

# Reasoning about Deltas — Even Doing Nothing is Difficult

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## Takeaway

- Problem: reasoning about specifications deltas is difficult

- Need: automated proof engineering

- Solution: typeCart, an analysis tool



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  - Compares two versions of specifications
  - Generates boilerplate code relating the two versions
  - Semi-automates reasoning about two versions of specifications



#### Takeaway

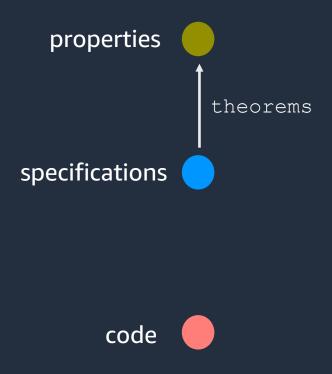
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- Application: relating two specification versions of AWS authorization engine



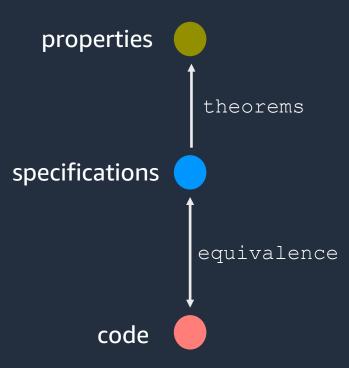
specifications

code

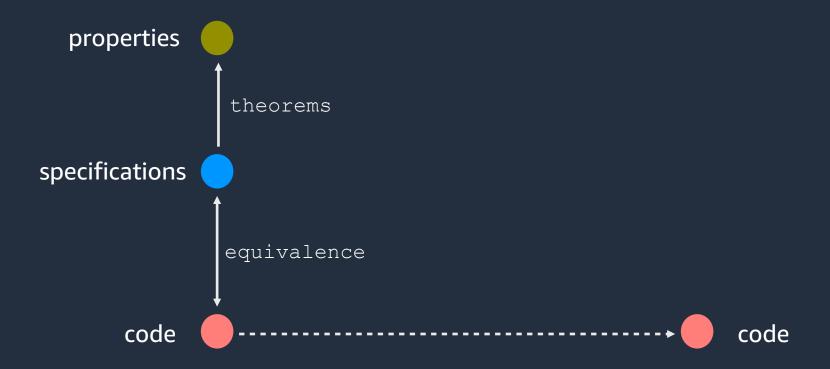




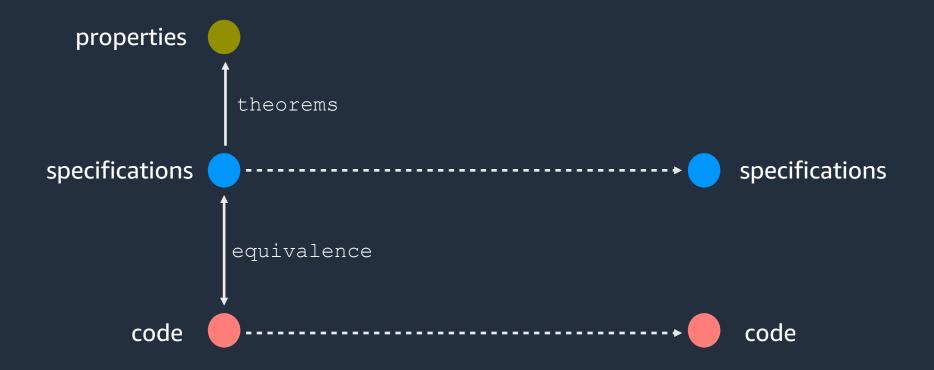




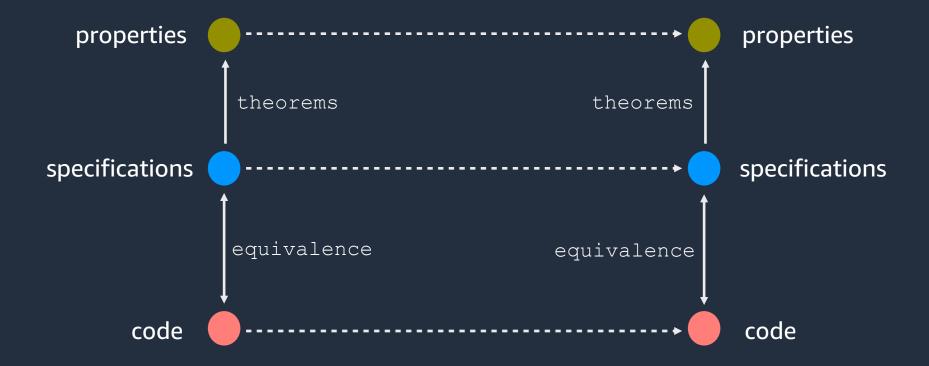




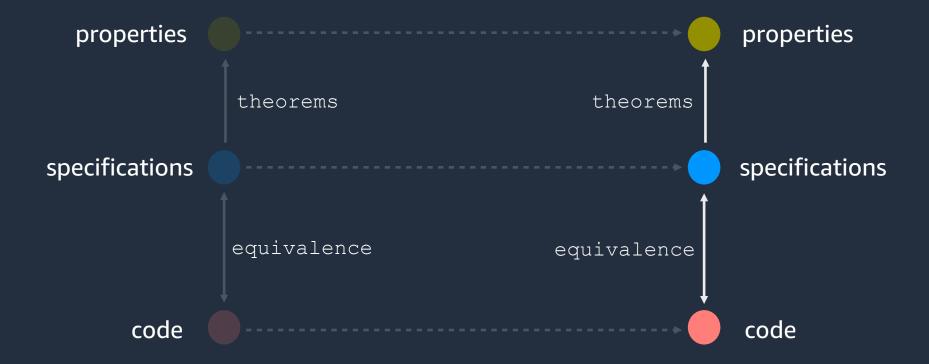




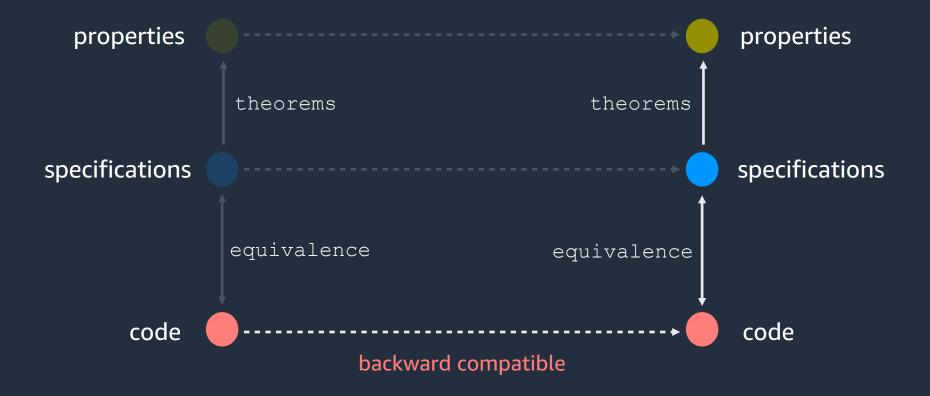




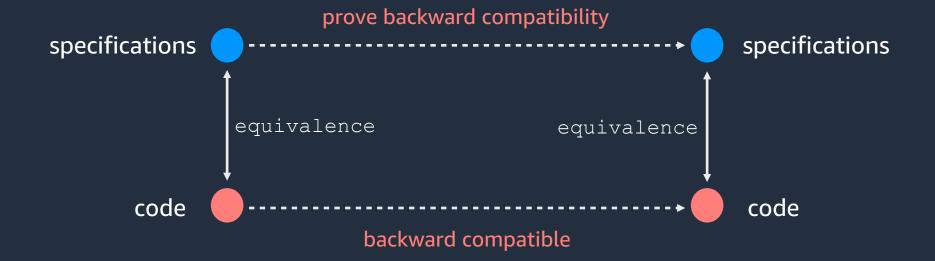












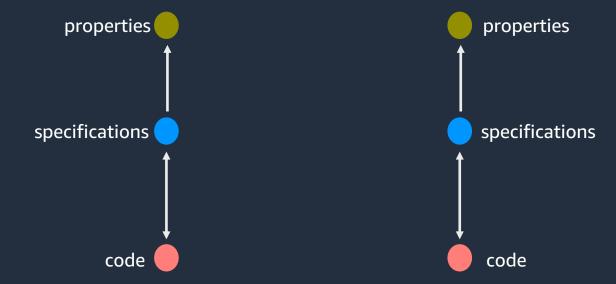


- Verifying a system once is already an achievement



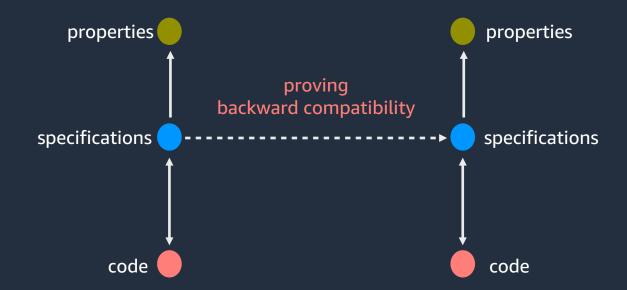


- Verifying a system once is already an achievement
- Maintaining verification artefacts is already a challenge



