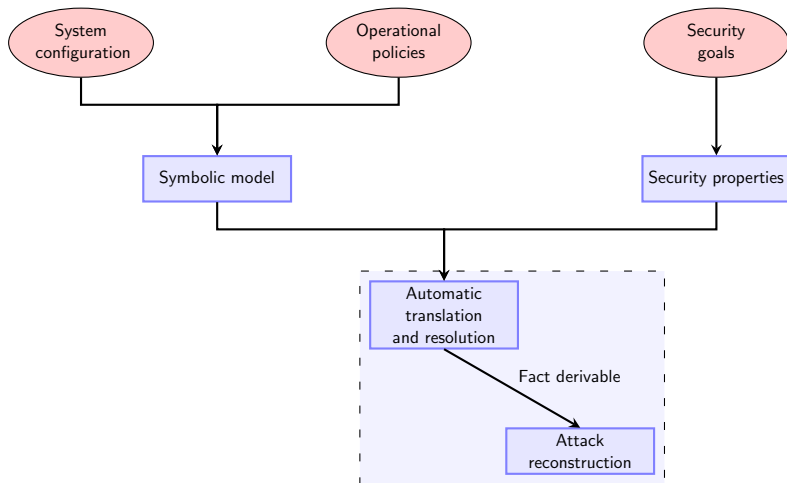
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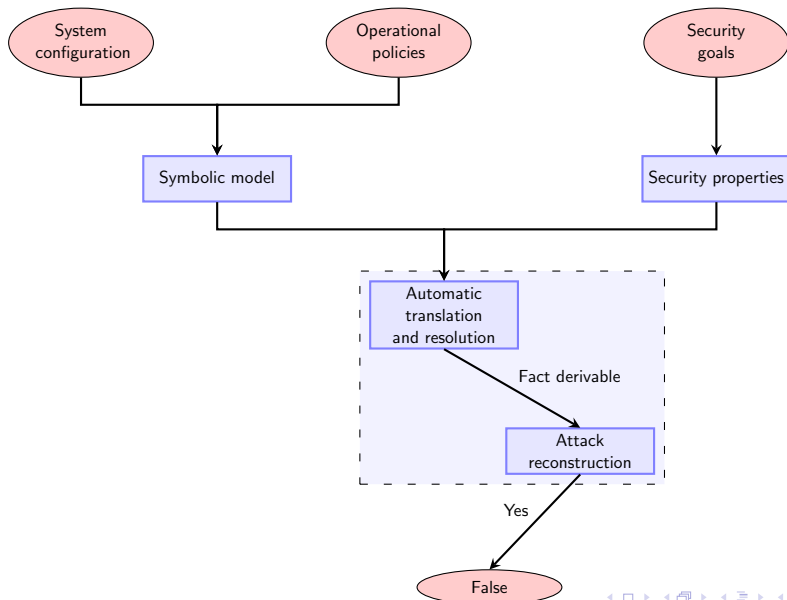
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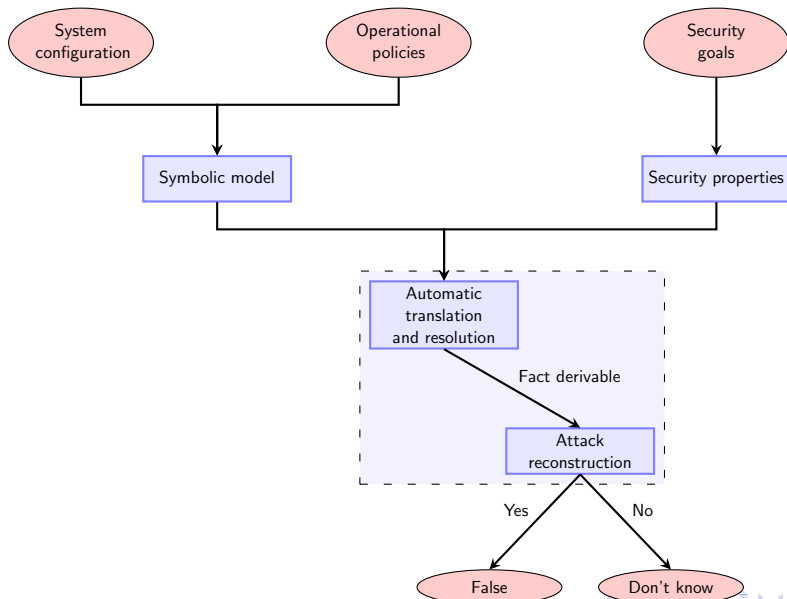
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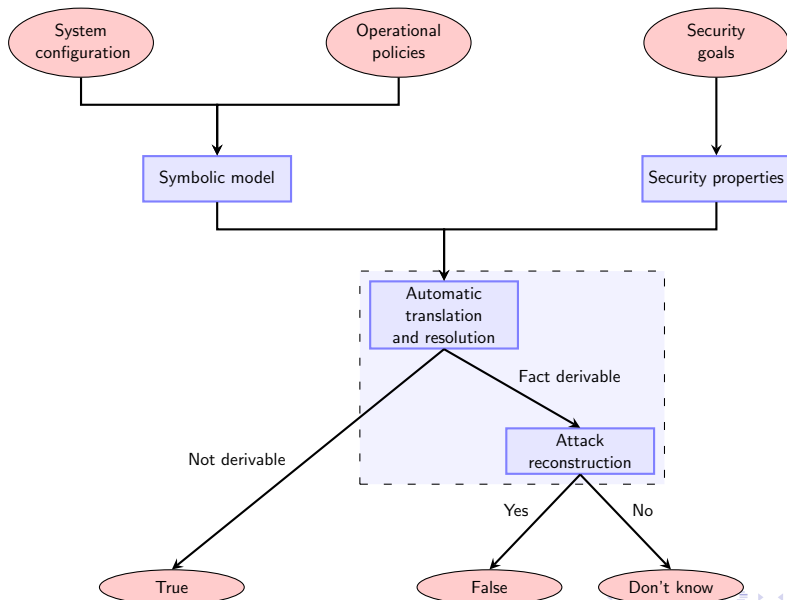
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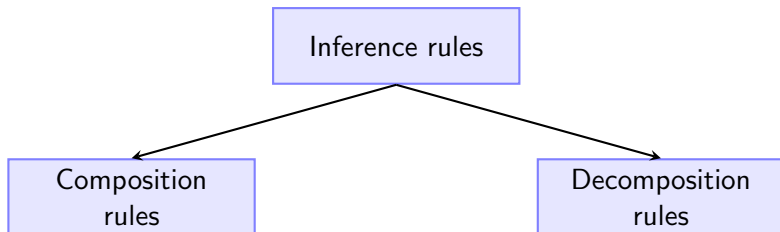