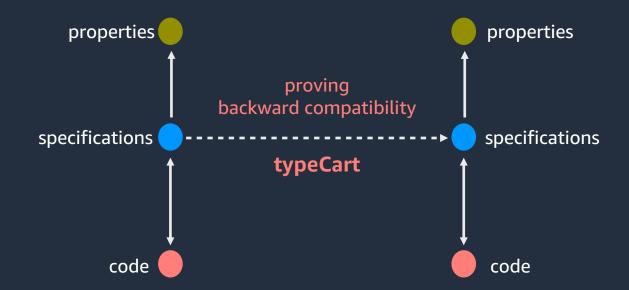


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- Verifying a system once is already an achievement
- Maintaining verification artefacts is already a challenge
- Proving backward compatibility is an add-on
- typeCart: automates the process of relating specifications







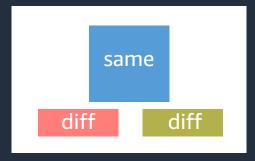
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  - The two versions represent the same implementation with some deltas
  - They model and specify relatable objects







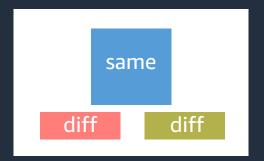
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- Relating equivalent objects between specifications should be trivial



version 1 and version 2



- Relating equivalent objects between specifications should be trivial
- Let's test our hypothesis!



version 1 and version 2



```
module Spec {
  const INT_MAX := 0x7fff_ffff
  newtype nat32 = x | 0 <= x <= INT_MAX

  datatype exp = Const (nat32) | Add (nat32,nat32)

  function eval (e:exp):(v:nat32)
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```

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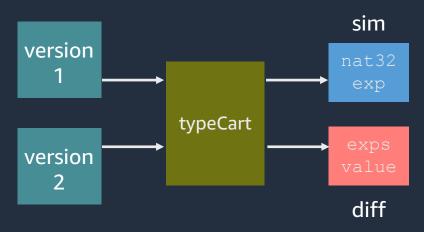
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backward compatibility

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#### both `eval` functions evaluate to the same answer

```
module SpecRel{
  import Spec
  import Spec
  // Duplicate name of import
```



- Relating equivalent objects between specifications should be trivial

```
module Old.Spec {
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namespaces for different versions

```
module SpecRel{
import Old
import New
```



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  }
}
```

```
module SpecRel{
  import Old
  import New

lemma trivial:
   e:Old.Spec.exp
  New.Spec.eval(e) = Old.Spec.eval(e)

  // typecheck fails: New.Spec.eval expects New.Spec.exp
```



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  }
}
```

type casting functions

```
module SpecRel{
  import Old
  import New

lemma trivial:
   New.Spec.eval(exp0ldtoNew(e)) = Old.Spec.eval(e)
  // typecheck fails: New.Spec.nat32 and Old.Spec.nat32
```



- Relating equivalent objects between specifications should be trivial

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type casting functions

```
module SpecRel{
  import Old
  import New

lemma trivial:
   New.Spec.eval(exp0ldtoNew(e)) = nat320ldtoNew(Old.Spec.eval(e))
  // typecheck succeeds
```



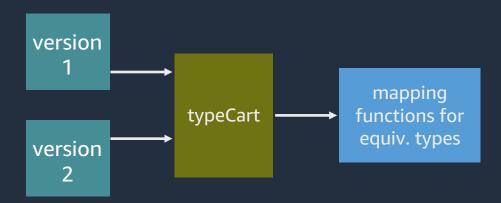
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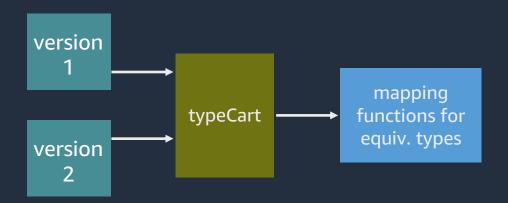


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  - Identifies syntactically equivalent types between Dafny files
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\*Fun Fact\*

**typeCart:** name-blend of **type** and **cartography** (practice of making maps)



#### - Recursive datatype

```
datatype recSimple =
    A
    B(b: recSimple)
```

```
function recSimpleOldToNew(r: Old.Spec.recSimple): New.Spec.recSimple
{
    match r
    case A =>
        New.Spec.recSimple.A
    case B(b: Old.Spec.recSimple) =>
        New.Spec.recSimple.B(recSimpleOldToNew(b))
}
```



Recursive datatype

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datatype recSimple =
   A
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```

```
function recSimpleOldToNew(r: Old.Spec.recSimple): New.Spec.recSimple
```

```
datatype refSimple =
   A(a: int)
   B(b: recSimple)
```

```
function refSimpleOldToNew(r: Old.Spec.refSimple): New.Spec.refSimple
{
    match r
    case A(a: int) =>
        New.Spec.refSimple.A(a)
    case B(b: Old.Spec.recSimple) =>
        New.Spec.refSimple.B(recSimpleOldToNew(b))
}
```



- Parametric types

```
datatype either<S, T> =
  Left(s: S)
  Right (t: T)
```

```
function eitherOldToNew<S, T, S', T'>
  (fS: S -> S', fT: T -> T', e: Old.Spec.either<S, T>):
  New.Spec.either<S', T'>
{
  match e
  case Left(s: S) =>
    New.Spec.either.Left(fS(s))
  case Right(t: T) =>
    New.Spec.either.Right(fT(t))
}
```



- Bounded types

```
const INT_MAX := 0x7fff_ffff
newtype nat32 = x | 0 <= x <= INT_MAX</pre>
```

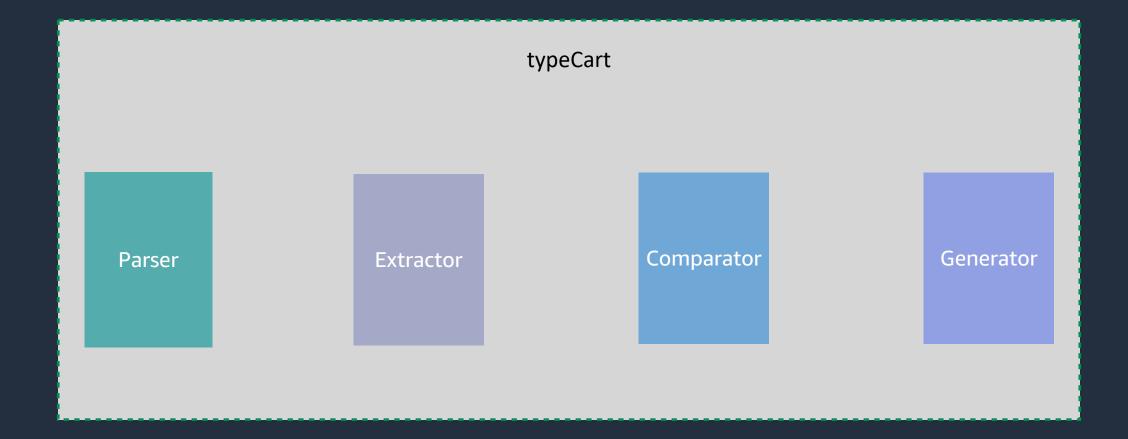
```
function nat320ldToNew(n: Old.Spec.nat32): (n': New.Spec.nat32)
  ensures n as int == n' as int
{
    n as int as New.Spec.nat32
}
```