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Inference System and Horn Clauses (Simplified)

- Composition rules

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 - pair $\frac{x \quad y}{\langle x, y \rangle} \quad att(x) \wedge att(y) \rightarrow att(\langle x, y \rangle)$

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- verifysign $\frac{vpk(sk) \quad m \quad signAppDet(sk, m)}{true}$

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- Identification of **discrepancies** including inconsistent information

⁷Blanchet et al., "Modeling and verifying security protocols with the applied pi calculus and ProVerif", 2016

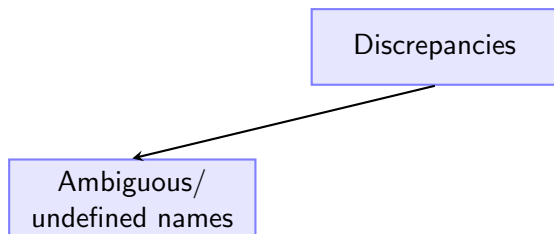
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- Precise **specification** of TD attestation protocol in ProVerif⁷

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- Identification of **discrepancies** including inconsistent information
- Precise **specification** of TD attestation protocol in ProVerif⁷
- Automated **verification** of confidentiality and authentication properties in ProVerif

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Discrepancies Identified

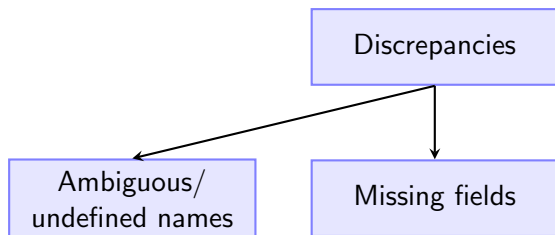


- SEAMINFO vs. TEE_TCB_INFO (e.g., p.2-8)⁸

⁸Intel, *Intel® Trust Domain CPU Architectural Extensions*, 2020

⁹Intel, *Architecture Specification: Intel® Trust Domain Extensions (Intel® TDX) Module*, 2020

Discrepancies Identified

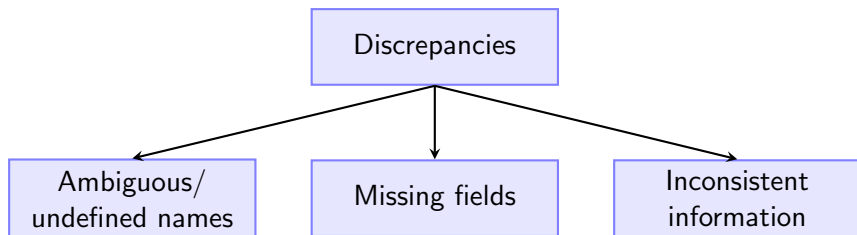


- MROWNERCONFIG missing in TDINFO (Fig. 10.1, p.85)⁹

⁸Intel, *Intel® Trust Domain CPU Architectural Extensions*, 2020

⁹Intel, *Architecture Specification: Intel® Trust Domain Extensions (Intel® TDX) Module*, 2020