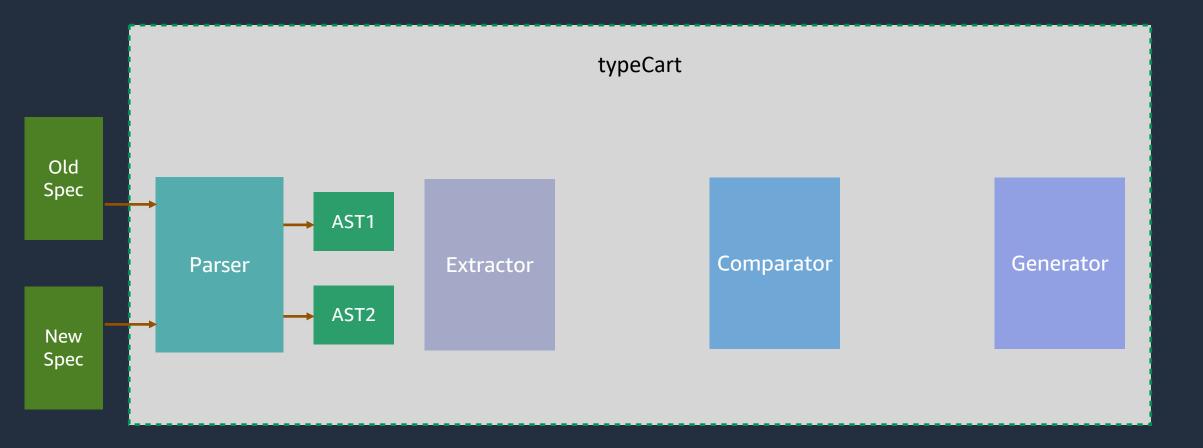
- Collection types

```
datatype collectionType<T> =
   A(a: T)
   B(b: seq<T>)
```

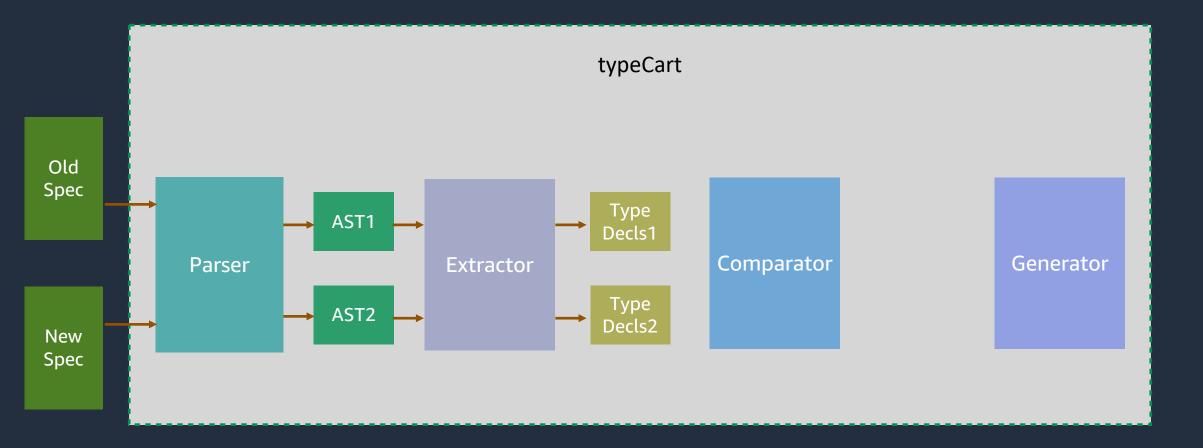




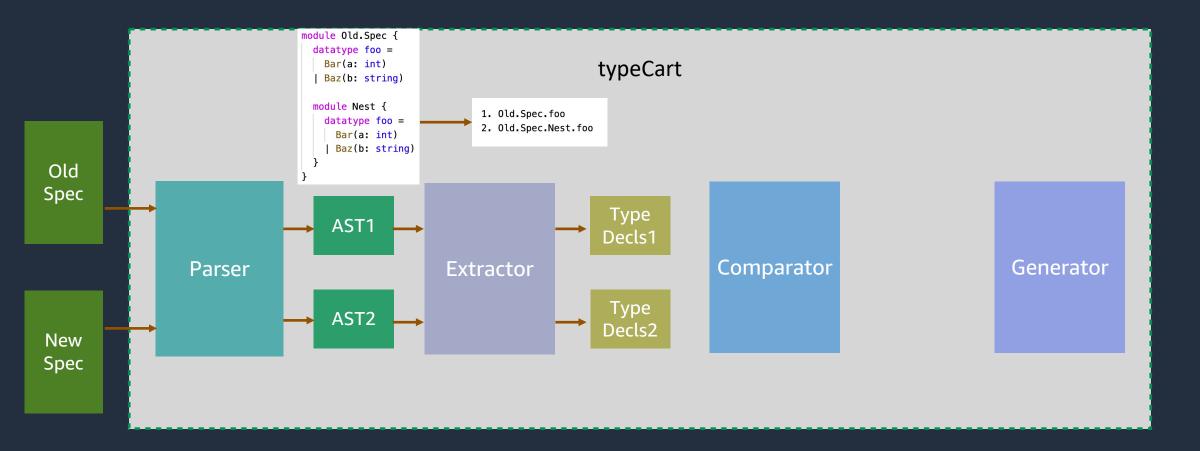




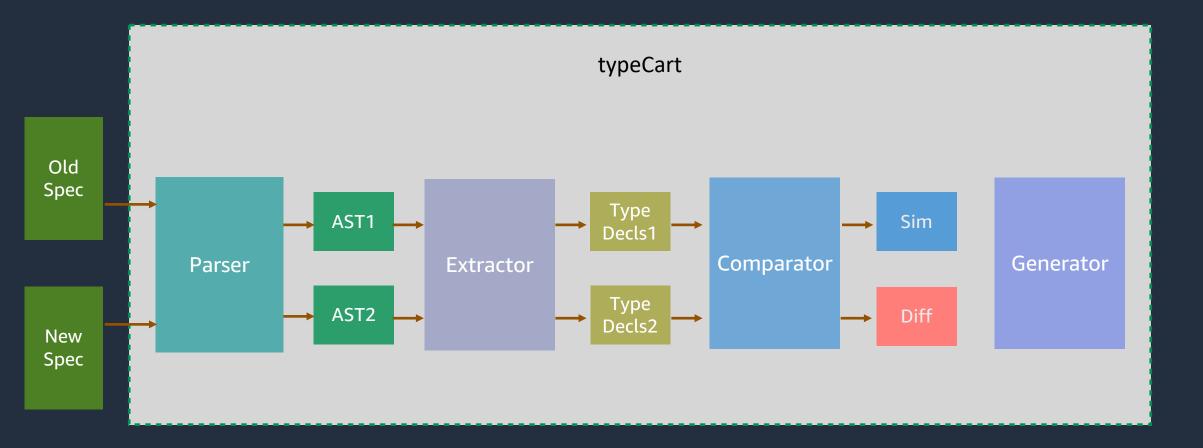




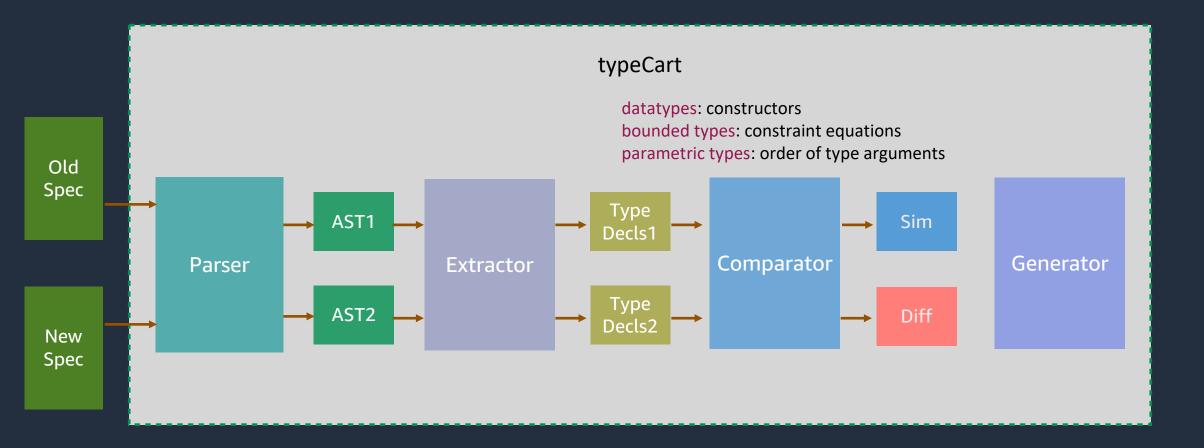




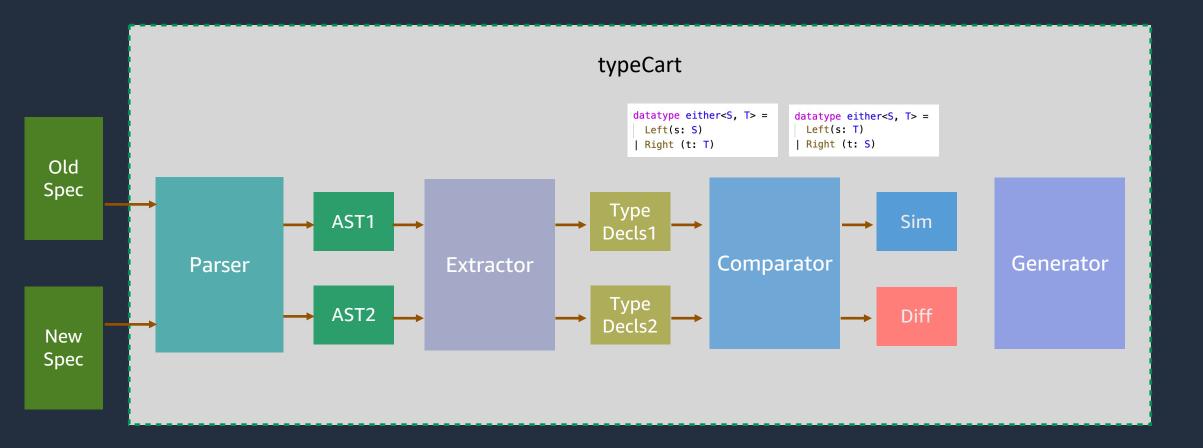




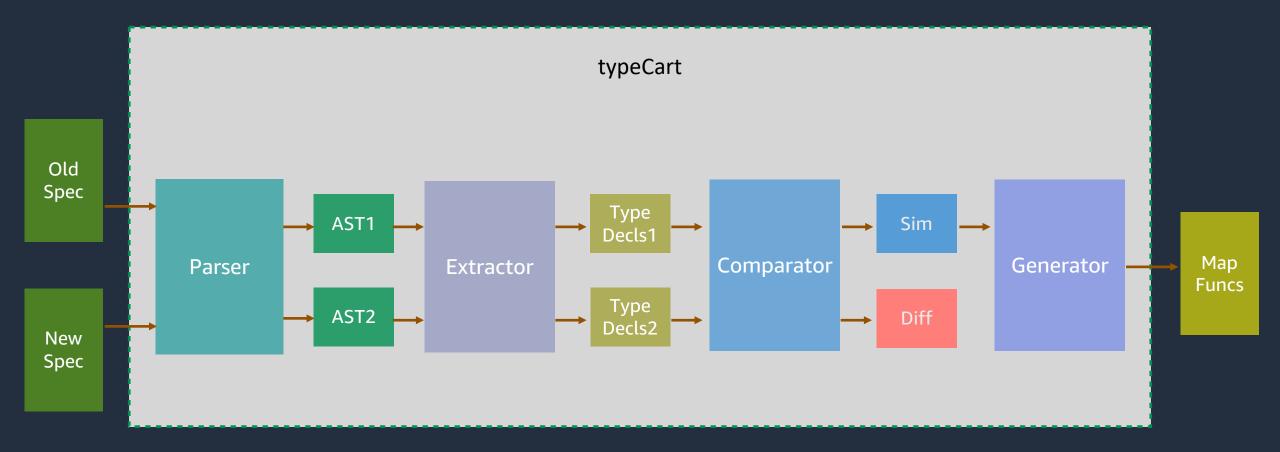






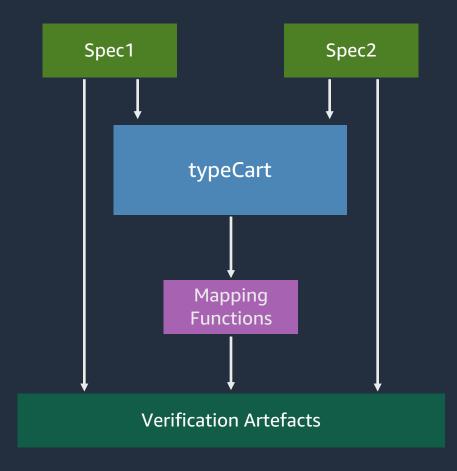








## Integrating typeCart to Verification





#### Recap

```
datatype exp = Const (nat32) | Add (nat32,nat32)
function eval (e:exp):(v:nat32)
```

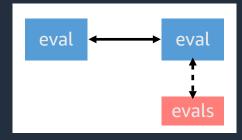
#### version1

```
datatype exps = One (exp) | Struct (seq<exp>)

datatype value = Val (nat32) | Vals (seq<nat32>)

function evals (es:exps):(v:value)
```

version2



backward compatibility



#### Recap

```
datatype exp = Const (nat32) | Add (nat32,nat32)

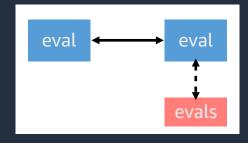
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backward compatibility

version2

Doing nothing is was difficult: typeCart





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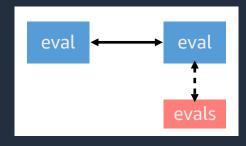
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backward compatibility

- version2
- Doing nothing is was difficult: typeCart
- Let's do something!
  - Reasoning about deltas







- Specifications deltas involve evolving types
- Evolving types require mapping relations instead of mapping functions



- Specifications deltas involve evolving types
- Evolving types require mapping relations instead of mapping functions