Discrepancies Identified



⁸Intel, *Intel® Trust Domain CPU Architectural Extensions*, 2020

⁹Intel, *Architecture Specification: Intel® Trust Domain Extensions (Intel® TDX) Module*, 2020

Inconsistent Information: Example 1¹⁰

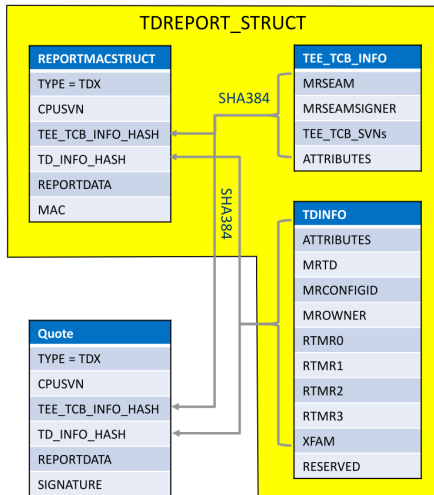


Figure 10.1: TDX Measurement Reporting

¹⁰Intel, Architecture Specification: Intel® Trust Domain Extensions (Intel® TDX) Module, 2020

Inconsistent Information: Example 1¹¹

`tmp_seamreport.REPORTMACSTRUCT.TEE_TCB_INFO_HASH = SHA384(tmp_seamreport.TEE_TCB_INFO);`

Table 2-3. TEE_TCB_INFO Structure

Name	Offset (Bytes)	Size (Bytes)	Description
VALID	0	8	Indicates TEE_TCB_INFO fields which are valid. <ul style="list-style-type: none">1 in the i-th significant bit reflects that the 8 bytes starting at offset (8 * i) are valid.0 in the i-th significant bit reflects that either 8 bytes starting at offset (8 * i) is not populated or reserved, and is set to zero.
TEE_TCB_SVN	8	16	TEE_TCB_SVN array.
MRSEAM	24	48	Measurement of the Intel TDX module.
MRSIGNERSEAM	72	48	Measurement of TDX module signer if valid.
ATTRIBUTES	120	8	Additional configuration ATTRIBUTES if valid.
RESERVED	128	111	Must be zero.

¹¹Intel, Intel (®) Trust Domain CPU Architectural Extensions, 2020

Inconsistent Information: Example 2¹²

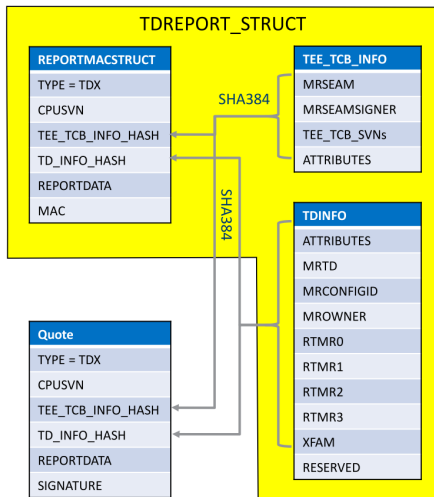



Figure 10.1: TDX Measurement Reporting

RESERVED is not a part of hash!