```
datatype expr = Const(x: int) | Add(e1: expr, e2: expr) | Sub(e1: expr, e2: expr)
// one added, one deleted constructor
datatype expr = Const(x: int) | Add(e1: expr, e2: expr) | Mul(e1: expr, e2: expr)
```

```
// relation between old and new type (same name as the type)
function expr(e0: Old.expr, eN: New.expr): bool
{
    match (e0,eN) {
      case (Const(n0),Const(nN)) => n0 == nN
      case (Add(x0, y0), Add(xN,yN)) => expr(x0,xN) && expr(y0,yN)
      case _ => false
    }
}
```



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```

#### upcoming typeCart V2:

- will identify evolving+equivalent types
- will generate mapping relations



- Type constructors with higher order arguments require nested translation



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```
datatype foo<A> = Foo (a:A)

datatype bar<A> = Bar (b:foo<foo<A>>)
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- Generic interface
  - TIL: typeCart Intermediate Language



- Generic interface
  - TIL: typeCart Intermediate Language
- Revision control integration

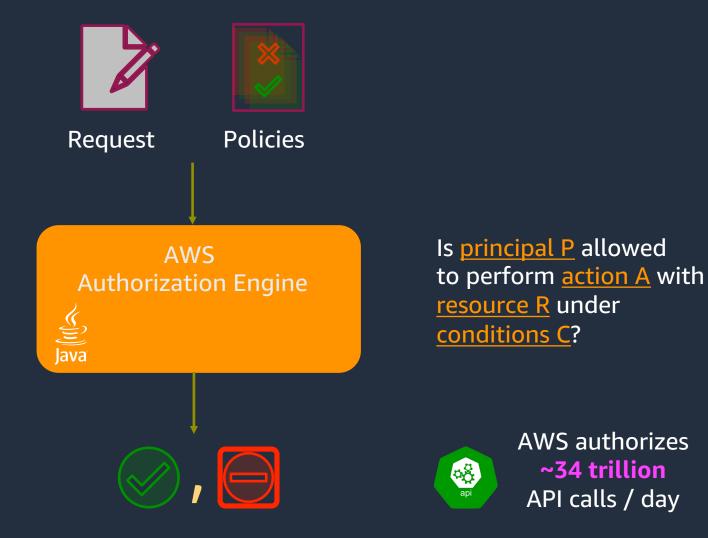






Is <u>principal P</u> allowed to perform <u>action A</u> with <u>resource R</u> under <u>conditions C?</u>







- Changes to the Authorization Engine
  - Performance updates
  - Feature updates
  - Algorithmic updates



- Changes to the Authorization Engine
  - Performance updates
  - Feature updates
  - Algorithmic updates

#### Engineers seek answers to:

- How does this change affect the existing implementation?
- Is it backward compatible?
- Does this change affect the existing authorization strategy?



# Yucca — Verified AWS Authorization Engine

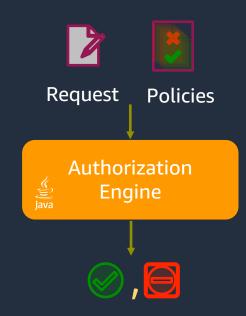
- The Yucca project started in the spring 2020
- Goals
  - Have a high level specifications of the Authorization Engine
  - Prove authorization properties
  - Prove backward compatibility going forward





# Yucca — Verified AWS Authorization Engine

- The Yucca project started in the spring 2020
- Goals
  - Have a high level specifications of the Authorization Engine
  - Prove authorization properties
  - Prove backward compatibility going forward
    - Apply typeCart to semi-automate the process





### Can You use typeCart?

- Absolutely yes!



- Want to analyze the effect of introducing changes to your Dafny development?
  - typeCart is here



### Can You use typeCart?

- Absolutely yes!



- Want to analyze the effect of introducing changes to your Dafny development?
  - typeCart is here
- Questions!