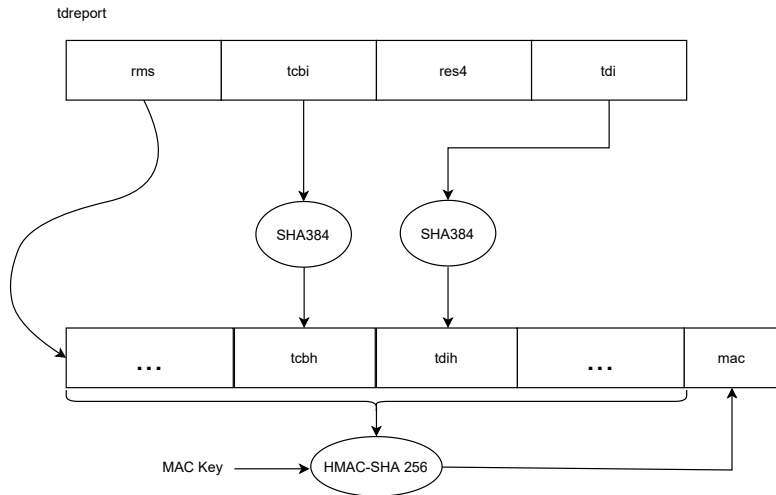
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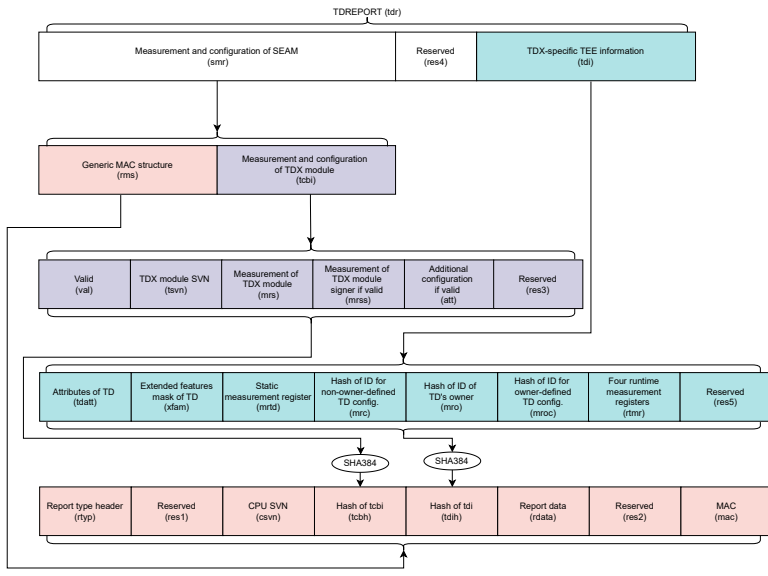
Software verifying a TEE report structure (for TDX, this includes TEE_TCB_INFO_STRUCT and TDINFO_STRUCT) should first confirm that its REPORTMACSTRUCT.TEE_TCB_INFO_HASH equals the hash of the TEE_TCB_INFO_STRUCT (if applicable) and that REPORTMACSTRUCT.TEE_INFO_HASH equals the hash of the TDINFO_STRUCT. Then, software uses

¹³Intel, *Architecture Specification: Intel® Trust Domain Extensions (Intel® TDX) Module*, 2020 

TD Report Structures (Simplified view)



TD Report Structures¹⁴

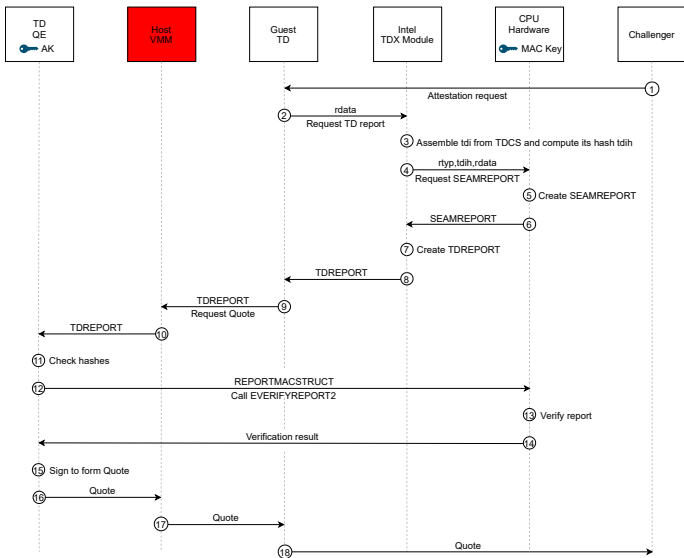


¹⁴Sardar, Musaeu, and Fetzer, "Demystifying Attestation in Intel Trust Domain Extensions via Formal Verification", 2021

Simplified View of Protocol

- Local attestation → Symmetric crypto → MAC
- Remote attestation → Asymmetric crypto → Digital signatures

TDX Attestation Flow for Quote Generation¹⁵



¹⁵Sardar, Musaeu, and Fetzer, "Demystifying Attestation in Intel Trust Domain Extensions via Formal Verification", 2021

Automated Verification

- Validation: reachability of all parts of code
- Confidentiality: reachability property
- Authentication properties, e.g.,
 $x \equiv \langle rtyp, res1, csvn, tcbh, tdih, rdata, res2 \rangle$

$\forall x.$

$\exists mac, tcbi.$

$event(QuoteVerified(x)) \Rightarrow event(CPUsentSMR(x, mac, tcbi))$

- 1 Introduction
- 2 Formal Security Analysis Approach
- 3 TDX
 - Discrepancies Identified
 - Formal Specification
 - Automated Verification
- 4 Summary

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- Shameless plug: we are hiring PhDs, post-docs
([muhammad_usama.sardar,christof.fetzer]@tu-dresden.de)

Key References I



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Key References II



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